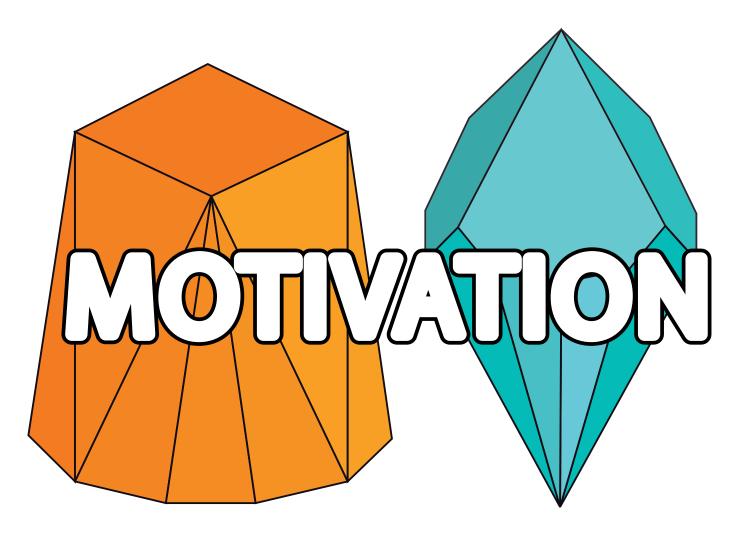


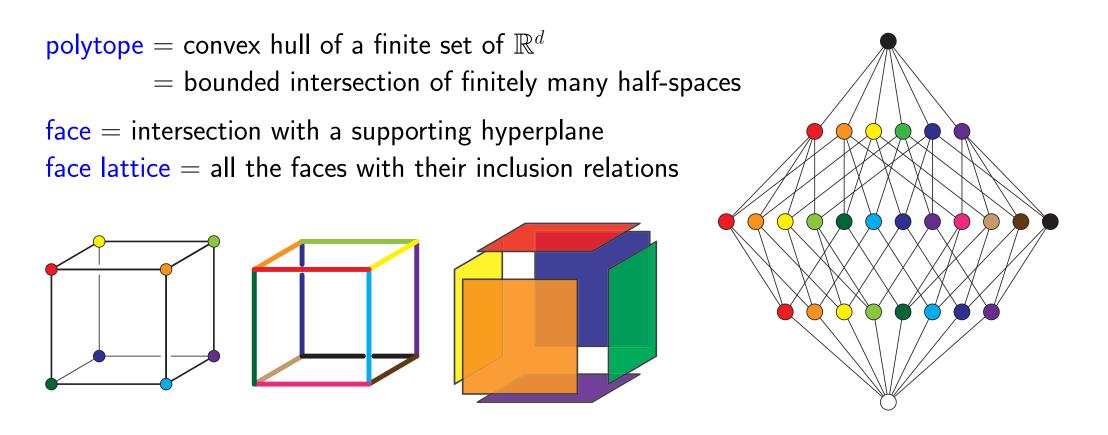
Vincent PILAUD (Université Paris 7)

Francisco SANTOS

(Universidad de Cantabria)



POLYTOPES WITH PRESCRIBED COMBINATORICS



Given a set of points, determine the face lattice of its convex hull.

Given a lattice, is there a polytope which realizes it?

POLYTOPES OF DIMENSION ≥ 4

Polytopes of dimension $3 \leftrightarrow planar 3$ -connected graphs

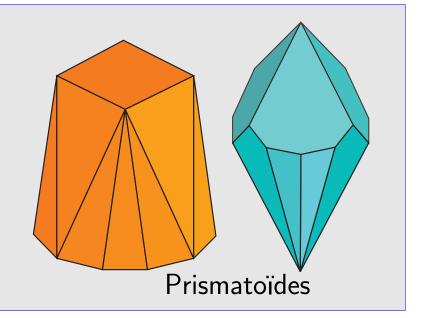
Various open conjectures in dimension 4:

Hirsch conjecture

diameter $\leq \#$ facets – dimension (Santos) complexity of the simplex algorithm

 3^d Conjecture (Kalai)

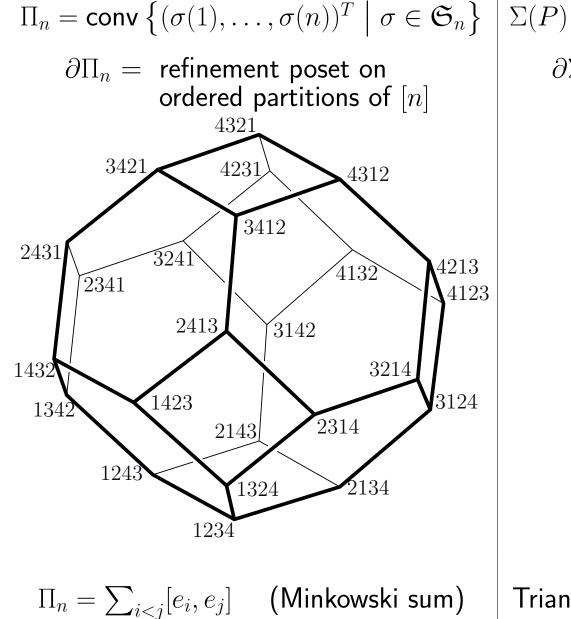
f-vecteur shape (Barany, Ziegler)



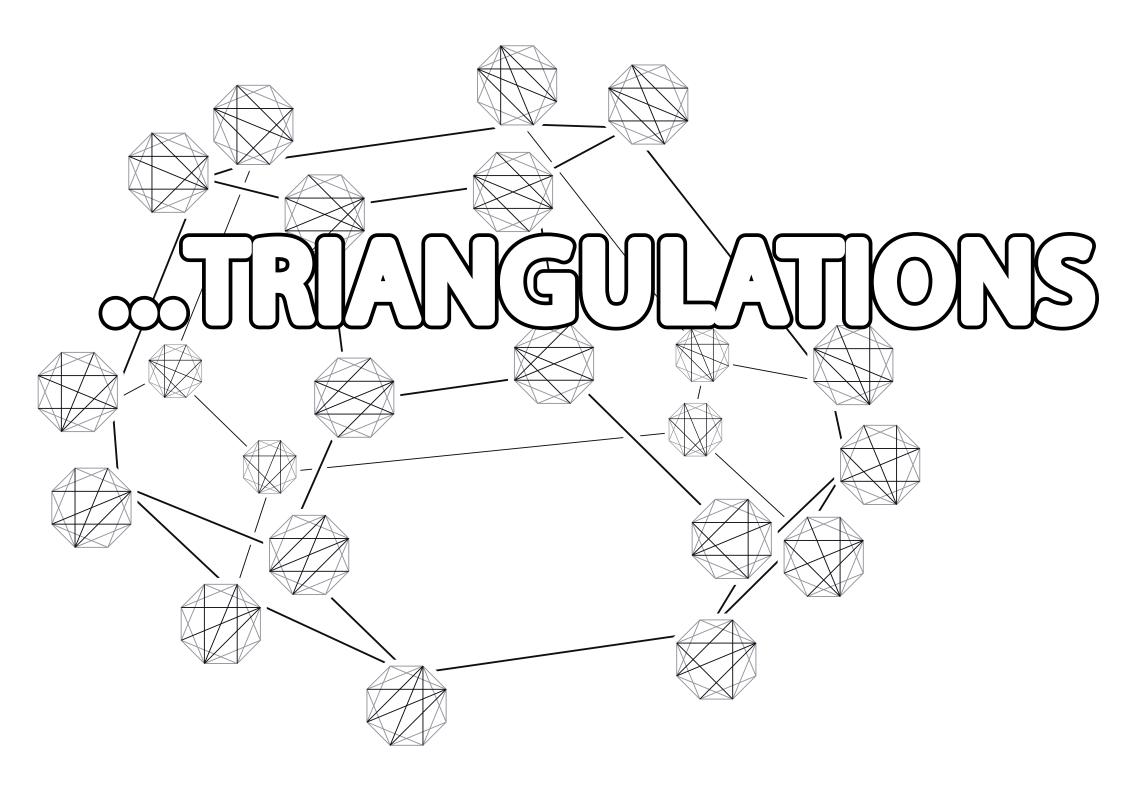
"Our main limits in understanding the combinatorial structure of polytopes still lie in our ability to raise the good questions and in the lack of examples, methods of constructing them, and means of classifying them." Kalai. Handbook of Discrete and Computational Geometry (2004)

