

# Biological Regulatory Networks (BRN) modeled as an institution: property preservation along embedding

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Biological regulatory networks (BRN) :

genes or derived products as proteins interact with others such that cell behaviours are regulated

Goal :

Understanding the functioning of biological regulatory networks in order to predict some knowledge about behaviours

Motivation:

- Focusing on an isolated part of the global system, perceived as having a particular biological function
- and studying in which cases properties of biological regulatory networks are preserved.

# Qualitative modeling frameworks for BRN

multivalued discrete approach developed by R. Thomas:

- constituent concentrations are abstracted by integers to denote **thresholds** from which they can act on other constituents
- biological systems are described by an **interaction graph** defining the static part
- a huge but finite set of **state transition graphs** defining all the possible dynamics, or models.
- discretization preserves qualitative biological observations/experiments expressed as **temporal properties**

# Decomposing BRN as sub-BRNs expressing biological functions

- As expected, model-checking technics are efficient to study small BRNs, but cannot cope with large BRNs.
- In practice, biologists study small BRNs, of particular importance w.r.t a **biological function**.
- Interactions between BRNs are studied afterwards, whether properties of sub-BRNs are preserved or not.

Following an institution-like approach, we provide a logic to characterize dynamic of BRNs

- **Syntax**

**Signatures** are interaction graphs

**Sentences** are CTL-X formulas

(Computational Tree Logic without the neXt temporal operator)

- **Semantics**

**Models** are a particular sub-class of state transition systems

**Satisfaction relation** is the usual one for CTL

and we establish **the satisfaction condition** with some restricted conditions on the embedding of a BRN within an other one.

## Definition

An **institution**  $\mathcal{I} = (Sig, Sen, Mod, \models)$  consists of

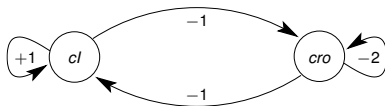
- a category  $Sig$  of **signatures**,
- a functor  $Sen : Sig \rightarrow Set$  giving for each signature  $\Sigma$  a set, element of **sentences**,
- a contravariant functor  $Mod : Sig^{op} \rightarrow Cat$  giving for each signature  $\Sigma$  a category of  **$\Sigma$ -models**
- a  $|Sig|$ -indexed family of **satisfaction relations**  
 $\models_{\Sigma} \subseteq |Mod(\Sigma)| \times Sen(\Sigma)$

such that the **satisfaction condition** holds:

$\forall \sigma : \Sigma \rightarrow \Sigma', \forall \mathcal{M}' \in |Mod(\Sigma')|, \forall \varphi \in Sen(\Sigma),$

$$\mathcal{M}' \models_{\Sigma'} Sen(\sigma)(\varphi) \Leftrightarrow Mod(\sigma)(\mathcal{M}') \models_{\Sigma} \varphi$$

## Biological regulatory graph : Example



- Vertex = genes
- Edges = interactions (activation or inhibition)
- Labels =
  - $\left\{ \begin{array}{l} S_n : \text{Sign } ((+) \text{ for activation, } (-) \text{ for inhibition}) \\ Th : \text{Threshold (interaction level)} \end{array} \right.$

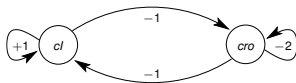
A **BRN-signature** is a labeled directed graph  $G = \langle V, F, Sn, Th \rangle$  where :

- 1  $V$  is a finite set of *variables*.
- 2  $F \subseteq V \times V$  denotes the set of edges.  
For any  $i \in V$ ,  $G_i^+$ , resp.  $G_i^-$ , denotes the set of successors, resp. predecessors, of  $i$  in  $\langle V, F \rangle$ .
- 3  $Sn$  is a mapping from  $F$  to  $\{+, -\}$ .
- 4  $Th$  is a mapping from  $F$  to  $\mathbb{N}^*$  such that:

$$\forall i \in V, \forall j \in G_i^+, Th(i, j) = c \wedge c \neq 1 \Rightarrow \exists k \in G_i^+ : Th(i, k) = c-1$$



Dynamic behaviors = traces associating at each time to each gene its concentration level.



