

Cut and Projection - Part II

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CNRS & Univ. Paris 13

CIMPA School
Tilings and Tessellations
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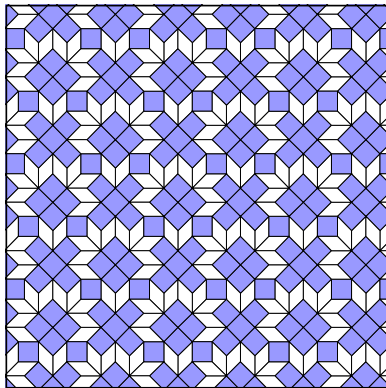
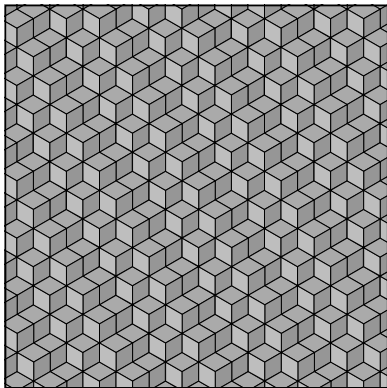
Outline

- 1 Local rules
- 2 Sufficient conditions
- 3 Necessary conditions
- 4 Colored local rules

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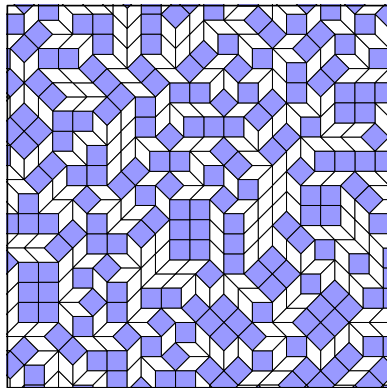
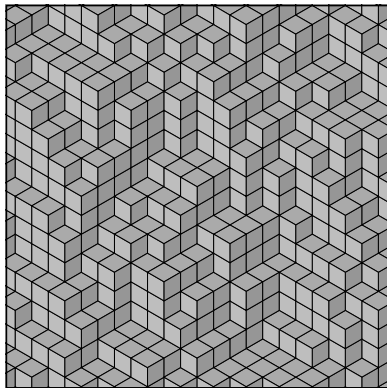
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General $n \rightarrow d$ tilings



Planar tilings are well ordered...

General $n \rightarrow d$ tilings



Planar tilings are well ordered... but they can easily be messed up!

Local rules

Definition (Local rules)

A planar tiling of slope E has *diameter* r and *thickness* t *local rules* if any tiling with a smaller or equal r -atlas lifts into $E + [0, t]^n$.



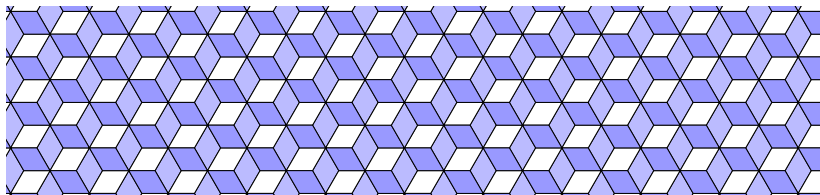
Main Open Question

Which planar tilings do admit local rules?

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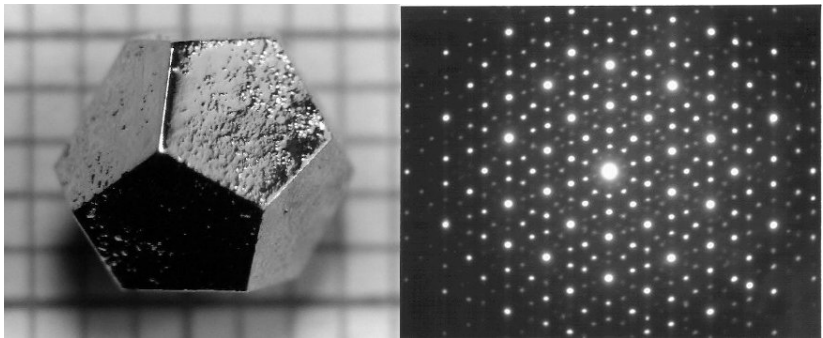


Main Open Question

Which planar tilings do admit local rules?

Link with quasicrystals

Planar $n \rightarrow d$ tilings aim to model the *structure* of *quasicrystals*.



Local rules aim to model their *stability* (i.e., energetic interactions).

