

# How to conjecture and prove that the generating function of the Yang-Zagier numbers is algebraic<sup>1</sup>

CAP21 (IHES, France)

Sergey Yurkevich

Inria Saclay and University of Vienna

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<sup>1</sup>Joint work with [Alin Bostan](#) and [Jacques-Arthur Weil](#).

# Two sequences

$$(a_n)_{n \geq 0} = (1, -48300, 7981725900, -1469166887370000, \dots)$$

$$(b_n)_{n \geq 0} = (1, -144900, 88464128725, -62270073456990000, \dots)$$

# Origin of $a_n$ and $b_n$

- In [Arithmetic and Topology of Differential Equations, 2018](#) by [Don Zagier](#):

$$c_{n-3} + 20(4500n^2 - 18900n + 19739)c_{n-2} + 80352000n(5n-1)(5n-2)(5n-4)c_n + 25(2592000n^4 - 16588800n^3 + 39118320n^2 - 39189168n + 14092603)c_{n-1} = 0,$$

with initial terms  $c_0 = 1$ ,  $c_1 = -161/(2^{10} \cdot 3^5)$  and  $c_2 = 26605753/(2^{23} \cdot 3^{12} \cdot 5^2)$ .

- Recursion comes from physics: integral over a moduli space (“topological ODE”) [[Bertola, et al, 2015](#)].

## Problem (Zagier, 2018)

Find  $(u, v) \in \mathbb{Q}^* \times \mathbb{Q}^*$  such that  $c_n \cdot (u)_n \cdot (v)_n \cdot w^n \in \mathbb{Z}$  for some  $w \in \mathbb{Z}^*$ .

$$(u)_n := u \cdot (u+1) \cdots (u+n-1).$$

- [[Yang and Zagier](#)]:  $a_n = c_n \cdot (3/5)_n \cdot (4/5)_n \cdot (2^{10} \cdot 3^5 \cdot 5^4)^n \in \mathbb{Z}$ ,
- [[Dubrovin and Yang](#)]:  $b_n = c_n \cdot (2/5)_n \cdot (9/10)_n \cdot (2^{12} \cdot 3^5 \cdot 5^4)^n \in \mathbb{Z}$ .

# Mystery of $a_n$ and $b_n$

- “Yang and I found a formula showing that the numbers  $a_n$  are integers of exponential growth and hence can be expected to have a generating series that is a **period**, although we have not succeeded in finding it” – [Zagier, 2018]
- “Dubrovin and Yang found that the numbers  $b_n$  are *also* integral and that in this case the generating function is not only of Picard-Fuchs type, but is actually **algebraic!**” – [Zagier, 2018]
- “So this is a very mysterious example [...] of numbers defined by recursions with polynomial coefficients.” – [Zagier, 2018]
- “My presumed arithmetic intuition [...] was entirely broken” – [Wadim Zudilin]

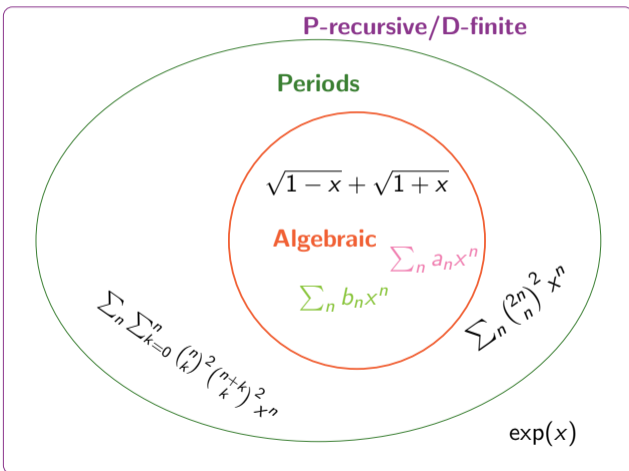
## Problem

Investigate the nature of  $(a_n)_{n \geq 0}$ ,  $(b_n)_{n \geq 0}$  and similar sequences.