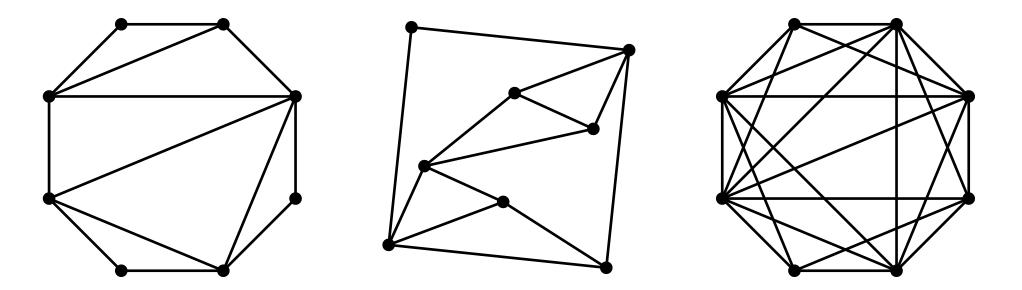
PERMUTAHEDRON

SECONDARY POLYTOPE



 $\Sigma(P) = \operatorname{conv}\left\{\sum_{p \in P} \operatorname{vol}(T, p) e_p \mid T \text{ triang. } P\right\}$ $\partial \Sigma(P) =$ refinement poset on regular polyhedral subdivisions of PTriangulations are non-regular

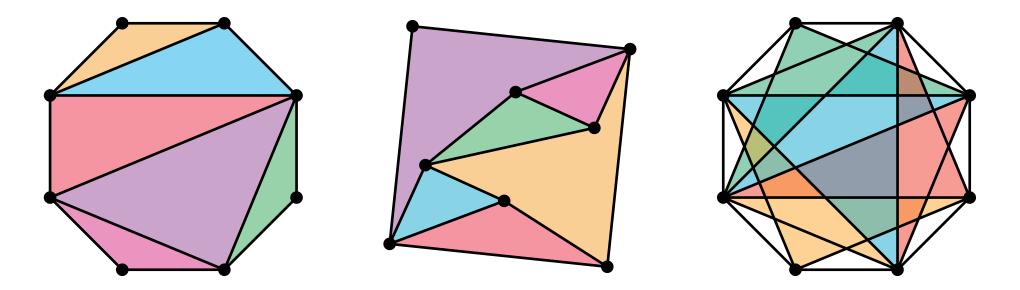




triangulation = maximal crossing-free set of edges

pseudotriangulation = maximal crossing-free pointed set of edges

k-triangulation = maximal (k + 1)-crossing-free set of edges

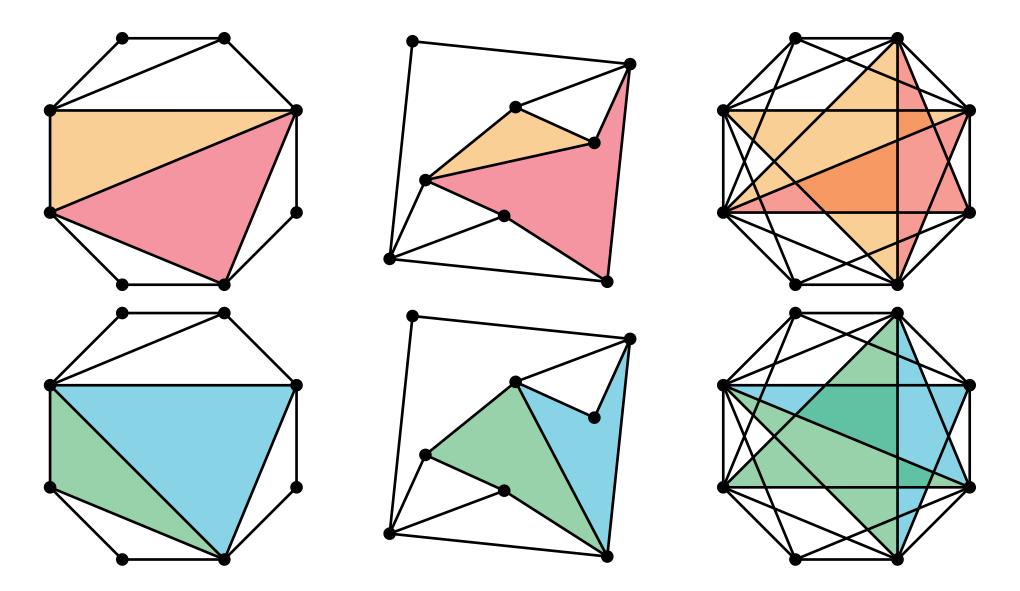


triangulation = maximal crossing-free set of edges

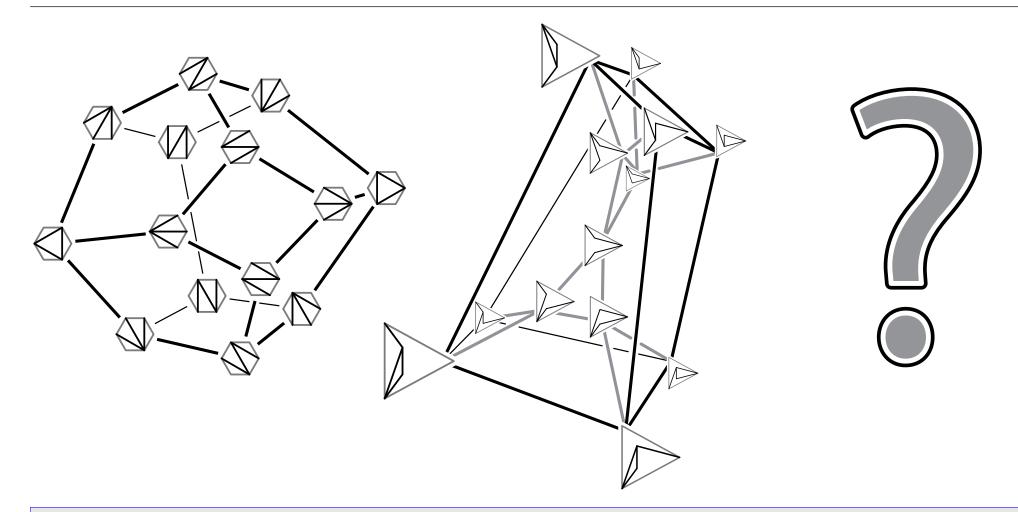
= decomposition into triangles

- pseudotriangulation = maximal crossing-free pointed set of edges = decomposition into pseudotriangles
 - k-triangulation = maximal (k + 1)-crossing-free set of edges = decomposition into k-stars

VP & F. Santos, Multitriangulations as complexes of star polygons, 2009.