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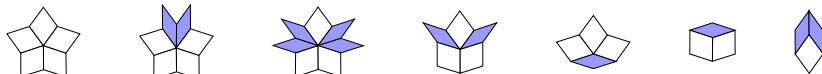
Penrose tilings

Definition (Penrose tiling)

A *Penrose tiling* is a planar $5 \rightarrow 2$ tiling with slope

$$\frac{1}{5}(1, 1, 1, 1, 1) + \mathbb{R} \left(\cos \frac{2k\pi}{5} \right)_{0 \leq k \leq 4} + \mathbb{R} \left(\sin \frac{2k\pi}{5} \right)_{0 \leq k \leq 4}.$$

It is the dualization of the multigrid with vectors $e^{\frac{2ik\pi}{5}}$ and shifts $\frac{1}{5}$.



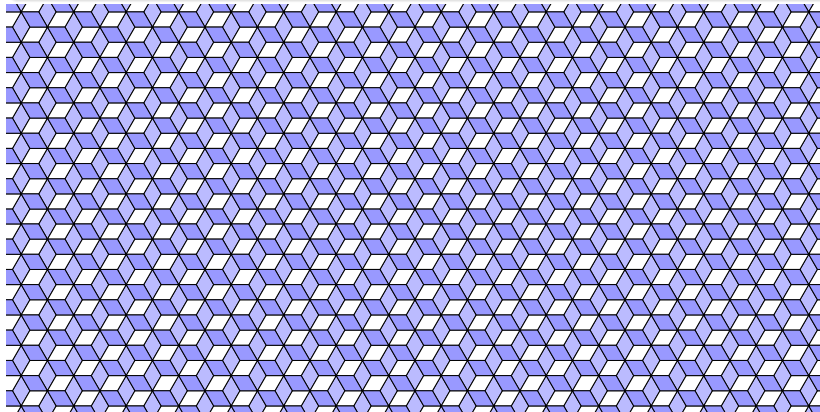
Theorem (de Bruijn, 1981)

Penrose tilings have local rules of diameter 0 and thickness 1.

n -fold tilings

Definition (n -fold tiling)

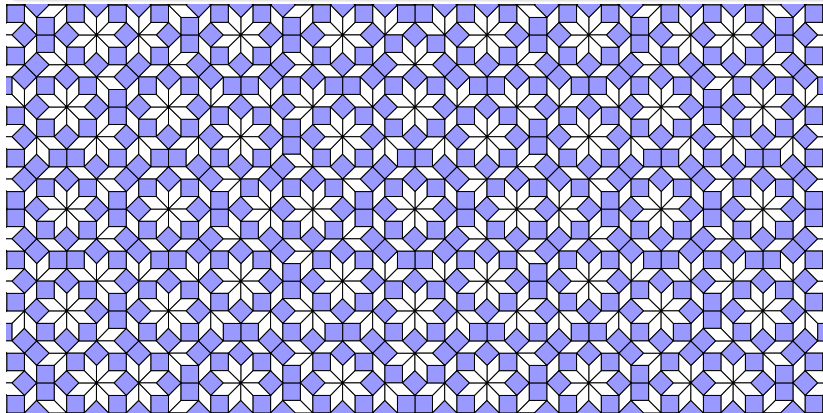
A n -fold tiling is a planar $n \rightarrow 2$ tiling which has the same finite patterns as its image under a rotation by $2\pi/n$.