## Theorem (Bostan, Weil, Y.)

*The generating functions of both  $(a_n)_{n \geq 0}$  and  $(b_n)_{n \geq 0}$  are algebraic.*

# Definitions and interactions

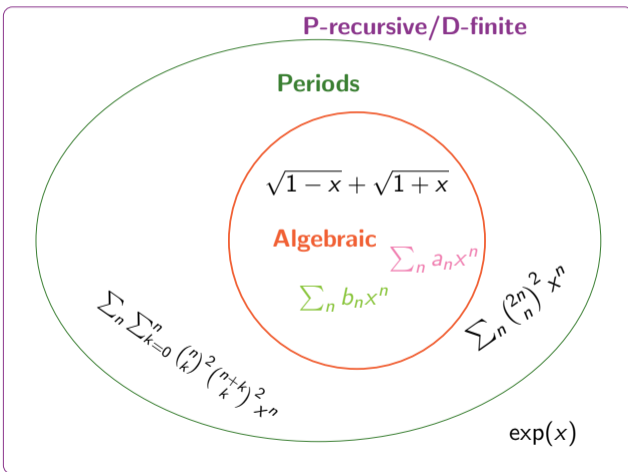


A sequence  $(u_n)_{n \geq 0}$  is **P-recursive**, if it satisfies a linear recurrence with polynomial coefficients:

$$c_r(n)u_{n+r} + \cdots + c_0(n)u_n = 0.$$

$u_n = 1/n!$  satisfies  $nu_n = u_{n-1}$ .

# Definitions and interactions



A power series  $f(x) \in \mathbb{Q}[[x]]$  is **D-finite** if it satisfies a linear differential equation with polynomial coefficients:

$$p_n(x)f^{(n)}(x) + \cdots + p_0(x)f(x) = 0.$$

This equation can be rewritten:  $L \cdot f = 0$ ,

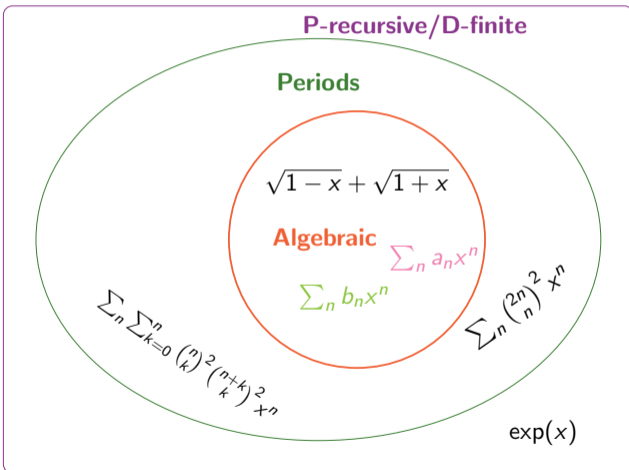
$$L = p_n(x)\partial^n + \cdots + p_0(x) \in \mathbb{Q}(x)[\partial],$$

where  $\partial := \frac{d}{dx}$ .

$\exp(x)$  satisfies  $\exp'(x) = \exp(x)$ .

$$L = \partial - 1.$$

# Definitions and interactions



A power series  $f(x) \in \mathbb{Q}[[x]]$  is called a **Period function** if it is an integral of a rational function in  $x$  and  $t_1, \dots, t_n$  over a semi-algebraic set.

$$p(e) = 4 \int_0^1 \sqrt{\frac{1-e^2 t^2}{1-t^2}} dt$$

$$= 4 \iint \frac{du dv}{1 - \frac{1-e^2 u^2}{(1-u^2)v^2}} \text{ and}$$

$$((e - e^3)\partial^2 + (1 - e^2)\partial + e) \cdot p = 0,$$

$$p(e) = 2\pi - \frac{\pi}{2}e^2 - \frac{3\pi}{32}e^4 - \dots$$

## Back to $a_n$ and $b_n$

- $(a_n)_n$  and  $(b_n)_n$  are **P-recursive** sequences  $\Rightarrow$  generating functions are **D-finite**.