We know that  $x_{cl} \in \{0, 1\}$  and  $x_{cro} \in \{0, 1, 2\}$

The **state space**  $S_G$  of  $G = \langle V, F, Sn, Th \rangle$  is the set of mappings  $s : V \rightarrow \mathbb{N}$  s.t.  $\forall i \in V, s(i) \in \{0, \dots, b_i\}$ .

with  $b_i = |\{s \in \mathbb{N}^* \mid \exists j \in G_i^+, Th(i, j) = s\}|$

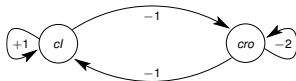
The concentration level of  $i \in V$  evolves over time depending on concentration levels of its resources, i.e.  $i$ 's predecessors having reached a concentration level to affect  $i$ 's one:

$$R_{G,i}(s) = \begin{cases} \{j \in G_i^- | (Sn(j, i) = + \text{ and } s(j) \geq Th(j, i))\} \\ \cup \\ \{j \in G_i^- | (Sn(j, i) = - \text{ and } s(j) < Th(j, i))\} \end{cases}$$

## Resources

$cl$	$cro$	$R_{G,cl}$	$R_{G,cro}$
0	0	$\{cro\}$	$\{cl, cro\}$
0	1	$\emptyset$	$\{cl, cro\}$
0	2	$\emptyset$	$\{cl\}$
1	0	$\{cl, cro\}$	$\{cro\}$
1	1	$\{cl\}$	$\{cro\}$
1	2	$\{cl\}$	$\emptyset$

## Graph



Hence, a resource is the presence of an activator or the absence of an inhibitor.

- No indication in the signature to decide the concentration level that  $i$  can reach.
- This degree of freedom gives rise to a class of possible **G-models**, so-called dynamics of  $G$ .

Let  $\kappa = \{(i, w) \mid i \in V \wedge w \subseteq G_i^-\}$  be the set of all subsets of predecessors in  $G$  for  $i \in V$ .

A **G-model** is a mapping  $p : \kappa \rightarrow \mathbb{N}$  s. t.:

$\forall i \in V, p((i, \emptyset)) = 0 \wedge (\forall (i, w \neq \emptyset) \in \kappa, p((i, w)) \in \{0, \dots, b_i\})$ .

# Models: asynchronous transition system (AST)

Since in the nature, several variables cannot cross a threshold simultaneously, we make evolve one variable  $i$  by one unit in the direction its concentration level specified by  $p$ .

The **asynchronous transition system** generated from  $p$  is a directed graph  $GTA((G, p)) = \langle S_G, T \rangle$  s.t.:

- $\forall s \in S_G, (s, s) \in T \Leftrightarrow \forall i \in V, s(i) = p((i, R_{G,i}(s)))$
- $\forall s \neq s' \in S_G, (s, s') \in T$  iff:
  - there exists  $i \in V$ , s.t.

$$s'(i) = \begin{cases} s(i) + 1 & \text{and } s(i) < p((i, R_{G,i}(s))) \\ s(i) - 1 & \text{and } s(i) > p((i, R_{G,i}(s))) \end{cases}$$

- and  $s'(j) = s(j)$  for every  $j \in V \setminus \{i\}$ .

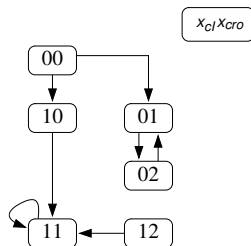
## Graph



## Resources and a model

$cl$	$cro$	$R_{G,cl}$	$p_{cl}$	$R_{G,cro}$	$p_{cro}$
0	0	$\{cro\}$	1	$\{cl, cro\}$	2
0	1	$\emptyset$	0	$\{cl, cro\}$	2
0	2	$\emptyset$	0	$\{cl\}$	1
1	0	$\{cl, cro\}$	1	$\{cro\}$	1
1	1	$\{cl\}$	1	$\{cro\}$	1
1	2	$\{cl\}$	1	$\emptyset$	0

## Asynchronous model



Sentences are CTL-X formulas whose atomic formulas are comparisons between a concentration level of a variable with some threshold values.