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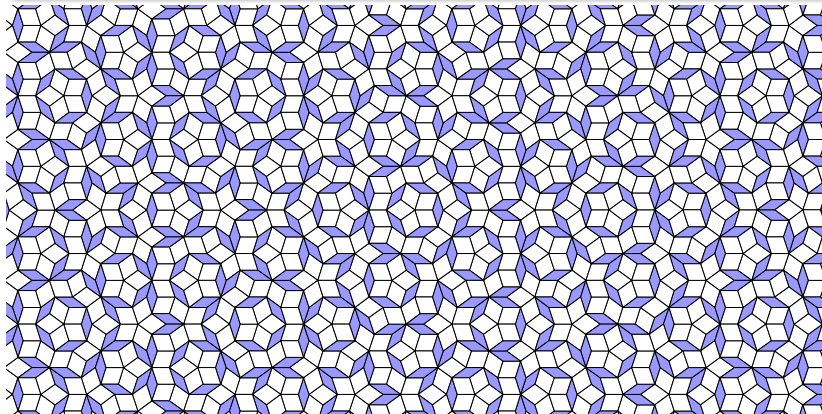
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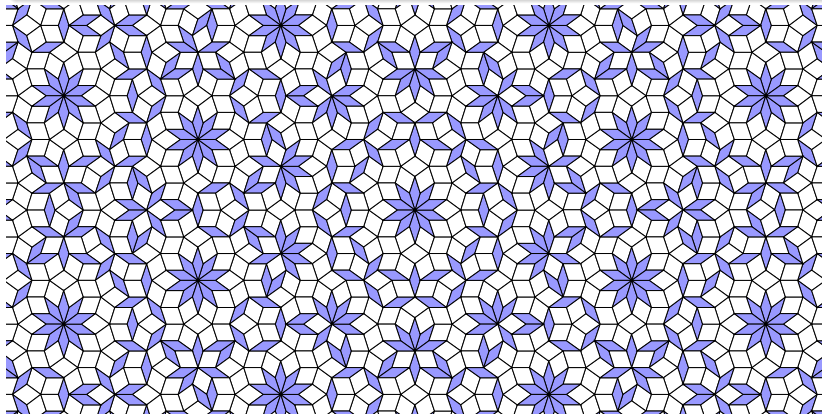
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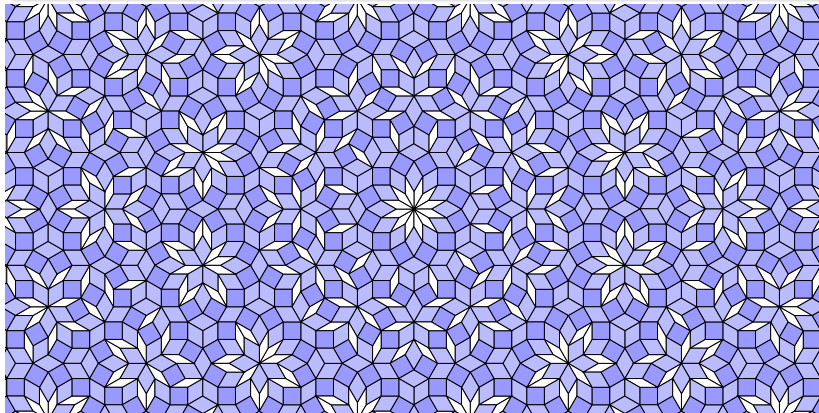
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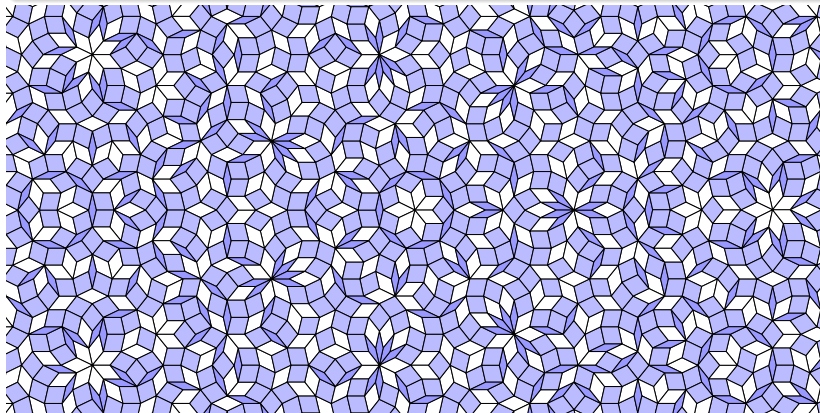
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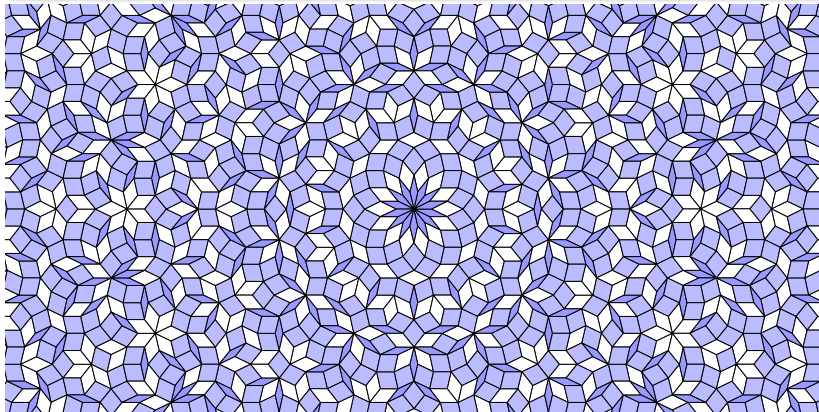
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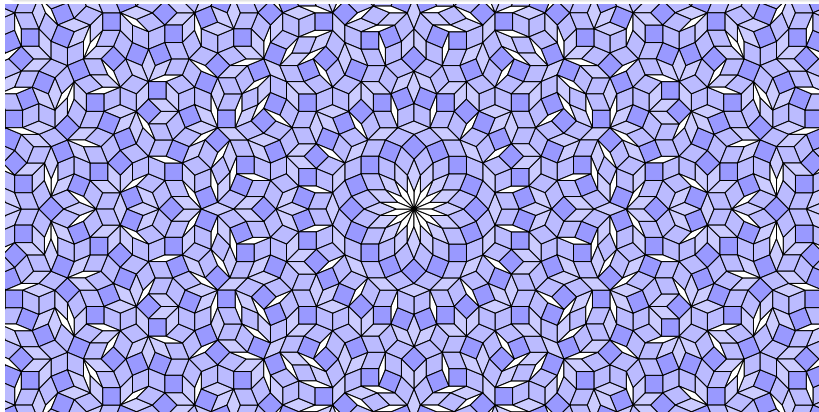
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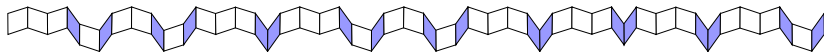