$$L_a = 1800x(7x - 62)(x^2 + 50x + 20)\partial^2 + 720(42x^3 + 173x^2 - 14230x - 620)\partial + 6048x^2 - 139453x - 249550 \in \mathbb{Q}(x)[\partial],$$

$$L_b = 90000x^3(2911x + 310)(x^2 + 50x + 20)\partial^4 + 18000x^2(154283x^3 + 5185005x^2 + 1675710x + 142600)\partial^3 + 50x(147290778x^3 + 2740219655x^2 + 566777510x + 37497600)\partial^2 + 5(919899288x^3 + 5629046605x^2 + 1348939210x + 10713600)\partial + 18(13937868x^2 - 1076845x + 1247750) \in \mathbb{Q}(x)[\partial].$$

- The generating functions of  $(a_n)_{n \geq 0}$  and  $(b_n)_{n \geq 0}$  solve  $L_a \cdot y = 0$  and  $L_b \cdot y = 0$ .



# Main problem

## Stanley's problem (1980)

Given a **D-finite** series how to prove or disprove that it is **algebraic**?

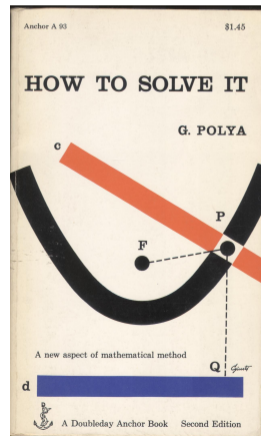
## Useful (sub-)question

Given a **D-finite** series how to **conjecture** whether it is **algebraic**?

- Guess & Prove approach – **but** algebraicity degree can be arbitrarily high.
- Algorithm for **rational solutions** of linear ODE [Liouville, 1833], [Barkatou, 1998].
- Solved **in theory** [Singer, 1979, 2014] – **but** usually not applicable in practice.
- New practical algorithm for **disproving algebraicity** [Bostan, Rivoal, Salvy, 2021].
- Several tests for justifying **algebraicity** based on **conjectures** or **numerics**:  
work well in practice but do not provide proofs.
- Applied differential Galois theory sometimes efficient proving **algebraicity**.

# The Guess-and-Prove approach

- Experimental mathematics and “Guess-and-Prove” propagated by G. Pólya.
- Extremely fruitful when using a computer.
- **P-recursive** sequences/**D-finite** functions: ideal data structure for guessing.
- Find new structure and simpler formulas.
- Very efficient and easy in practice (e.g. with Maple).
- For guessing: gfun [Salvy, Zimmermann 1992].
- For proving: Theory of  $\mathbb{Q}(x)[\partial]$  and effective properties.



# Guess & Prove for P-recursive sequences/D-finite functions

- Given  $u_0, u_1, \dots, u_N \in \mathbb{Q}$ , finding some  $c_0(n), \dots, c_r(n)$  for fixed  $r$  and of fixed maximal degree  $d$  such that

$$c_r(n)u_{n+r} + \dots + c_0(n)u_n = 0, \quad \text{for } n = 0, \dots, N - r$$

amounts to solving a *linear system*.

- Showing equality of two **P-recursive** sequences/**D-finite** functions is decidable and easy/efficient in practice.
- $\rightarrow$  Maple (for  $(a_n)_{n \geq 0}$  and  $(b_n)_{n \geq 0}$ )

## Guess & Prove for algebraicity: Toy example

- Let  $f(x) = 2 - x^2/4 - 5x^4/64 + \dots$  be a solution of

$$(4x^2 - 4) \cdot f''(x) + 4x \cdot f'(x) - f(x) = 0.$$

- How to prove that  $f(x)$  is algebraic?
- Idea: *guess* a minimal polynomial  $P(x, t)$  and then *prove* its correctness.
- Let  $g(x) = 2 - x^2/4 - 5x^4/64$ . Finding  $(c_1, c_2, \dots, c_9) \in \mathbb{Q}$  such that

$$P(x, f(x)) = (c_1 + c_2x + c_3x^2) + (c_4 + c_5x + c_6x^2)g^2 + (c_7 + c_8x + c_9x^2)g^4 = 0$$

results in a *linear system* which we can easily solve:  $c_3 = 4, c_4 = -4, c_7 = 1$ .

