

A Self-Simulating Tileset

Thomas Fernique

Goal

We define step by step

- ▶ a finite set τ of decorated squares, where every edge is endowed with a (red,green,blue) triple of indices,
- ▶ a bijection ϕ from 3×3 squares of τ -tiles to τ -tiles,

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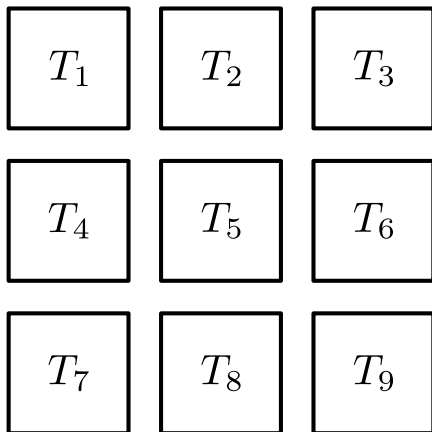
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such that:

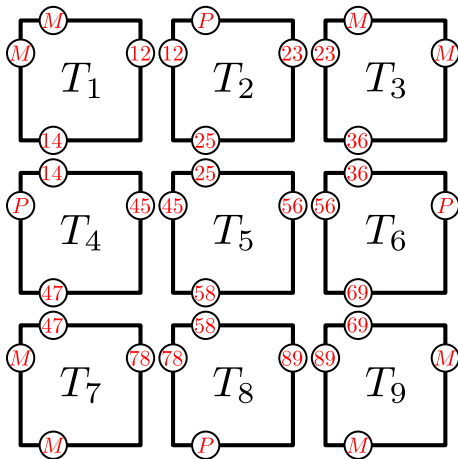
- ▶ there exist τ -tilings,
- ▶ every τ -tiling can be uniquely partitioned into 3×3 squares,
- ▶ applying ϕ on these 3×3 squares (+scaling) yields a τ -tiling.

Step 1: macro-tiles



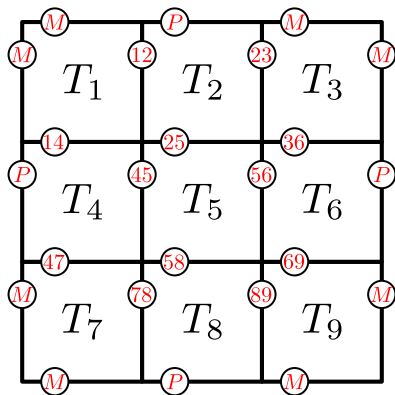
Start with tiles T_1, \dots, T_9 .

Step 1: macro-tiles



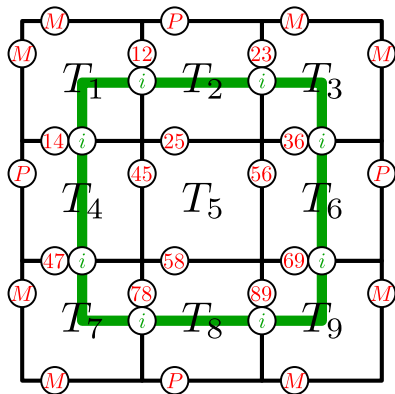
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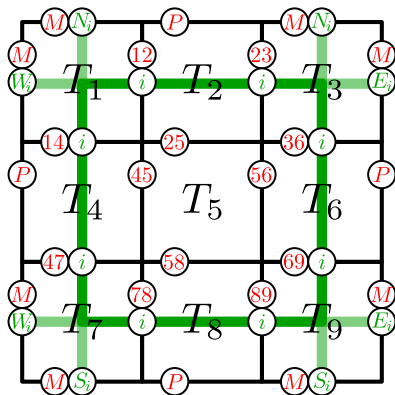
Start with tiles T_1, \dots, T_9 . Endow each edge with a red index enforcing the tiles to assemble into *macro-tiles* aligned along a grid.

Step 2: rings and macro-macro-tiles



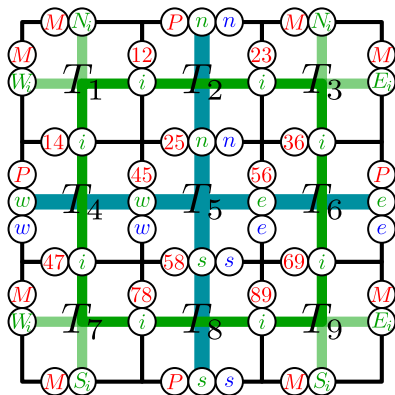
A green index $i \in \{1, \dots, 9\}$ runs along a *ring* in every macro-tile.