Local rules for n -fold tilings

Theorem (Socolar 1990)

An n -fold tiling has local rules when n is not a multiple of 4.

Local rules actually enforce an *alternation condition*:

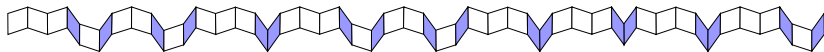


Local rules for n -fold tilings

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Local rules actually enforce an *alternation condition*:



When n is a multiple of 4, there are square tiles...

Subperiods

Definition (Subperiod)

A planar $n \rightarrow d$ tiling has a *subperiod* if one gets a periodic tiling by an orthogonal projection onto $d + 1$ well-chosen basis vectors.

For example, a Penrose tiling has 10 subperiods (video).

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For example, a Penrose tiling has 10 subperiods (video).

This translates in linear rational dependencies between Grassmann coordinates over $d + 1$ indices. For Penrose:

$$G_{12} = G_{23} = G_{34} = G_{45} = G_{51}, \quad G_{13} = G_{35} = G_{52} = G_{24} = G_{41}.$$

Planarity issues

Proposition

The subperiods of a planar tiling can be enforced by local rules.

But these local rules may not suffice to enforce planarity. . .

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Theorem (Bédaride-Fernique 2015)

A planar $4 \rightarrow 2$ tiling has local rules iff its slope is characterized by its subperiods. In particular the slope is quadratic (or rational).

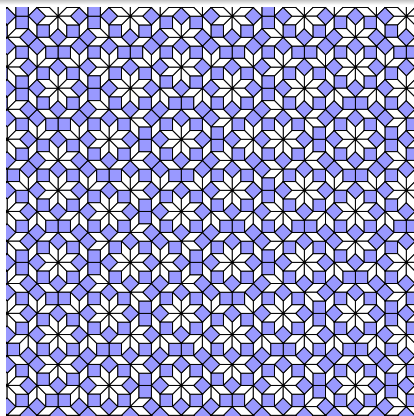
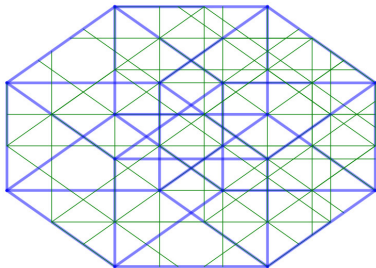
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$4p$ -fold tilings

Theorem (Bédaride-Fernique 2015)

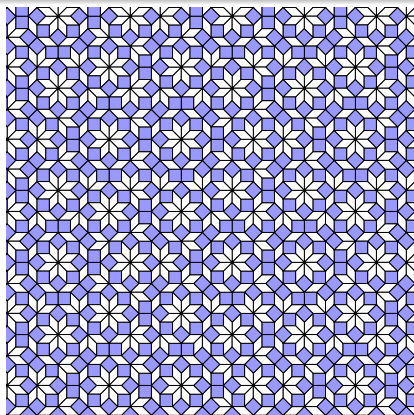
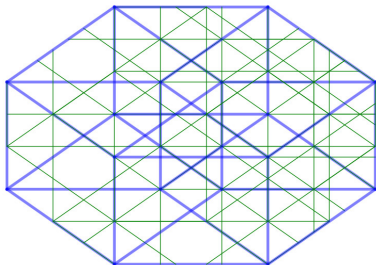
The $4p$ -fold tilings do not have local rules.