- Atomic formulas are of the form  $(i \sim s)$  where  $i \in V$ ,  $s \in \{0, \dots, b_i\}$  and  $\sim \in \{=, <, >\}$ .
- Formulas are of the form:

*ATOM* | *For*  $\Rightarrow$  *For* | *For*  $\wedge$  *For* | *For*  $\vee$  *For* |  $\neg$ *For*  
*AG For* | *EG For* | *AF For* | *EF For* | *A[For U For]* | *E[For U For]*

A (for All path), E (there exists one path), F (there exists one state in the path), G (for all states in the path), U (until).

$i \geq s$  (resp.  $i \leq s$ ) will denote  $i = s \vee i > s$  (resp.  $i = s \vee i < s$ ).

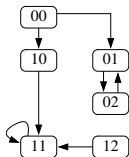
Satisfaction relation between models and sentences for BRN is derived from the usual one between transition systems and CTL-X formulas.

For a model  $p$  with  $GTA((G, p))$  as associated AST:

$$p \models_G \varphi \Leftrightarrow GTA((G, p)) \models \varphi$$

with for  $s \in S_G$ ,

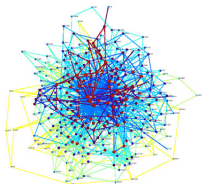
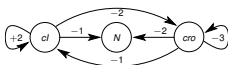
$$L(s) = \{ i > l, i < l', i = l'' \mid \left. \begin{array}{l} i \in V, l, l', l'' \in \{0, 1, \dots, b_i\}, \\ s(i) > l, s(i) < l', s(i) = l'' \end{array} \right\}$$



$$\models x_{cl} = 1 \wedge x_{cro} = 0 \Rightarrow EG(x_{cl} = 1)$$

# Embedding : motivation and objectives

Biologists study small BRNs viewed as a biologic function.

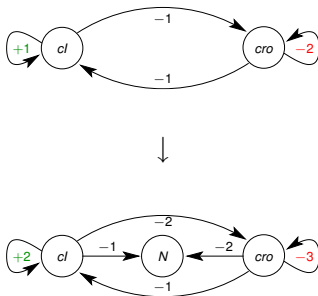




- Embedding = inclusion of a graph  $G$  in a graph  $G'$
- Embedding preserves genes, interactions, and order between thresholds in relation to each gene.

**Effect** : Shifting the thresholds

**Example**



$G = \langle V, F, Sn, Th \rangle$  and  $G' = \langle V', F', Sn', Th' \rangle$  signatures.

An embedding  $G \rightarrow G'$  is an injective mapping  $\sigma : V \rightarrow V'$  s.t.:

- 1  $\forall i, j \in V, (i, j) \in F \Leftrightarrow (\sigma(i), \sigma(j)) \in F'$
- 2  $\forall i, j \in V, (i, j) \in F, Sn(i, j) = Sn'(\sigma(i), \sigma(j))$
- 3  $\forall i \in V, \forall j, k \in G_i^+,$   
 $Th(i, j) = Th(i, k) \Leftrightarrow Th'(\sigma(i), \sigma(j)) = Th'(\sigma(i), \sigma(k))$
- 4  $\forall i \in V, \forall j, k \in G_i^+,$   
 $Th(i, j) < Th(i, k) \Leftrightarrow Th'(\sigma(i), \sigma(j)) < Th'(\sigma(i), \sigma(k))$
- 5  $\forall j \in V, \forall k' \in V',$   
 $(k', \sigma(j)) \in F' \Rightarrow \exists i \in V, (i, j) \in F \wedge \sigma(i) = k'$