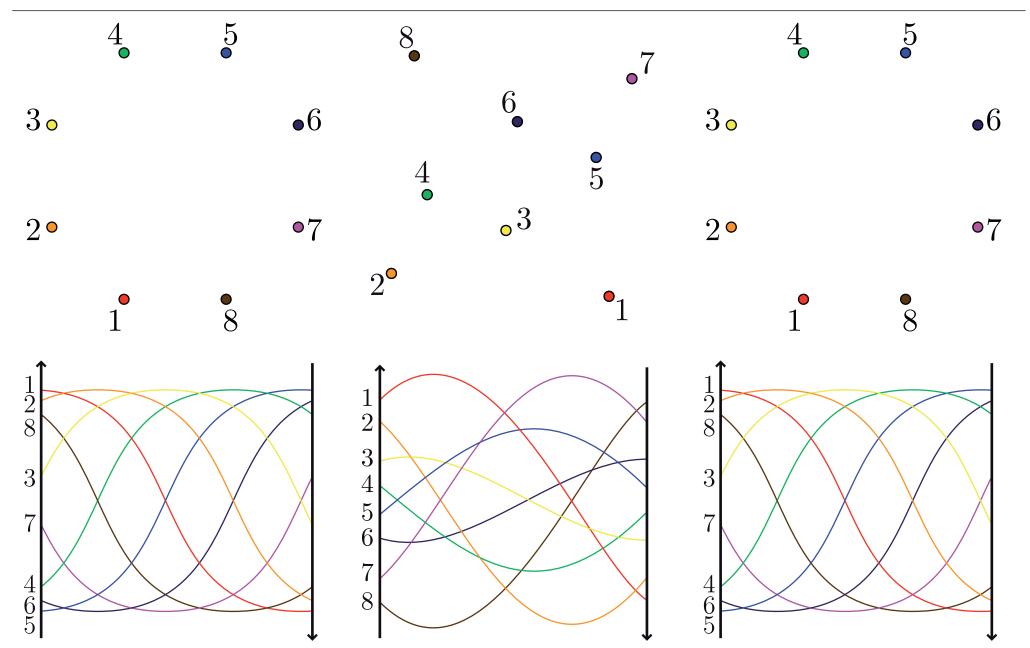


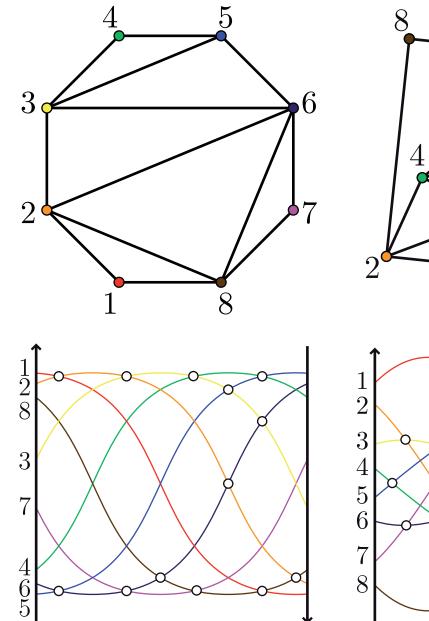
flip = exchange an internal edge with the common bisector of the two adjacent cells

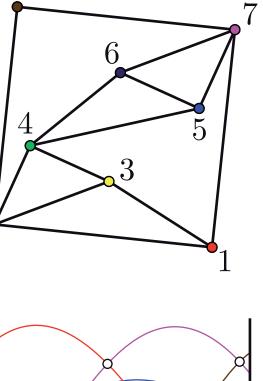


- \rightarrow crossing-free sets of internal edges
 - pointed crossing-free sets of internal edges
 - (k+1)-crossing-free sets of k-internal edges
- associahedron \longleftrightarrow
- pseudotriangulations polytope $\leftrightarrow \rightarrow$
 - multiassociahedron \leftrightarrow









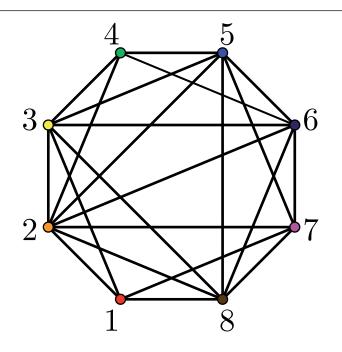
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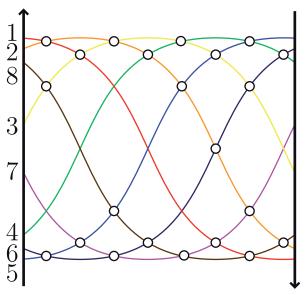
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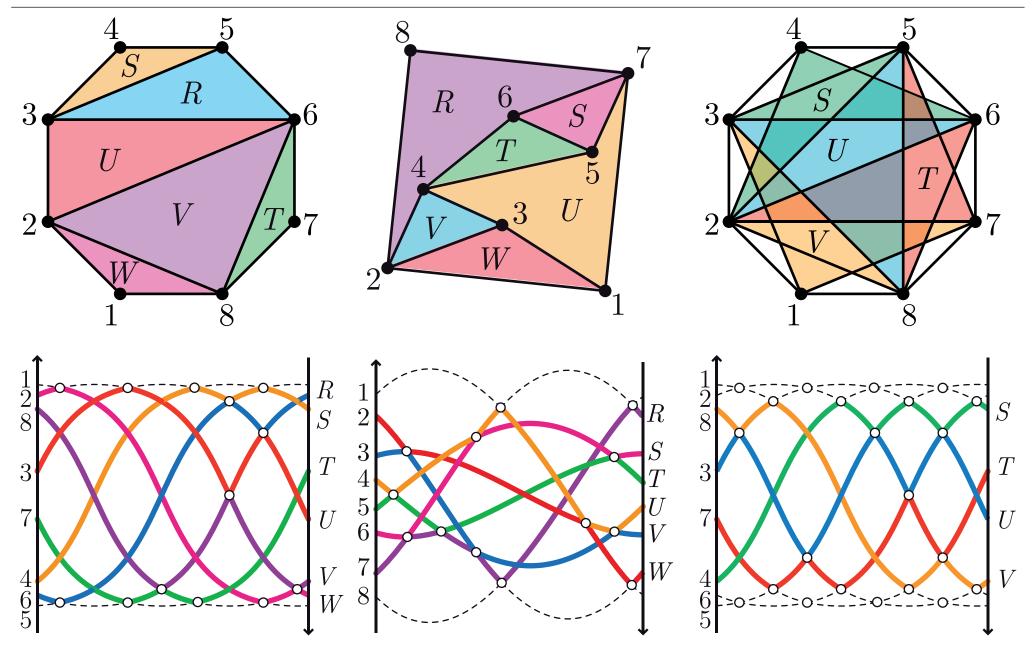
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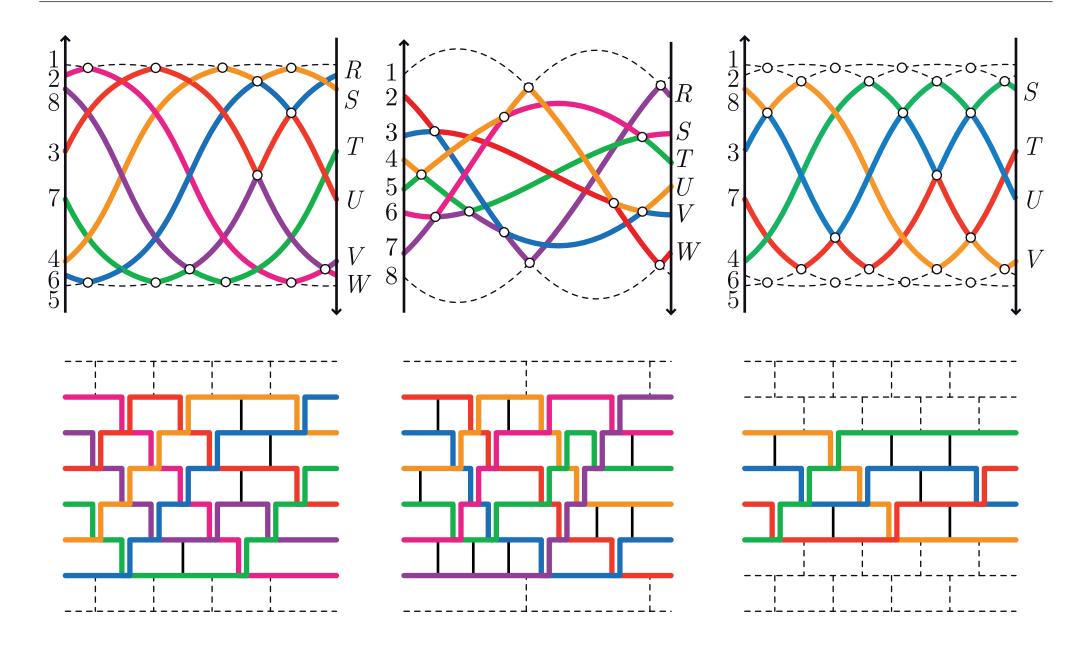
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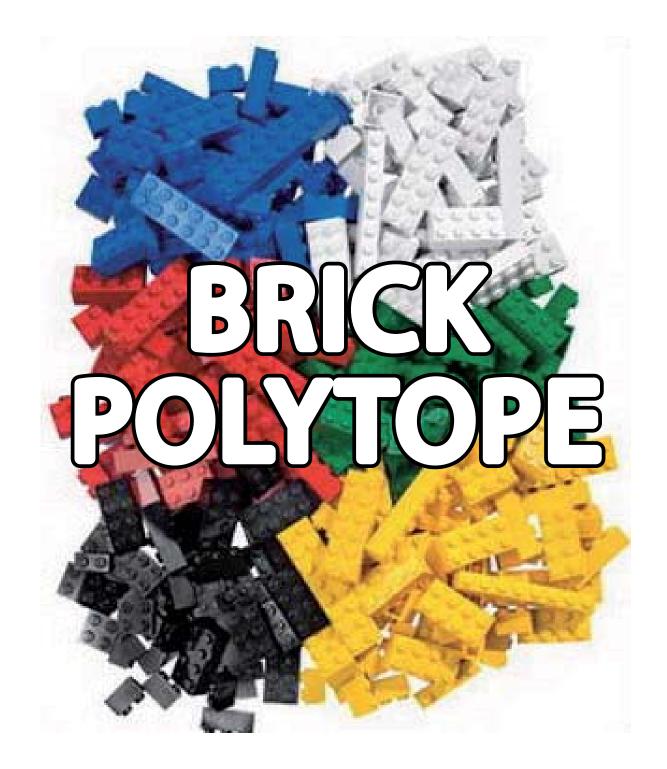