- Guess:  $P(x, t) = 4x^2 - 4t^2 + t^4$ .
- Effective version of Abel's Theorem: Any solution  $h(x)$  of  $P(x, h(x)) = 0$  satisfies:

$$(4x^2 - 4) \cdot h''(x) + 4x \cdot h'(x) - h(x) = 0.$$

- Proof: Conclude with uniqueness.
- In general algebraicity degree can be arbitrarily high:  $N(1 + x)f'(x) = f(x)$ .

# Grothendieck-Katz conjecture: “testing” algebraicity

- $L \cdot y = 0$  is equivalent to  $Y' = A(x)Y$ , where  $A(x) \in M^{n \times n}(k)$  and  $k = \mathbb{Q}(x)$ .
- The  $p$ -curvature is the matrix  $A_p(x) \in \mathbb{Q}(x)$ , where

$$A_0(x) = \text{Id}_n, \quad \text{and} \quad A_{\ell+1}(x) = A'_\ell(x) + A_\ell(x)A(x) \quad \text{for} \quad \ell \geq 0.$$

- It holds that  $\partial^k Y = A_k Y$  for  $k = 0, 1, \dots$

## Conjecture (Grothendieck 1960's; Katz, 1972)

*All solutions of  $Y' = A(x)Y$  are algebraic if and only if  $A_p = 0 \pmod p$  for almost all primes  $p$ .*

- If  $A(x)$  is given and we find that  $A_{p_i} = 0 \pmod{p_i}$  for all primes  $p_1, p_2, \dots, p_N$ , we can conjecture that all solutions of  $Y' = A(x)Y$  are algebraic.
- $A_p \pmod p$  can be efficiently computed [[Bostan, Caruso, Schost, 2015](#)].

# “Testing” algebraicity for $L_a$ and $L_b$

- It holds that  $\partial^k Y = A_k Y$  for  $k = 0, 1, \dots$
- The right Euclidean division of  $\partial^k$  by  $L$  in  $\mathbb{Q}(x)[\partial]$  reads:

$$\partial^k = (\dots) \cdot L + d_{n-1}(x)\partial^{n-1} + \dots + d_0(x),$$

for  $d_0(x), \dots, d_{n-1}(x) \in \mathbb{Q}(x)$  depending on  $k$ .

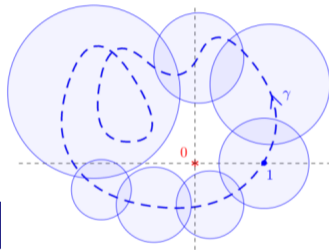
- Hence, the first row of  $A_k$  is  $d_0(x), \dots, d_{n-1}(x)$ .
- $\rightarrow$  Maple (for  $L_a$  and  $L_b$ )

# Monodromy group: quantifying **algebraicity**

- $L \cdot y = 0$  for  $L \in \mathbb{Q}(x)[\partial]$  has  $n = \text{ord}(L)$  linearly independent solutions.
- Assume  $f_1, \dots, f_n$  are linearly independent solutions at 0. If we analytically continue them along a closed loop in  $\mathbb{C}$ , we find  $\tilde{f}_1, \dots, \tilde{f}_n$  possibly different.
- There exists  $M_{\underline{f}} \in \text{GL}(n, \mathbb{C})$  such that

$$\begin{pmatrix} \tilde{f}_1 \\ \vdots \\ \tilde{f}_n \end{pmatrix} = M_{\underline{f}} \begin{pmatrix} f_1 \\ \vdots \\ f_n \end{pmatrix}.$$

- The matrices  $M_{\underline{f}}$  define the so-called monodromy group  $M$ .