Step 2: rings and macro-macro-tiles



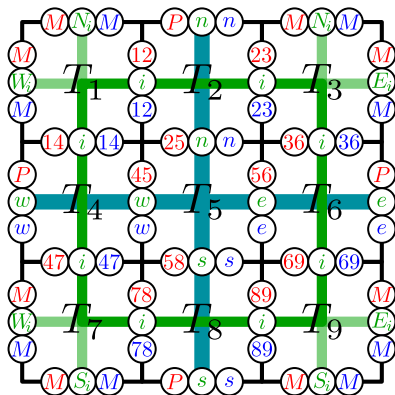
Further green indices force rings to order as T_i 's in a macro-tile.
(X_i denotes the red index on the X -edge of T_i , $X = N, W, S, E$)

Step 3: the network



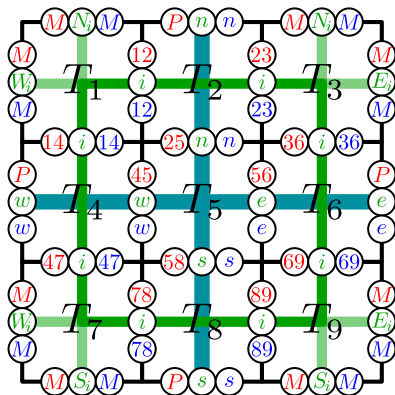
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Step 3: the network



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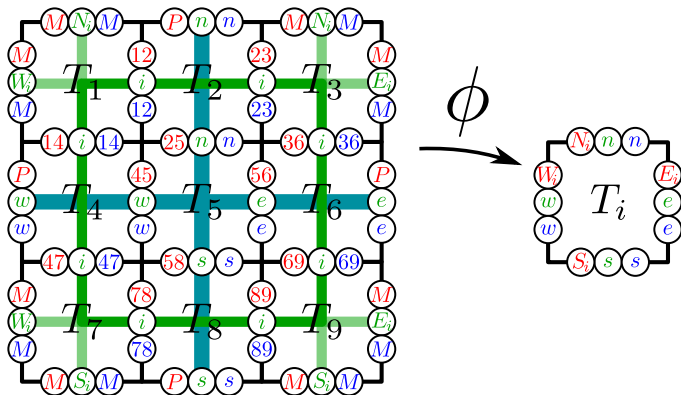
Step 3: the network



Duck Lemma

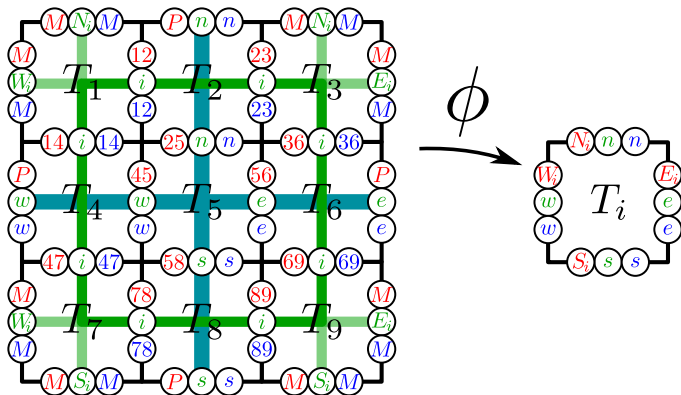
The **green/blue** indices of every $T_i \neq T_5$ determine the **red** indices.

Interlude: the map ϕ



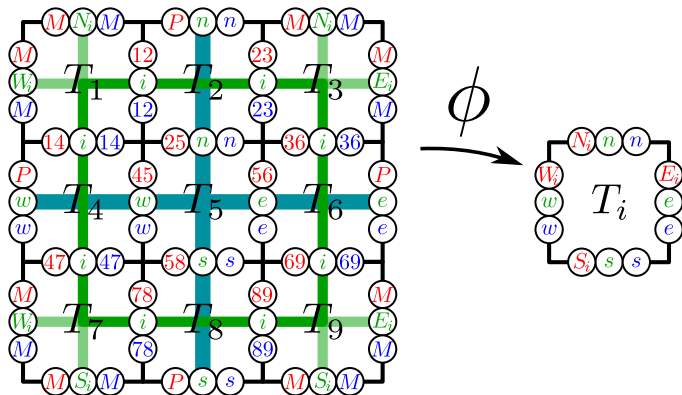
ϕ maps the green indices forcing ring ordering onto red indices and copies the green/blue indices from network branches.