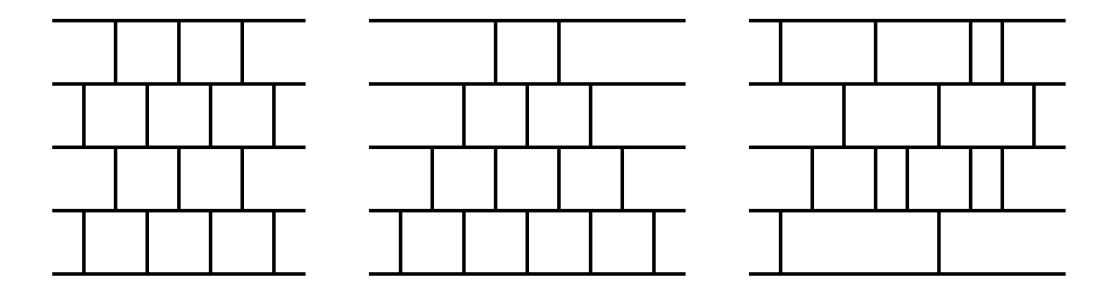


VP & M. Pocchiola, Multitriangulations, pseudotriangulations and primitive sorting networks, 2011.



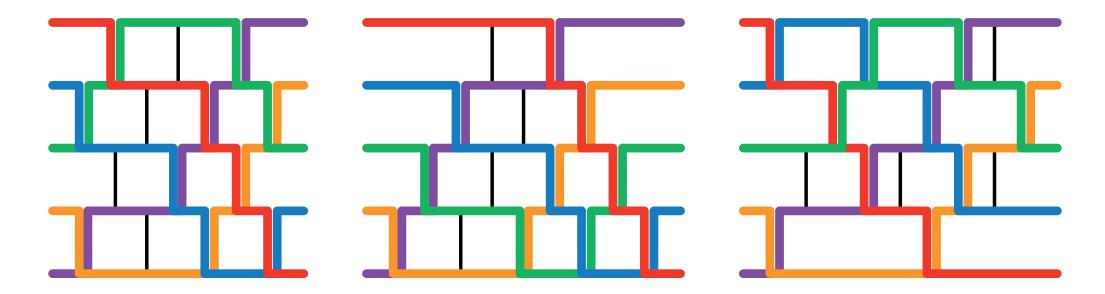


NETWORKS & PSEUDOLINE ARRANGEMENTS

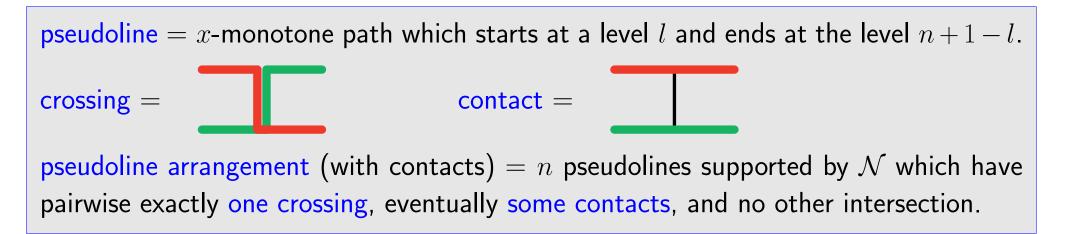


network $\mathcal{N} = n$ horizontal levels and m vertical commutators. bricks of $\mathcal{N} =$ bounded cells.

NETWORKS & PSEUDOLINE ARRANGEMENTS

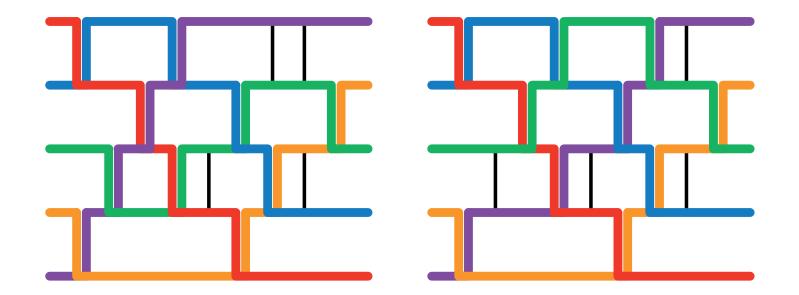


network $\mathcal{N} = n$ horizontal levels and m vertical commutators. bricks of $\mathcal{N} =$ bounded cells.



FLIPS

flip = exchange a contact with the corresponding crossing.