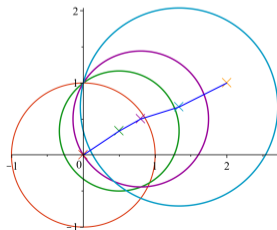
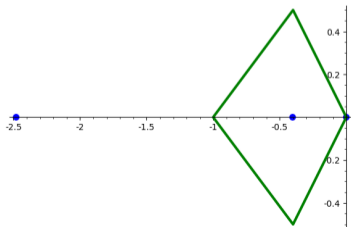
## Theorem (Singer, Ulmer, 1993)

*Let  $f$  be a solution of  $L \cdot y = 0$ . The algebraicity degree of  $f$  is equal to the cardinality of the orbit of  $f$  under the action of  $M$ .*

- Analytic continuation of **D-finite** functions can be efficiently computed numerically [Chudnovsky<sup>2</sup>, 1987], [van der Hoeven, 1999, 2001], [Mezzarobba, 2010].

# Quantifying **algebraicity** for $L_a$ and $L_b$

- Very efficient analytic continuation implemented by Mezzarobba in SageMath.
- → SageMath.
- Numerical computations suggest: solutions of  $L_a$  and  $L_b$  have alg. degree 120.





# Differential Galois theory: proving **algebraicity**

- $L \cdot y = 0$  is equivalent to  $Y' = A(x)Y$ , where  $A(x) \in M^{n \times n}(k)$  and  $k = \overline{\mathbb{Q}}(x)$ .
- Picard-Vessiot extension:  $K = k(U)$ , where  $U$  is a fundamental solution matrix.
- The differential Galois group  $G$  is the group of field automorphisms of  $K$  which commute with the derivation and leave all elements of  $k$  invariant:

$$G := \text{Aut}_{\partial}(K/k) = \{\sigma \in \text{Aut}(K) : \sigma|_k \equiv \text{id}_k \text{ and } \sigma \circ \partial \equiv \partial \circ \sigma\}.$$

- $G$  is a linear algebraic subgroup of  $\text{GL}_n(\overline{\mathbb{Q}})$ .
- $G$  stabilizes the ideal of differential relations between solutions. Moreover:

Theorem (Kolchin, 1948)

$L \cdot y = 0$  has a basis of **algebraic** solutions if and only if  $G$  is finite.

- In practice  $G$  is difficult to compute [Hrushovski, 2002], [Feng, 2015], [van der Hoeven, 2007], [Amzallag, Minchenko, Pogudin, 2018], [Sun, 2019].

# Differential Galois theory: proving **algebraicity**

- $L \cdot y = 0$  is equivalent to  $Y' = A(x)Y$ , where  $A(x) \in M^{n \times n}(k)$  and  $k = \overline{\mathbb{Q}}(x)$ .
- $G$  is the differential Galois group.
- Galois-Lie algebra  $\mathfrak{g} := \text{Lie}(G)$ : Lie algebra of  $G$ , i.e. the tangent space of  $G$  at  $\text{id}$ .
- $\mathfrak{g}$  measures the transcendence of  $K$  over  $k$ :

## Theorem (Kolchin, 1948)

*If  $K$  is the Picard-Vessiot extension of  $Y' = A(x)Y$  and  $\mathfrak{g} = \text{Lie}(G)$ , then*

$$\dim_{\mathbb{C}}(G) = \dim_{\mathbb{C}}(\mathfrak{g}) = \text{trdeg}(K/k).$$

- Theory and algorithm for computing  $\mathfrak{g}$  [Barkatou, Cluzeau, Di Vizio, Weil, 2020].
- Idea: Compute symmetric powers of  $L$  and find **rational solutions** of them. These solutions yield information for  $\mathfrak{g}$  via solving a **linear** system.

## Toy example

- The operator  $L = (4x^2 - 4)\partial^2 + 4x\partial - 1$  has a basis of algebraic solutions:

$$\sqrt{1+x} + \sqrt{1-x} \text{ and } \sqrt{1+x} - \sqrt{1-x}.$$

- $L \cdot y = 0$  is equivalent to  $Y' = A(x)Y$  where  $A(x) = \begin{pmatrix} 0 & 1 \\ \frac{1}{4x^2-4} & \frac{-4x}{4x^2-4} \end{pmatrix}$ .