$4p$ -fold tilings

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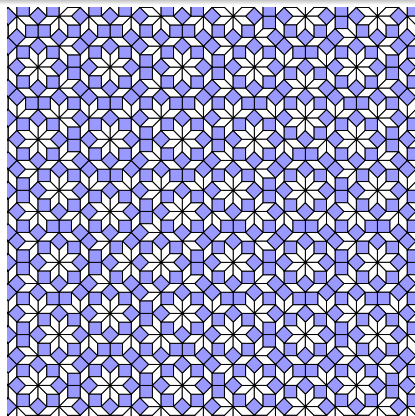
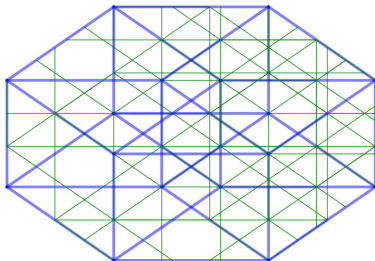
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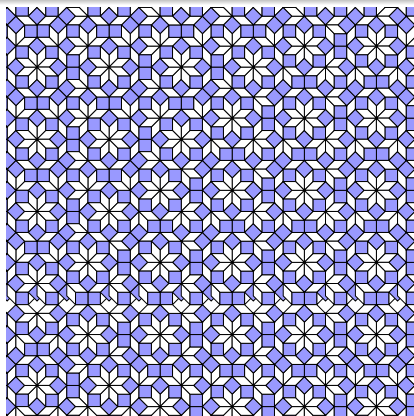
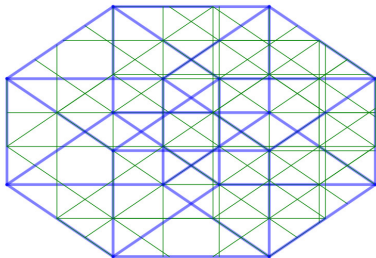
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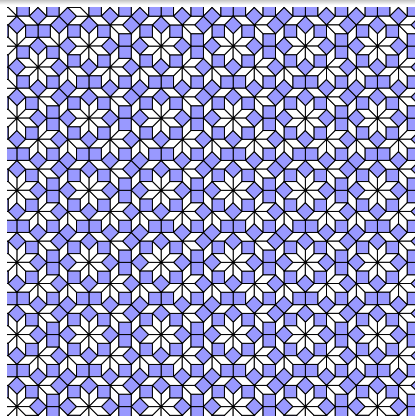
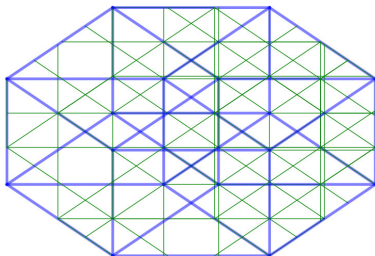
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Local rules of thickness 1

Full subperiods: any projection on $d + 1$ basis vector is periodic.

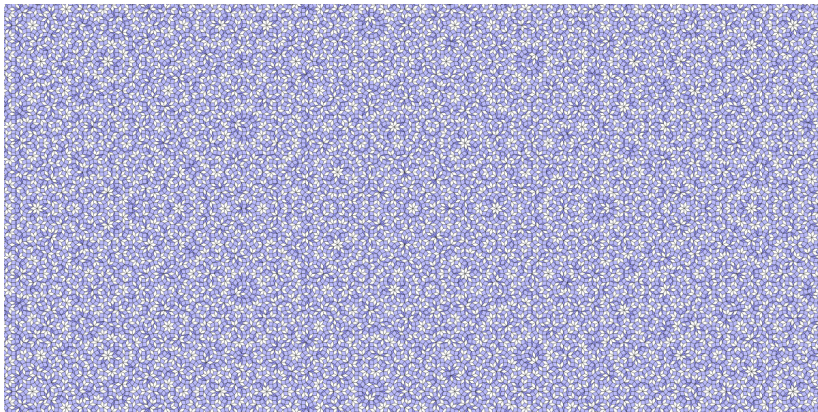
Theorem (Levitov 1988)

A planar tiling with thickness 1 local rules has full subperiods.

For n -fold tilings, this yields $n \in \{4, 6, 8, 10, 12\}$.