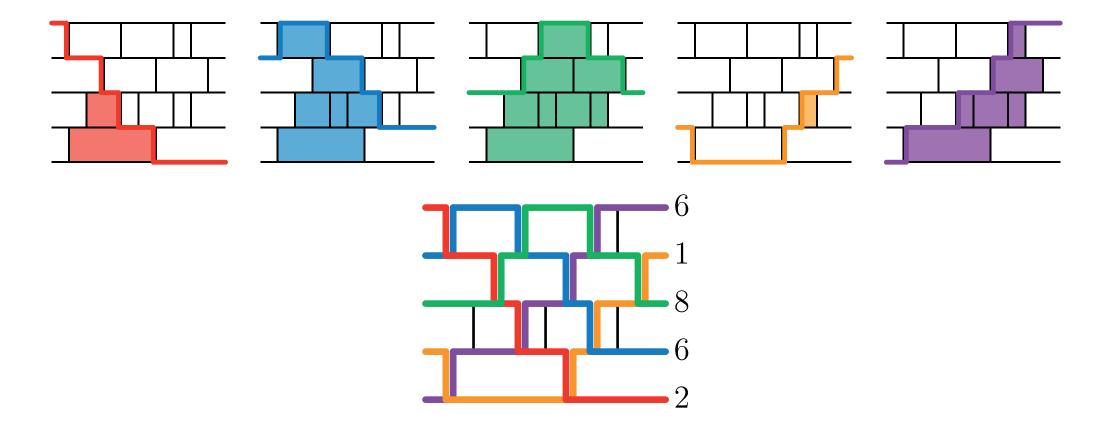
THEOREM. Let \mathcal{N} be a sorting network with n levels and m commutators. The graph of flips $G(\mathcal{N})$ is $\left(m - \binom{n}{2}\right)$ -regular and connected.

QUESTION. Is $G(\mathcal{N})$ the graph of a simple $\left(m - \binom{n}{2}\right)$ -dimensional polytope?

A. Knutson & E. Miller, Subword complexes in Coxeter groups, 2004.

BRICK POLYTOPE

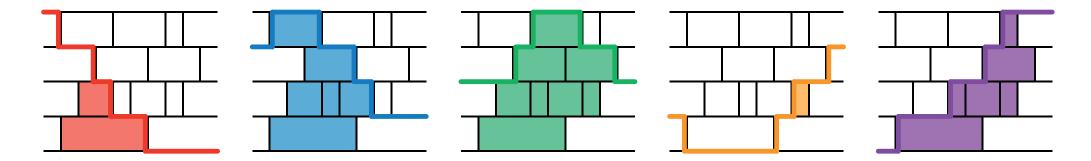
Λ pseudoline arrangement supported by \mathcal{N} → brick vector $ω(Λ) ∈ \mathbb{R}^n$. $ω(Λ)_j =$ number of bricks of \mathcal{N} below the *j*th pseudoline of Λ.



Brick polytope $\Omega(\mathcal{N}) = \operatorname{conv} \{ \omega(\Lambda) \mid \Lambda \text{ pseudoline arrangement supported by } \mathcal{N} \}.$

BRICK POLYTOPE

Λ pseudoline arrangement supported by \mathcal{N} → brick vector $ω(Λ) ∈ \mathbb{R}^n$. $ω(Λ)_j =$ number of bricks of \mathcal{N} below the *j*th pseudoline of Λ.



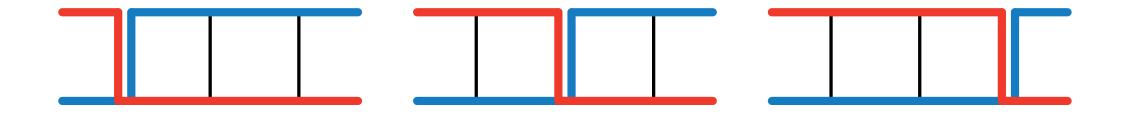
Brick polytope $\Omega(\mathcal{N}) = \operatorname{conv} \{ \omega(\Lambda) \mid \Lambda \text{ pseudoline arrangement supported by } \mathcal{N} \}.$

REMARK. The brick polytope is not full-dimensional:

$$\Omega(\mathcal{N}) \subset \left\{ \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \middle| \sum_{i=1}^n x_i = \sum_{b \text{ brick of } \mathcal{N}} \operatorname{depth}(b) \right\}$$

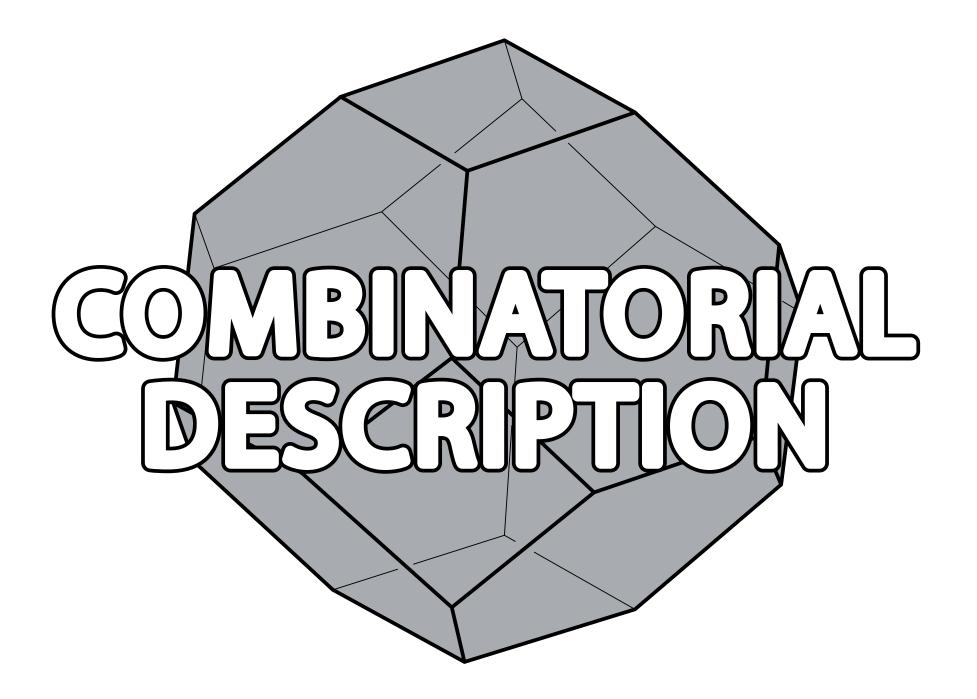
EXAMPLE: 2-LEVELS NETWORKS

 \mathcal{X}_m = network with two levels and m commutators.

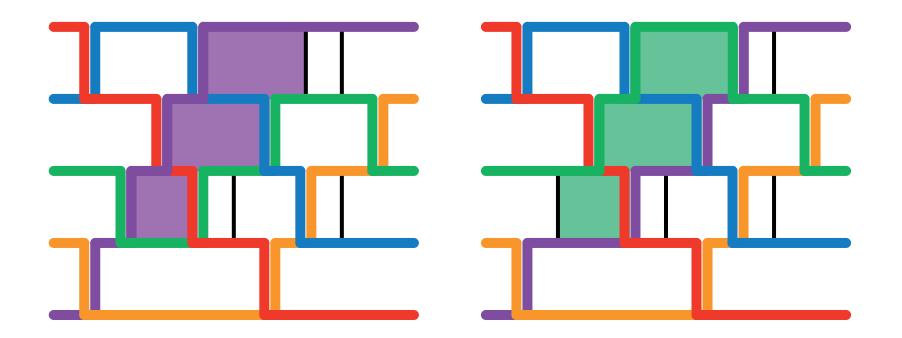


Graph of flips $G(\mathcal{X}_m) = \text{complete graph } K_m$.