# Toy example

- The operator  $L = (4x^2 - 4)\partial^2 + 4x\partial - 1$
- $L \cdot y = 0$  is equivalent to  $Y' = A(x)Y$  where  $A(x) = \begin{pmatrix} 0 & 1 \\ \frac{1}{4x^2-4} & \frac{-4x}{4x^2-4} \end{pmatrix}$ .
- If  $Y = (y_1, y_2)^t$  is a solution to  $Y' = A(x)Y$  then  $Y = (y_1^2, 2y_1y_2, y_2^2)^t$  is a solution to the symmetric square system  $Y' = A^{(2)}(x)Y$ , where now

$$A^{(2)}(x) = \frac{1}{4(x^2 - 1)} \begin{pmatrix} 0 & 4x^2 - 4 & 0 \\ 2 & -4x & 8x^2 - 8 \\ 0 & 1 & -8x \end{pmatrix}.$$

- It has rational solutions!  $F_1 = (4x, 4, x/(x^2 - 1))^t$ ,  $F_2 = (-4, 0, 1/(x^2 - 1))^t$ .
- If  $M \in \mathfrak{g}^{(2)}$  then  $MF = 0$  and  $M$  comes from a symmetric square. I.e.  $M$  satisfies

$$\begin{pmatrix} 2m_{1,1} & m_{1,2} & 0 \\ 2m_{2,1} & m_{1,1} + m_{2,2} & 2m_{1,2} \\ 0 & m_{2,1} & 2m_{2,2} \end{pmatrix} \cdot F_\ell = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}, \quad m_{i,j} \in \mathbb{Q}(x), \ell = 1, 2.$$

- The only solution is  $m_{i,j} = 0$ . Hence  $\mathfrak{g}^{(2)} = \mathfrak{g} = 0$ . All solutions of  $L$  are algebraic. 

# The generating sequence of $(b_n)_n$ is algebraic (known to Dubrovin & Yang)

- For  $L_b$  same method as in the toy example works.
- $L_b \cdot y = 0$  equivalent to  $Y' = A(x)Y$  for  $A(x) \in M^{4 \times 4}(\mathbb{Q}(x))$ .
- The fifth symmetric power  $Y' = A^{(5)}(x)Y$  has rational solutions.
- $A^{(5)}(x) \in M^{N \times N}(\mathbb{Q}(x))$ , where  $N = \binom{4+5-1}{4-1} = 56$ .
- Finding the rational solutions takes  $\approx 2$  min on a regular PC.
- The corresponding system in  $m_{i,j}$  has no non-zero solutions in  $\mathbb{Q}(x)$  ( $\approx 15$  sec).
- $\Rightarrow g_b = 0$ , therefore  $L_b$  has only algebraic solutions.

# The generating sequence of $(a_n)_n$ is algebraic (new)

- For the generating function of  $(a_n)_{n \geq 0}$  same method as for  $(b_n)_{n \geq 0}$  works.
- The 20th symmetric power has rational solutions ( $\approx 1$  sec).
- $A^{(20)} \in M^{N \times N}(\mathbb{Q}(x))$ , where  $N = \binom{2+20-1}{2-1} = 21$ .
- The corresponding system in  $m_{i,j}$  has no non-zero solutions in  $\mathbb{Q}(x)$  ( $\approx 4$  sec).
- $\Rightarrow g_a = 0$ , therefore  $L_a$  has only algebraic solutions.

■  $\rightarrow$  Maple

# Experimental mathematics: more similar examples

## Problem

Find  $(u, v) \in \mathbb{Q}^* \times \mathbb{Q}^*$  such that  $c_n \cdot (u)_n \cdot (v)_n \cdot w^n \in \mathbb{Z}$  for some  $w \in \mathbb{Z}^*$ .

$$(u)_n := u \cdot (u + 1) \cdots (u + n - 1).$$

#	$u$	$v$	ODE order	degree	#	$u$	$v$	ODE order	degree
$a_n$	3/5	4/5	2	120	$f_n$	19/60	49/60	4	155520
$b_n$	2/5	9/10	4	120	$g_n$	19/60	59/60	4	46080
$c_n$	1/5	4/5	2	120	$h_n$	29/60	49/60	4	46080
$d_n$	7/30	9/10	4	155520	$i_n$	29/60	59/60	4	155520
$e_n$	9/10	17/30	4	155520					

- “Test”: 0  $p$ -curvatures for primes  $< 100 \rightarrow$  expect **algebraic** generating functions.
- Quantify: Guesses for degrees based on numerics.
- Proof: Done:  $a_n, b_n, c_n$ . In progress:  $d_n, e_n, f_n, g_n, h_n, i_n$ .