# Translation of formulas along signature embedding

$$\sigma : g \rightarrow G$$

**Idea** : Translating a threshold into an interval of values.

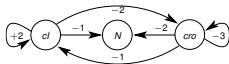
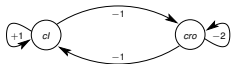
Notation:  $\sigma_{g_1}(0) = 0$  and

for  $l \neq 0$ ,  $\sigma_{g_1}(l) = Th_G(g_1, g_2)$  for  $g_2$  s.t.  $Th_g(g_1, g_2) = l$ ,

Example  $\sigma_{cl}(1) = 2$

- For all  $(i = l)$ ,  $Sen(\sigma)(x_g = l) = x_g \geq \sigma_g(l) \wedge x_g < \sigma_g(l + 1)$
- For all  $(i > l)$ ,  $Sen(\sigma)(x_g > l) = x_g \geq \sigma_g(l + 1)$
- For all  $(i < l)$ ,  $Sen(\sigma)(x_g < l) = x_g < \sigma_g(l + 1)$
- Other symbols are handled as usual

## Example



$$Sen(\sigma)(AG(x_{cro} = 1)) = AG(1 \leq x_{cro} < 3)$$

# Reduced model along a signature embedding

$$\sigma : G \rightarrow G'$$

The reduction of a model along a signature embedding is defined up to some restrictions on thresholds.

Given a signature embedding  $\sigma : G \rightarrow G'$  and a  $G'$ -model  $p'$ , the reduced  $G$ -model from  $p'$  denoted  $p'_{|\sigma}$  is defined as follows:

$\forall (i, w) \in \kappa,$

$$p'_{|\sigma}((i, w)) = \begin{cases} Th(i, j) & \text{if } \exists j \in V, \\ & Th'(\sigma(i), \sigma(j)) = \max_{(i,k) \in F} \{ Th'(\sigma(i), \sigma(k)) \mid \\ & Th'(\sigma(i), \sigma(k)) \leq p'((\sigma(i), \sigma(w))) \} \\ 0 & \text{otherwise} \end{cases}$$

CTL-X formulas are preserved through embedding of biological regulatory networks.

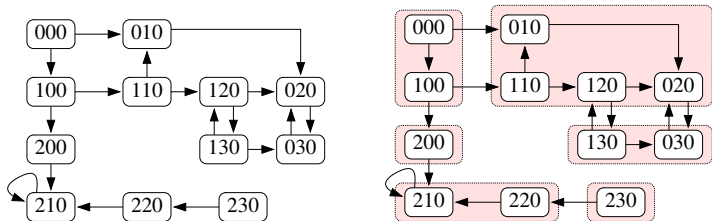
## Theorem

For  $\sigma : G \rightarrow G'$  embedding,  $p'$   $G'$ -model and  $\varphi \in \text{Sen}(G)$ ,

$$p' \models \sigma(\varphi) \iff p'|_{\sigma} \models \varphi$$

Let us consider  $\sigma : G \rightarrow G'$ ,  
 $p'$  a  $G'$ -model,  
 $(S_{G'}, T') = GTA((G', p'))$  its associated ATS  
 $\varphi \in Sen(G)$ .

Let us define a partition of the state space of  $GTA((G', p'))$ ,  
 taking into account shifting of thresholds through the  
 embedding  $\sigma$ .



Let us define the mapping  $B : S_G \rightarrow 2^{S_{G'}}$  as follows:

$\forall s \in S_G$ ,  $B(s) \subseteq S_{G'}$  verifying:  $s' \in B(s)$  if for every  $i$  in  $V$ :

- if  $s(i) = 0$ , then  $s'(\sigma(i)) \geq 0$  and  
 $s'(\sigma(i)) < \min_{(i,k) \in F} \{Th'(\sigma(i), \sigma(k)) \mid Th'(\sigma(i), \sigma(k)) > 0\}$
- otherwise, let  $j$  be any variable in  $G_i^+$  s.t.  $s(i) = Th(i, j)$   
then  $s'(\sigma(i)) \geq Th'(\sigma(i), \sigma(j))$  and  
 $s'(\sigma(i)) < \min_{(i,k) \in F} \{Th'(\sigma(i), \sigma(k)) \mid Th'(\sigma(i), \sigma(k)) > Th'(\sigma(i), \sigma(j))\}$

### Proposition 1

The mapping  $B$  makes a partition of  $S_{G'}$ , i.e.