Brick polytope
$$\Omega(\mathcal{X}_m) = \operatorname{conv}\left\{ \begin{pmatrix} m-i\\ i-1 \end{pmatrix} \middle| i \in [m] \right\} = \left[\begin{pmatrix} m-1\\ 0 \end{pmatrix}, \begin{pmatrix} 0\\ m-1 \end{pmatrix} \right].$$



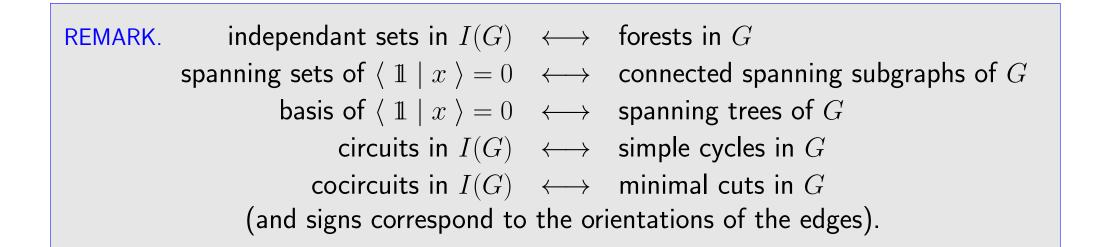
BRICK VECTORS AND FLIPS



REMARK. If Λ and Λ' are two pseudoline arrangements supported by \mathcal{N} and related by a flip between their *i*th and *j*th pseudolines, then $\omega(\Lambda) - \omega(\Lambda') \in \mathbb{N}_{>0}(e_j - e_i)$.

COROLLARY. The cone generated by the vector configuration

 $\{e_j - e_i \mid \text{there is a contact between the } ith and jth pseudolines of \Lambda\}$ is contained in the cone of the brick polytope $\Omega(S)$ at the brick vector $\omega(\Lambda)$. $G \text{ directed (multi)graph} \mapsto \text{Incidence configuration } I(G) = \{e_j - e_i \mid (i, j) \in G\}, \\ \mapsto \text{Incidence cone } C(G) = \text{cone generated by } I(G).$



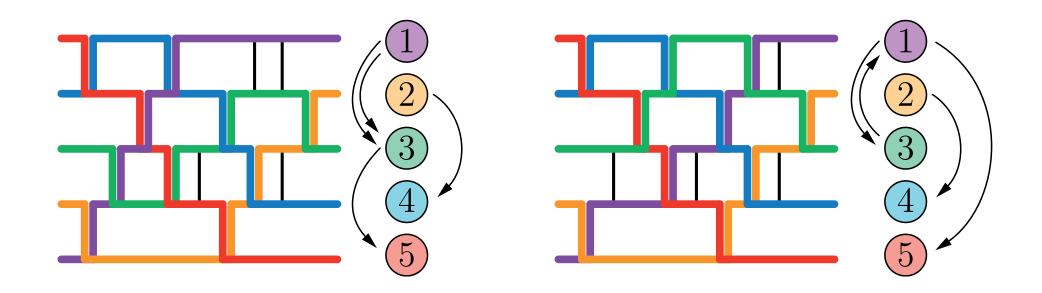
REMARK. *H* subgraph of *G*. Then I(H) forms a *k*-face of $C(G) \iff H$ has n - k connected components and G/H is acyclic. In particular:

 $\begin{array}{rcl} C(G) \text{ is pointed } & \longleftrightarrow & G \text{ is acyclic} \\ \text{ facets of } C(G) & \longleftrightarrow & \text{ complements of the minimal directed cuts of } G \end{array}$

CONTACT GRAPH OF A PSEUDOLINE ARRANGEMENT

Contact graph $\Lambda^{\#}$ of a pseudoline arrangement $\Lambda =$

- a node for each pseudoline of $\Lambda,$ and
- an arc for each contact point of Λ oriented from top to bottom.



THEOREM. The cone of the brick polytope $\Omega(S)$ at the brick vector $\omega(\Lambda)$ is the incidence cone $C(\Lambda^{\#})$ of the contact graph of Λ .

COMBINATORIAL DESCRIPTION

THEOREM. The cone of the brick polytope $\Omega(S)$ at the brick vector $\omega(\Lambda)$ is the incidence cone $C(\Lambda^{\#})$ of the contact graph of Λ .

VERTICES OF $\Omega(\mathcal{N})$

The brick vector $\omega(\Lambda)$ is a vertex of $\Omega(\mathcal{N}) \iff$ the contact graph $\Lambda^{\#}$ is acyclic.

$\mathsf{GRAPH}\ \mathsf{OF}\ \Omega(\mathcal{N})$

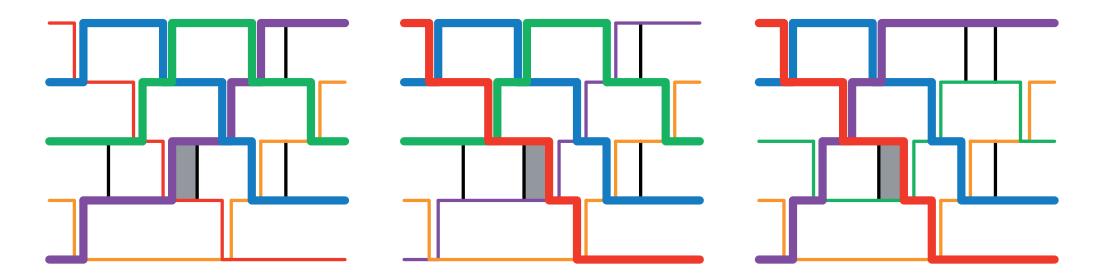
The graph of the brick polytope is a subgraph of $G(\mathcal{N})$ whose vertices are the pseudoline arrangements with acyclic contact graphs.

FACETS OF $\Omega(\mathcal{N})$

The facets of $\Omega(\mathcal{N})$ correspond to the minimal directed cuts of the contact graphs of the pseudoline arrangements supported by \mathcal{N} .

BRICK POLYTOPES AND MINKOWSKI SUMS

 \mathcal{N} network with n levels, b a brick of \mathcal{N} , Λ pseudoline arrangement supported by \mathcal{N} . $\omega(\Lambda, b) \in \mathbb{R}^n$ characteristic vector of the pseudolines of Λ passing above b. $\Omega(\mathcal{N}, b) = \operatorname{conv} \{\omega(\Lambda, b) \mid \Lambda \text{ pseudoline arrangement supported by } \mathcal{N} \} \subset \mathbb{R}^n$.



THEOREM. The brick polytope $\Omega(\mathcal{N})$ is the Minkowski sum of the polytopes $\Omega(\mathcal{N}, b)$ associated to the bricks of \mathcal{N} :

$$\Omega(\mathcal{N}) = \sum_{b \text{ brick of } \mathcal{N}} \Omega(\mathcal{N}, b).$$

BRICK POLYTOPES AND GENERALIZED PERMUTOHEDRA

Generalized permutohedra = polytope whose inequality description is of the form

$$Z\left(\{z_I\}_{I\in[n]}\right) = \left\{ \left(\begin{array}{c} x_1\\ \dots\\ x_n \end{array} \right) \in \mathbb{R}^n \ \middle| \ \sum_{i=1}^n x_i = z_{[n]} \text{ and } \sum_{i\in I} x_i \ge z_I \text{ for } I \subset [n] \right\}$$

for some family $\{z_I\}_{I \subset [n]} \in \mathbb{R}^{2^{[n]}}$.