# Summary

- Both sequences  $(a_n)_{n \geq 0}$  and  $(b_n)_{n \geq 0}$  have algebraic generating functions, hence they are particular **periods**.
- Guess & Prove approach often provides useful insight but is sometimes infeasible.
- The Grothendieck-Katz conjecture allows efficient “testing” whether a **D-finite** series is **algebraic**.
- Numerical monodromy group calculations allow efficient quantifying **algebraicity** of **D-finite** series.
- Differential Galois theory allows efficient proving that **D-finite** series is **algebraic**.



## Bonus: explicit solution for $\sum_{n \geq 0} a_n x^n$

We saw that  $\sum_{n \geq 0} a_n x^n$  is a solution of

$$q_2(x)y''(x) + q_1(x)y'(x) + q_0(x)y(x) = 0, \quad \text{where} \quad (1)$$

$$q_2(x) = 5x(302400x - 31)(373248000x^2 + 216000x + 1),$$

$$q_1(x) = 1354442342400000x^3 + 64571904000x^2 - 61473600x - 31,$$

$$q_0(x) = 300(902961561600x^2 - 240974784x - 4991).$$

Maple's `dsolve(deq)` shows that every solution of (1) is a linear combination of

$$u_1(x) \cdot {}_2F_1 \left[ \begin{matrix} -1/60 & 11/60 \\ 2/3 \end{matrix}; \frac{p_1(x)}{p_2(x)} \right] \quad \text{and} \quad u_2(x) \cdot {}_2F_1 \left[ \begin{matrix} 19/60 & 31/60 \\ 4/3 \end{matrix}; \frac{p_1(x)}{p_2(x)} \right],$$

where  ${}_2F_1 \left[ \begin{matrix} a & b \\ c \end{matrix}; x \right]$  is the Gaussian hypergeometric function

$${}_2F_1 \left[ \begin{matrix} a & b \\ c \end{matrix}; x \right] := \sum_{n=0}^{\infty} \frac{(a)_n (b)_n}{(c)_n} \frac{x^n}{n!}, \quad (u)_j := u(u+1) \cdots (u+j-1).$$

## Bonus: Origin of $(c_n)_{n \geq 0}$

- For a simple Lie-algebra  $(\mathfrak{g}, [\cdot, \cdot])$  [Bertola, Dubrovin, Yang, 2015] define the co-called *topological ordinary differential equation*

$$\frac{d}{d\lambda} M = [M, \Lambda],$$

where  $M = M(\lambda)$  and  $\Lambda = I_+ + \lambda E_{-\theta}$ , for a principal nilpotent element  $I_+ = \sum_{i=1}^n E_i$  and (normalized)  $E_{-\theta} \in \mathfrak{g}_{-\theta}$ .

- For  $\mathfrak{g} = \mathfrak{sl}_{n+1}(\mathbb{C})$  one finds

$$\Lambda = \begin{pmatrix} 0 & I_n \\ \lambda & 0 \end{pmatrix}, \quad I_n \text{ is the } n \times n \text{ identity matrix.}$$

and the (normalized) (dominant) ODE reads

$$64800000x^3(x+155)y^{(iv)}(x) + (x^2 - 1220x + 623875)y(x) + 7200(10x^2 + 3209x + 133920)y'(x) + 18000x(5x^2 + 6091x + 1874880)y''(x) + 12960000x^2(18x + 3565)y'''(x) = 0$$

- Then  $\sum_{n \geq 0} c_n x^n$  is the unique power series solution.