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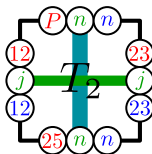


ϕ maps a tiling onto a tiling. Does it map a τ -tiling onto a τ -tiling?
 The indices on the **network** will have to be chosen so that it holds.

Step 4: ring/network intersections

When a **ring** along which runs j crosses an **X-branch**, it checks that the pair carried by the branch is allowed on the X-edge of T_j .

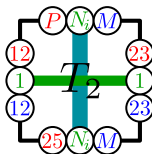
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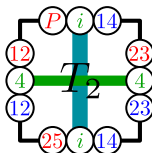


- ▶ $j \in \{1, 3\} \rightsquigarrow \{M, P, 14, 25, 36, 47, 58, 69\} \times \{M\};$

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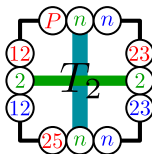


- ▶ $j \in \{1, 3\} \rightsquigarrow \{M, P, 14, 25, 36, 47, 58, 69\} \times \{M\};$
- ▶ $j \in \{4, 6, 7, 9\} \rightsquigarrow \{1, \dots, 9\} \times \{N_j\};$

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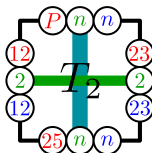


- ▶ $j \in \{1, 3\} \rightsquigarrow \{M, P, 14, 25, 36, 47, 58, 69\} \times \{M\}$;
- ▶ $j \in \{4, 6, 7, 9\} \rightsquigarrow \{1, \dots, 9\} \times \{N_j\}$;
- ▶ $j \in \{2, 5, 8\} \rightsquigarrow$ every pair already defined on any North-edge!

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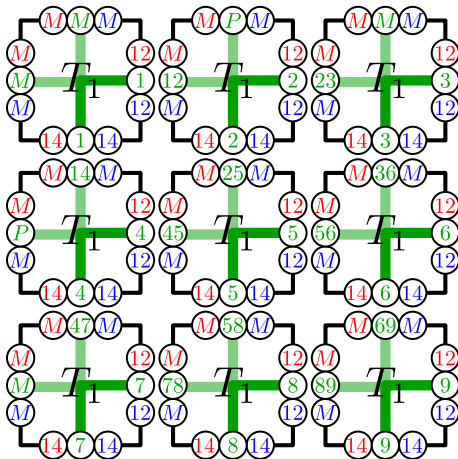
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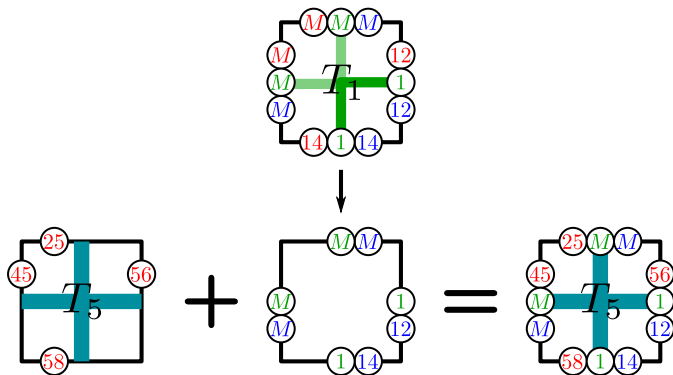
This yields $2 \times 8 + 4 \times 9 + 3 \times (2 \times 8 + 4 \times 9) = 208$ decorated T_2 .
Together with $T_{4,6,8}$ (208 each) and $T_{1,3,7,9}$ (9 each): 868 tiles.