A. Postnikov, Permutohedra, associahedra and beyond, 2009.

THEOREM. Any generalized permutahedron is a Minkowski sum of simplices:

$$Z\left(\{z_I\}_{I\in[n]}\right) = \sum_{I\subset[n]} y_I \Delta_I \quad \text{where} \quad y_I = \sum_{J\subset I} (-1)^{|I\smallsetminus J|} z_J \quad \left(\text{ie.} \quad z_I = \sum_{J\subset I} y_J\right).$$

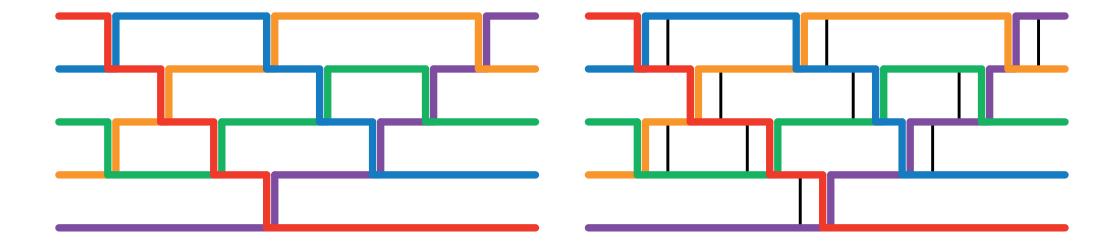
F. Ardila, C. Benedetti & J. Doker, Matroid polytopes and their volumes, 2010.

REMARK. All brick polytopes are generalized permutohedra. Compute $\{y_I\}_{I \subset [n]}$. Which generalized permutohedra are brick polytopes?

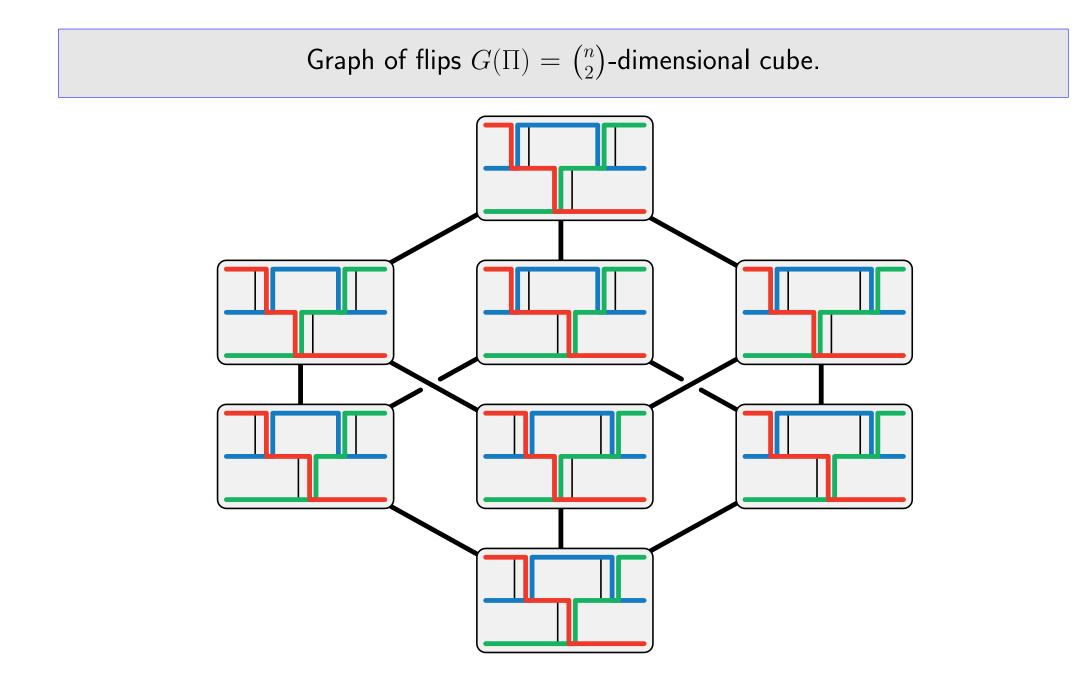


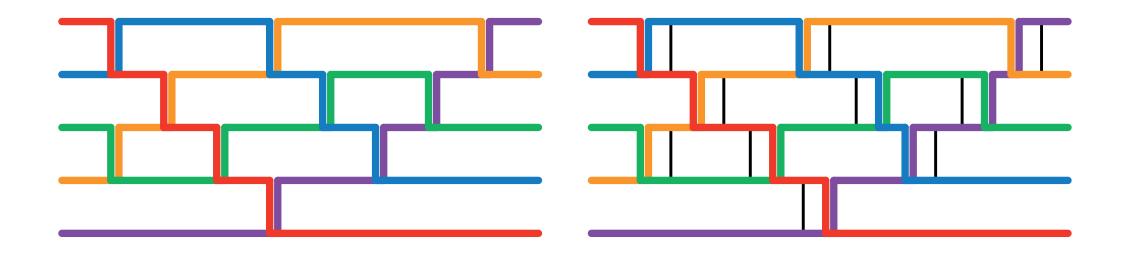
Reduced network = network with n levels and $\binom{n}{2}$ commutators. It supports only one pseudoline arrangement.

Duplicated network Π = network with n levels and $2\binom{n}{2}$ commutators obtained by duplicating each commutator of a reduced network.

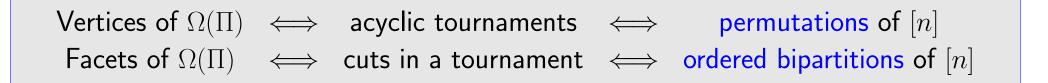


Any pseudoline arrangement supported by Π has one contact and one crossing among each pair of duplicated commutators.

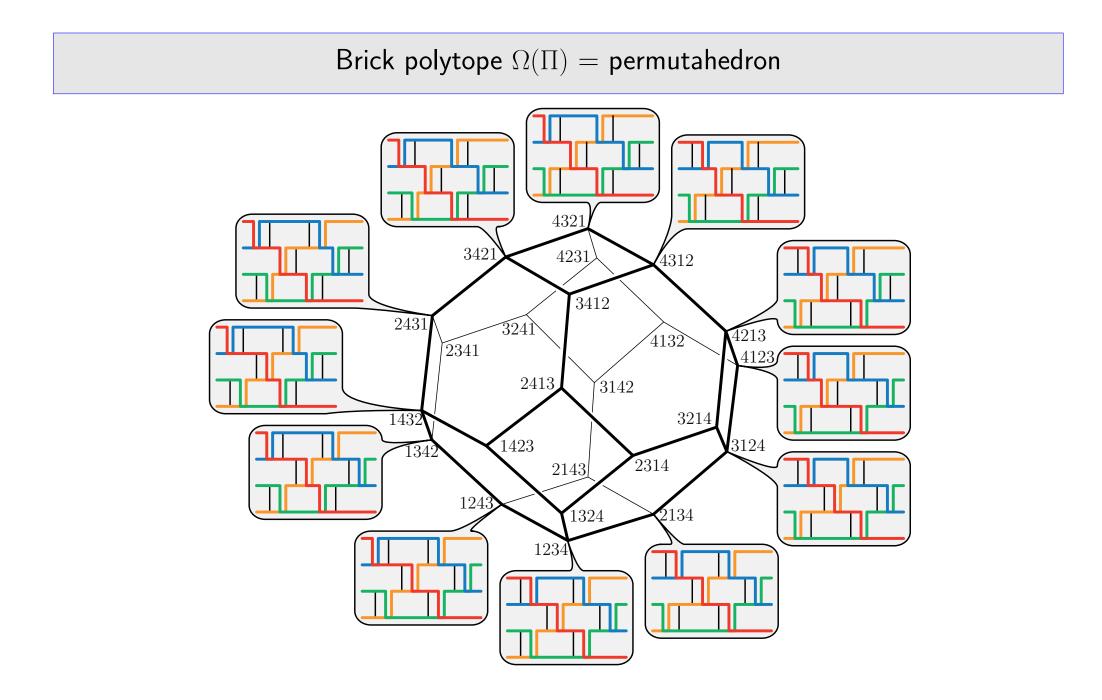


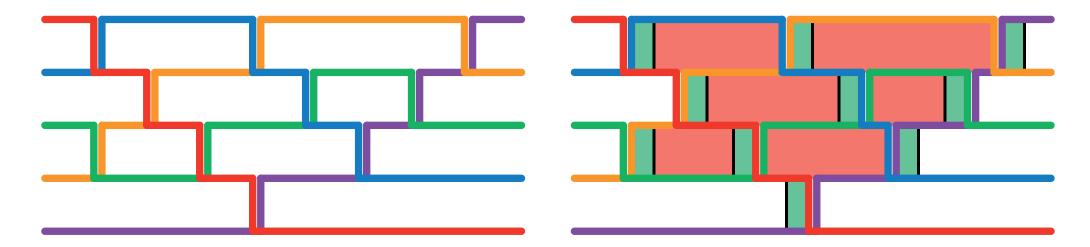


Any pseudoline arrangement supported by Π has one contact and one crossing among each pair of duplicated commutators. \implies The contact graph $\Lambda^{\#}$ is a tournament.