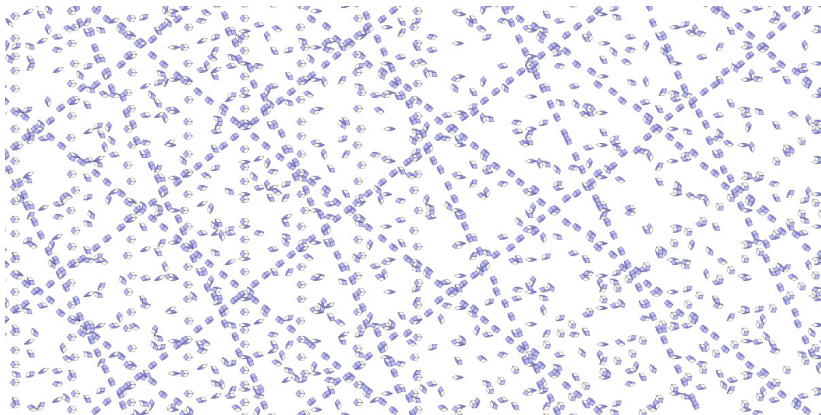
These are the only symmetries yet observed in real quasicrystals. . .

Proof sketch



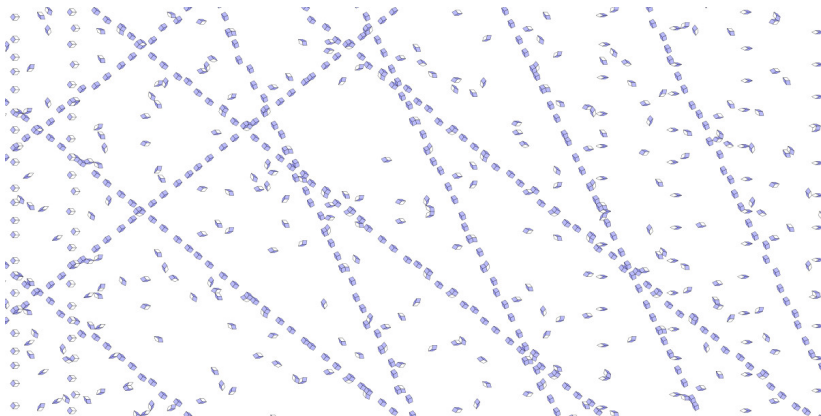
Consider a planar tiling which does not have full subperiods.

Proof sketch



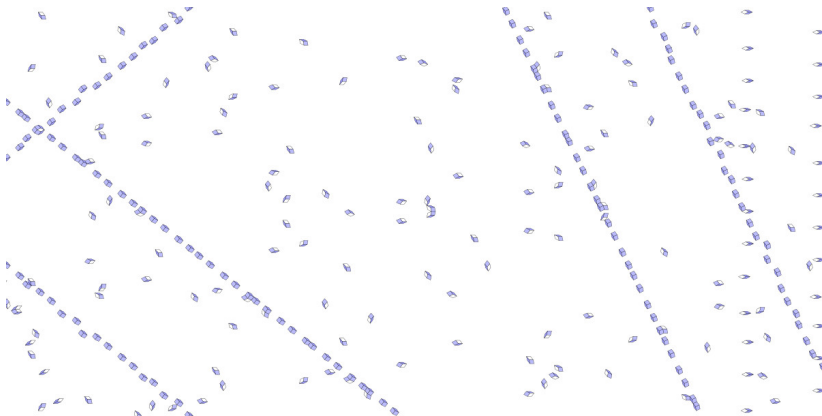
Shifting the slope creates *flips*. We shift without creating patterns.

Proof sketch



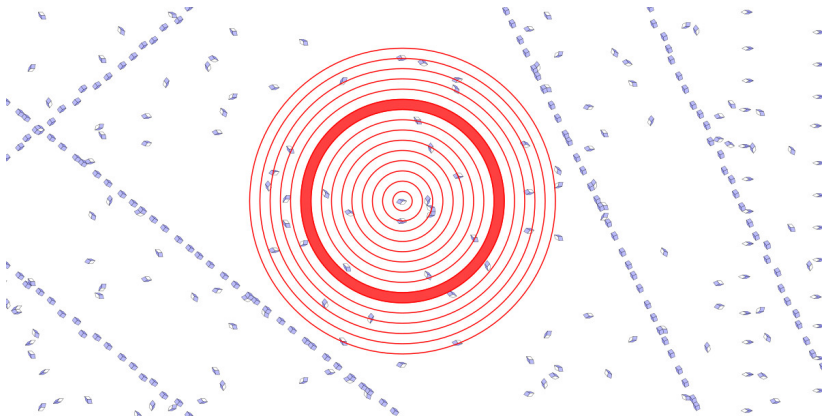
There are lines of flips (corresp. to subperiods) and isolated flips.