Step 5: synchronizing network branches



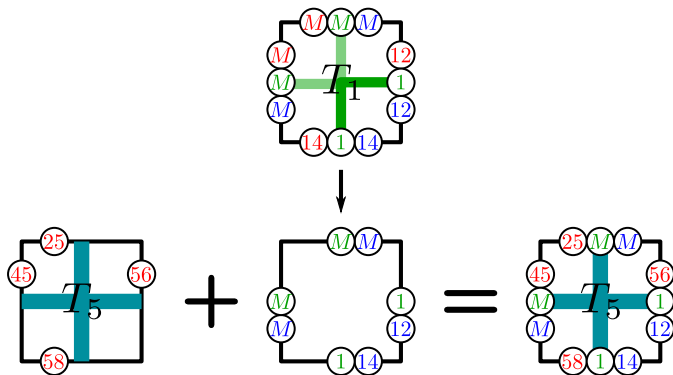
Pairs on X - and Y -branches could be allowed on X - and Y -edges of different decorated T_j . The branches have to be *synchronized*.

Step 5: synchronizing network branches



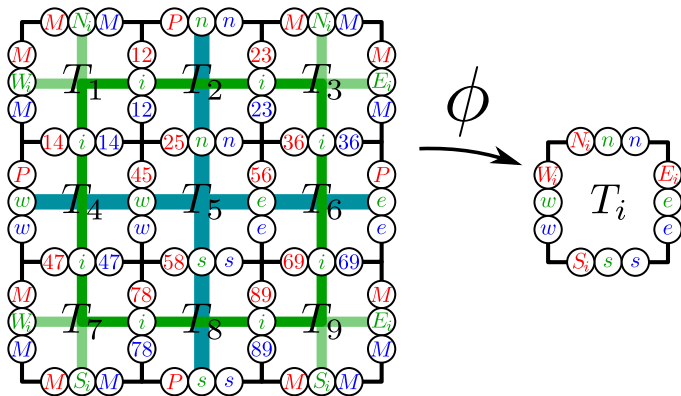
This is done by allowing on T_5 the pairs of every non-central tile.

Step 5: synchronizing network branches



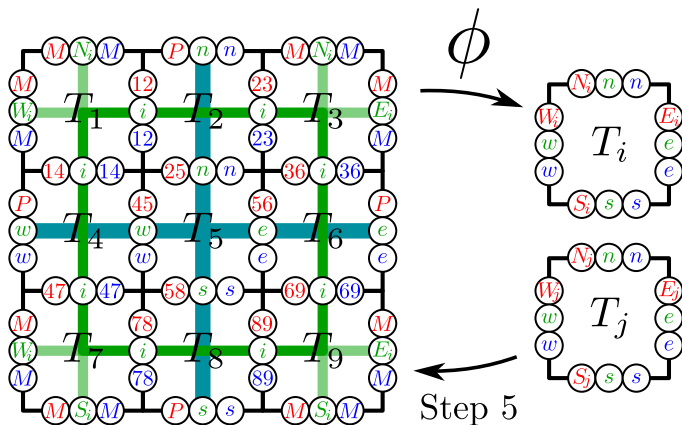
This is done by allowing on T_5 the pairs of every non-central tile. This double the number of tiles: there are thus 1736 τ -tiles in all.

Back to the map ϕ



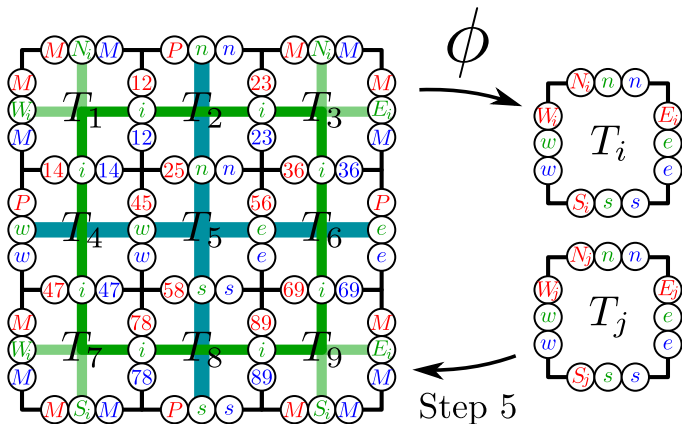
ϕ maps any τ -macro-tile with i on the ring onto a decorated T_i .
But why should it be a τ -tile?