- 1  $\forall s, s' \in S$ ,  $B(s) \cap B(s') = \emptyset$ , and
- 2  $\bigcup_{s \in S_G} B_s = S_{G'}$ .

A binary relation  $R$  is called a *divergence blind stuttering* (dbs) relation iff it is symmetric and

$$r R s \iff \begin{cases} L(r) = L(s) \\ (r, r') \in T \Rightarrow \exists s_0, s_1, \dots, s_n \text{ finite path, } n \geq 0, (s_0 = s) \\ \quad \wedge (\forall i < n, r R s_i) \wedge r' R s_n \end{cases}$$

The largest dbs relation is an equivalence relation noted  $\simeq_{dbs}$ .

## Proposition 2

Note  $P = \{B(s) \mid s \in S_G\}$ . Then, we have:  
 $P$  is a dbs equivalence.



## Preliminary

The quotient of a transition system  $(S, T)$  by  $\simeq_{dbs}$  is denoted  $(S, T)_{/\simeq_{dbs}}$ .

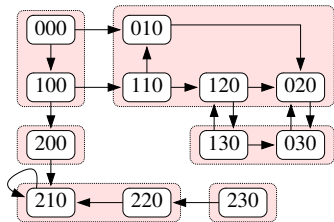
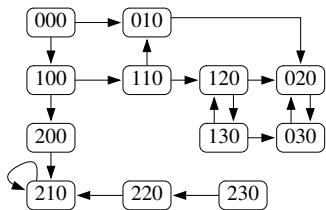
The equivalence relation  $\simeq_{dbs}$  preserves CTL-X formulas, i.e.  $(S, T)$  and  $(S, T)_{/\simeq_{dbs}}$  satisfy the same formulas.

## Proposition 3

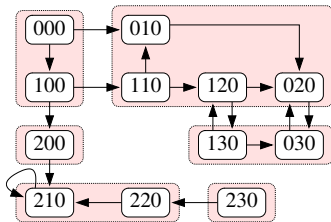
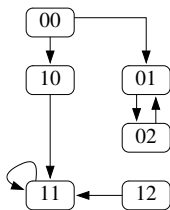
$(S_{G'}, T')_{/\simeq_{dbs}}$  and  $GTA(G, p'_{|\sigma})$  are isomorphic.

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# Sketch of the proof:illustration



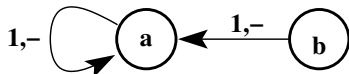
# Sketch of the proof:illustration



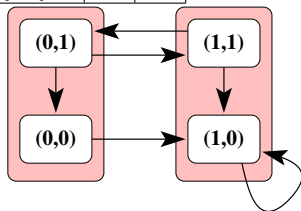
# Counter-example: necessity of restrictive conditions on signature embeddings



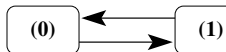
resource $\omega$	$p_a$
$\emptyset$	0
$\{a\}$	1



resource $\omega'$	$p'_a$	$p'_b$
$\emptyset$	0	0
$\{a\}$	1	
$\{b\}$	1	
$\{a, b\}$	1	



resource $\omega'$	$p'_{\sigma_a}$
$\emptyset$	0
$\{a\}$	1



$AG(AF(a = 0))$  is satisfied by  $p'_{\sigma}$  but not by  $p'$ .

- Result :
  - Preservation of properties along simple embedding of biological regulatory networks:  
no new entering edge when embedding a network within a larger one
  - discrete models for BRN modelled as an institution
- Current work : Investigation of some other loose conditions about property preservation for BRN
- Future work : BRN as an application domain to study complex systems.

A system is considered as complex according to the fact that properties of sub-systems are not preserved at the level of the global system.