Brick polytope $\Omega(\Pi) = permutahedron$





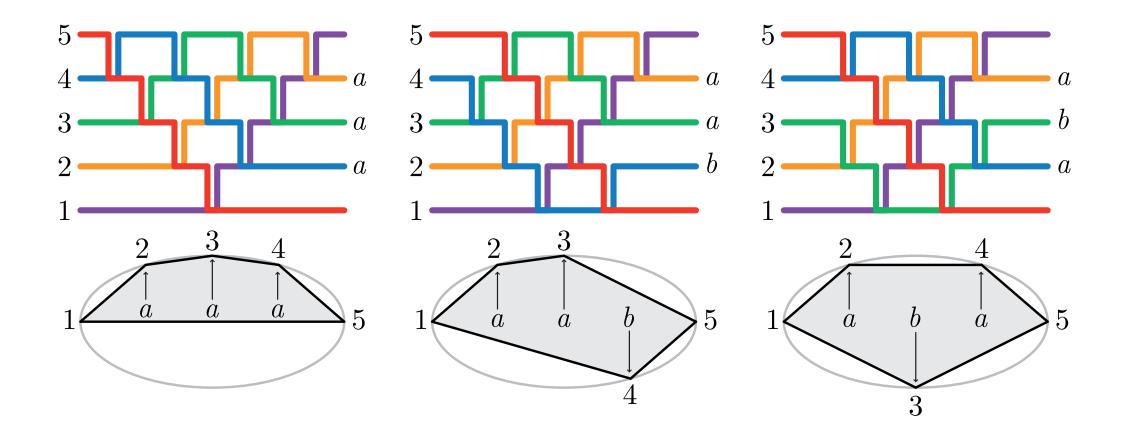
Minkowski sum decomposition

$$\Omega(\Pi) = \sum_{b \text{ brick of } \Pi} \Omega(\Pi, b) = \sum_{i < j} \text{segment } [e_i - e_j] + \sum \text{vertices} = \text{permutahedron}$$

$$\begin{split} P(0,1,\ldots,n-1) &= \mathsf{Newton}\left(\det\left[t_i^{j-1}\right]_{i,j\in[n]}\right) = \mathsf{Newton}\left(\prod_{1\leq i< j\leq n}(t_j-t_i)\right) \\ &= \sum_{1\leq i< j\leq n}\mathsf{Newton}\left(t_j-t_i\right) = \sum_{1\leq i< j\leq n}[e_j-e_i] \end{split}$$

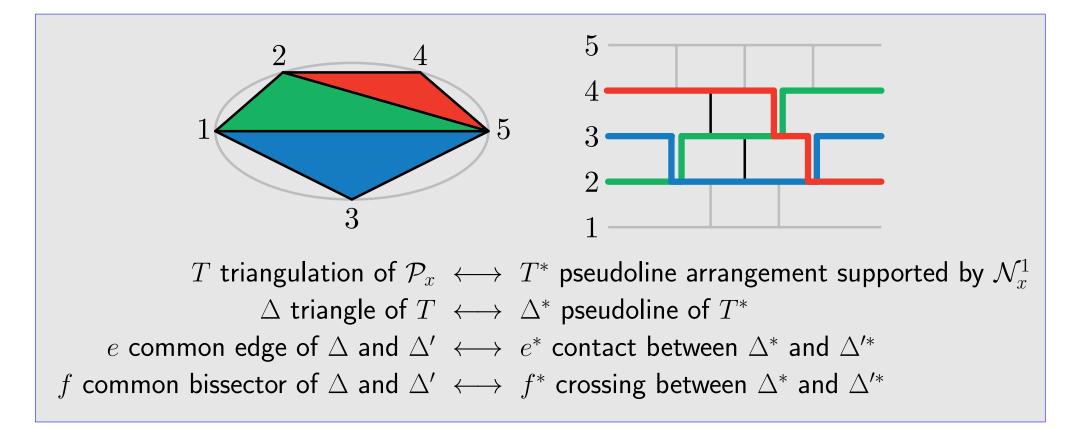
ALTERNATING NETWORK: ASSOCIAHEDRA

For $x \in \{a, b\}^{n-2}$, we define a reduced alternating network \mathcal{N}_x and a polygon \mathcal{P}_x .



 \mathcal{N}_x is the dual pseudoline arrangement of the polygon \mathcal{P}_x .

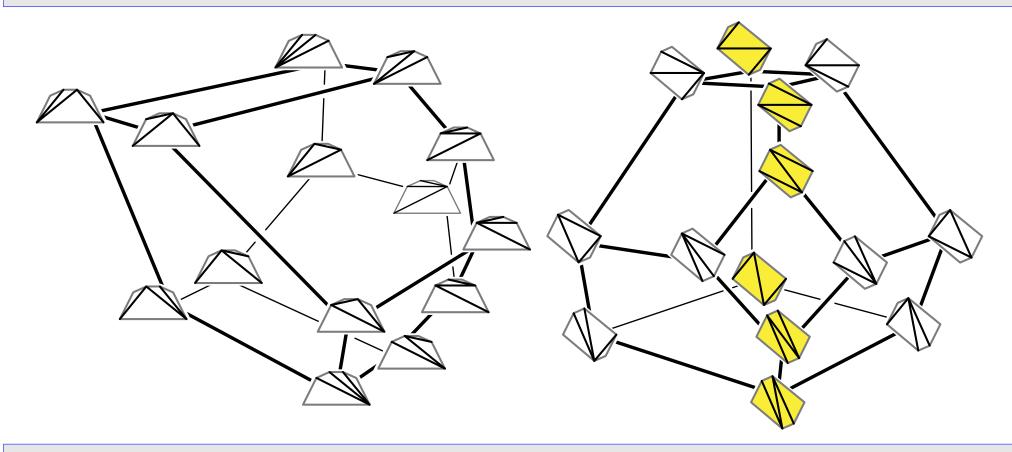
THEOREM. There is a duality between the pseudoline arrangements supported by \mathcal{N}_x^1 and the triangulations of the polygon \mathcal{P}_x .



COROLLARY. (i) The graph of flips $G(\mathcal{N}_x^1)$ is (isomorphic to) the graph of flips $G(\mathcal{P}_x)$. (ii) The contact graph $(T^*)^{\#}$ is (isomorphic to) the dual binary tree of T.

HOHLWEG & LANGE'S ASSOCIAHEDRA

THEOREM. For any word $x \in \{a, b\}^{n-2}$, the simplicial complex of crossing-free sets of internal diagonals of the convex *n*-gon \mathcal{P}_x is (isomorphic to) the boundary complex of the polar of the brick polytope $\Omega(\mathcal{N}_x^1)$.



REMARK. Up to translation, we obtain Hohlweg & Lange's associahedra.

C. Hohlweg & C. Lange, Realizations of the associahedron and cyclohedron, 2007.

arXiv:1103.2731