Back to the map ϕ



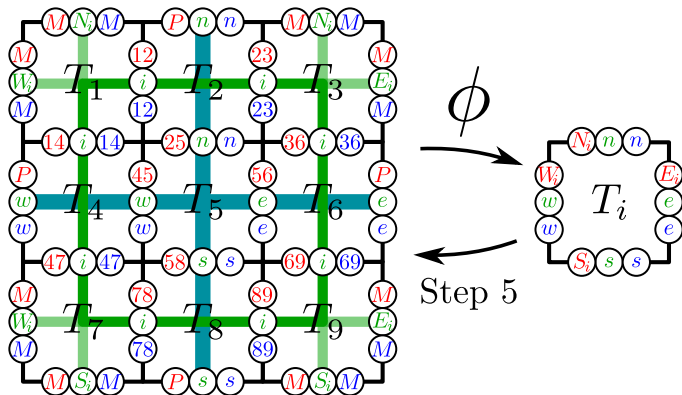
The decorated T_j the central tile is derived from (Step 5) is in τ . We claim that it is one and the same tile, except if $i = 5$.

Back to the map ϕ



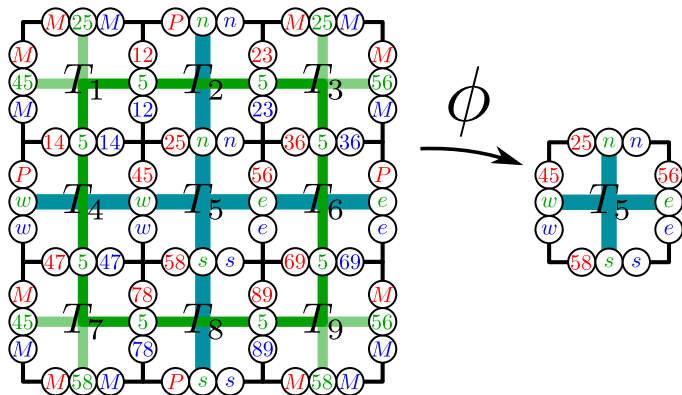
When the ring intersects an X -branch of the network, it forces the green/blue pair to be allowed on some decorated T_i (Step 4).

Back to the map ϕ



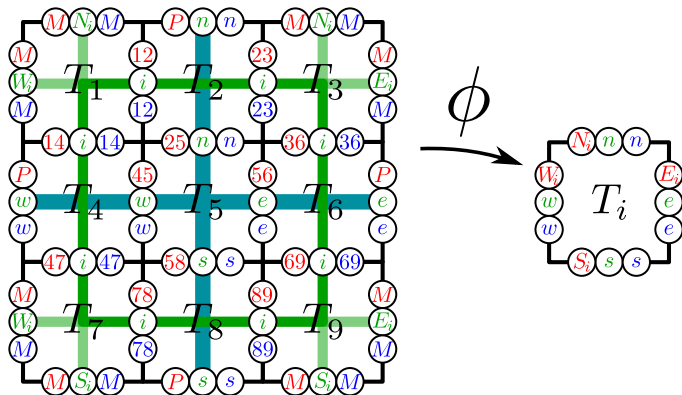
The green/blue indices of $T_{i \neq 5}$ determine i (Duck lemma) $\rightsquigarrow i = j$.

Back to the map ϕ



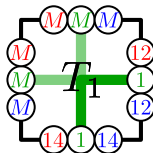
For $i = 5$, ϕ simply maps the macro-tile onto its central tile!

Back to the map ϕ



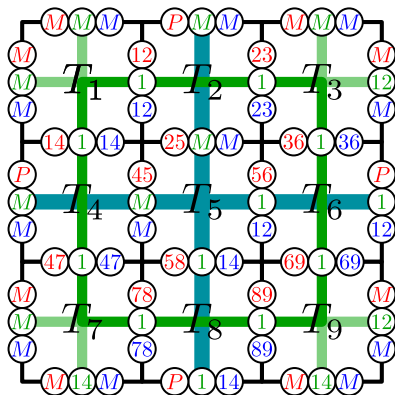
Whatever i on the ring, ϕ thus maps the τ -macro-tile onto a τ -tile. It is moreover a bijection: the inverse function is straightforward.

Back to the map ϕ



In particular, applying *ad infinitum* ϕ^{-1} to any τ -tile yields arbitrarily large τ -patches, hence a τ -tiling by compacity.

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