Proof sketch



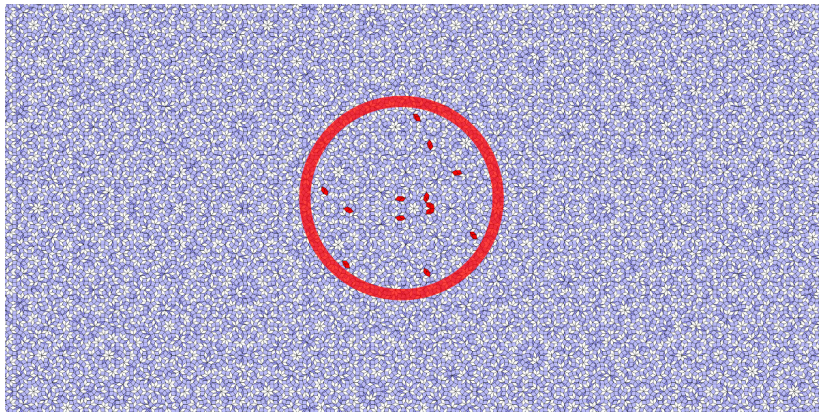
The smaller the shift is, the sparser these flips are.

Proof sketch



Given r , we eventually find a ring of thickness r without any flip.

Proof sketch

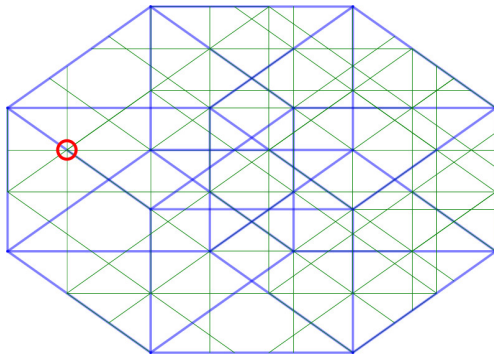


This yields a planar tiling of thickness $t > 1$ with the same r -atlas.

Algebraic obstruction

Theorem (Le 1995)

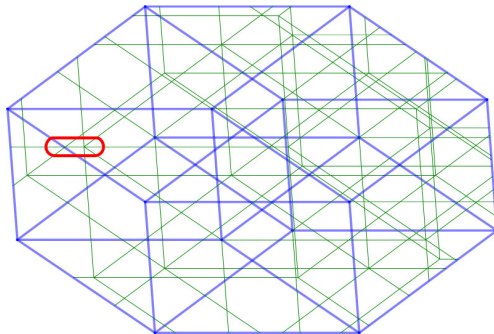
The slope of a planar tiling with local rules is algebraic.



Algebraic obstruction

Theorem (Le 1995)

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From colors to computability

Definition (Colored local rules)

A planar tiling has *colored local rules* if it is obtained by removing the colors of a colored planar tilings which has local rules.

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This allows to go far beyond the previous algebraic obstruction:

Theorem (Fernique-Sablik 2012)

A planar tiling has colored local rules iff its slope is computable.

We would like such a characterization for uncolored local rules. . .

Computable slope?

Definition

A number is *computable* if it can be computed to within any desired precision by a finite, terminating algorithm.

Q. Are rational numbers computable? Algebraic numbers? e ? π ?

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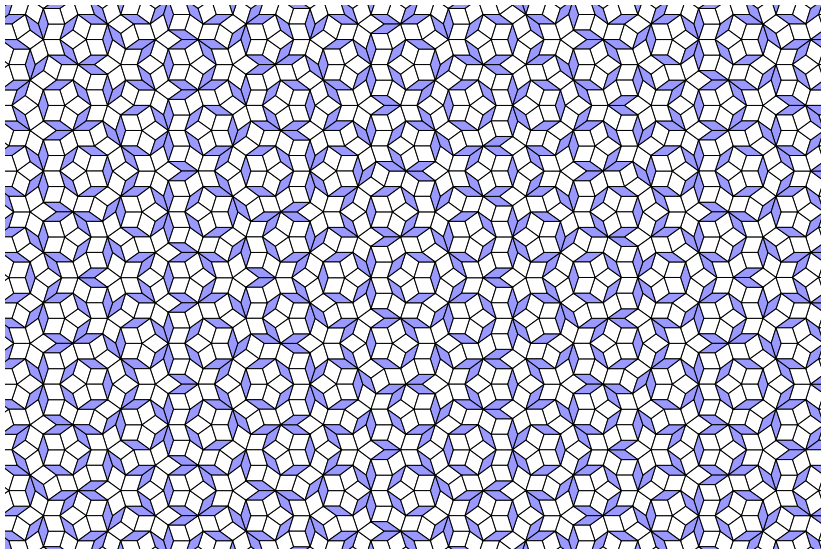
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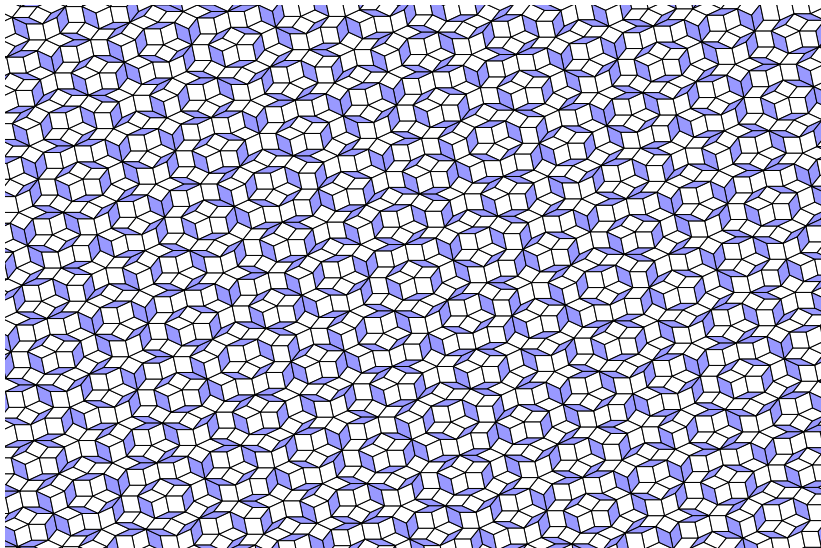
Q. Do there exist any non-computable number!?

A slope is computable if it has computable Grassmann coordinates.

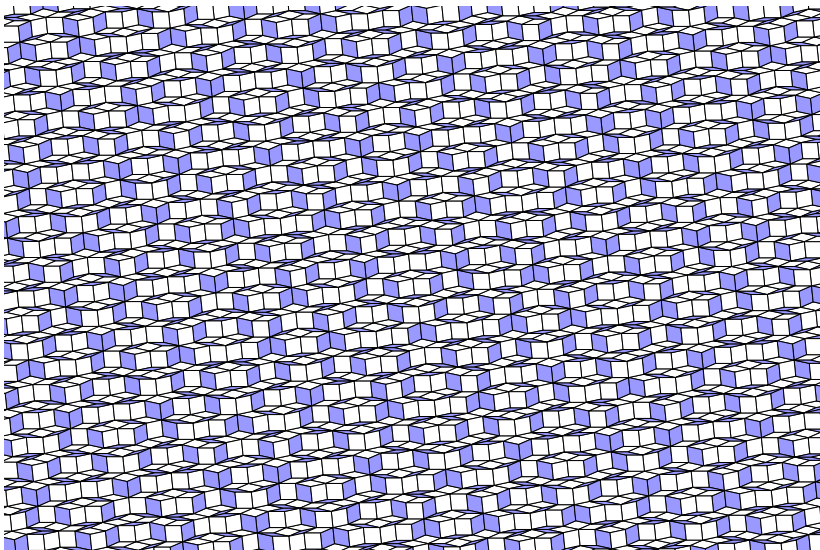
Proof sketch



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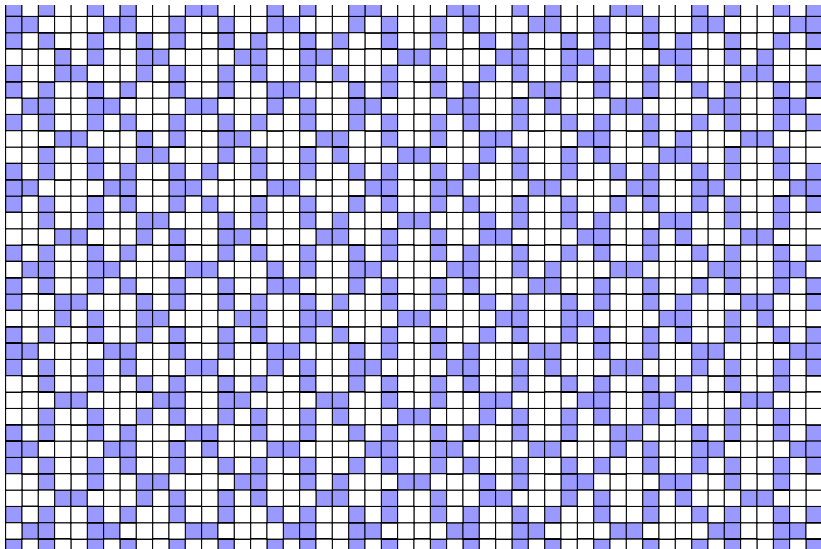
Local rules
ooo

Sufficient conditions
ooooo

Necessary conditions
oooo

Colored local rules
oo●

Proof sketch



The End