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## POLYTOPES WITH PRESCRIBED COMBINATORICS

polytope $=$ convex hull of a finite set of $\mathbb{R}^{d}$
$=$ bounded intersection of finitely many half-spaces face $=$ intersection with a supporting hyperplane face lattice $=$ all the faces with their inclusion relations


Given a set of points, determine the face lattice of its convex hull.

Given a lattice, is there a polytope which realizes it?

## POLYTOPES OF DIMENSION $\geq 4$

## Polytopes of dimension $3 \longleftrightarrow$ planar 3-connected graphs

Various open conjectures in dimension 4:

## Hirsch conjecture

 diameter $\leq \#$ facets - dimension (Santos) complexity of the simplex algorithm$3^{d}$ Conjecture (Kalai)
$f$-vecteur shape (Barany, Ziegler)


Prismatoïdes
"Our main limits in understanding the combinatorial structure of polytopes still lie in our ability to raise the good questions and in the lack of examples, methods of constructing them, and means of classifying them." Kalai. Handbook of Discrete and Computational Geometry (2004)


$$
\Pi_{n}=\sum_{i<j}\left[e_{i}, e_{j}\right]
$$

(Minkowski sum)

## SECONDARY POLYTOPE

$\Sigma(P)=\operatorname{conv}\left\{\sum_{p \in P} \operatorname{vol}(T, p) e_{p} \mid T\right.$ triang. $\left.P\right\}$ $\partial \Sigma(P)=$ refinement poset on regular polyhedral subdivisions of $P$


Triangulations
 are non-regular


## THREE GEOMETRIC STRUCTURES


triangulation $=$ maximal crossing-free set of edges
pseudotriangulation $=$ maximal crossing-free pointed set of edges
$k$-triangulation $=$ maximal $(k+1)$-crossing-free set of edges

## THREE GEOMETRIC STRUCTURES



$$
\begin{aligned}
\text { triangulation } & =\text { maximal crossing-free set of edges } \\
& =\text { decomposition into triangles } \\
\text { pseudotriangulation } & =\text { maximal crossing-free pointed set of edges } \\
& =\text { decomposition into pseudotriangles } \\
k \text {-triangulation } & =\text { maximal }(k+1) \text {-crossing-free set of edges } \\
& =\text { decomposition into } k \text {-stars }
\end{aligned}
$$

## THREE GEOMETRIC STRUCTURES


flip $=$ exchange an internal edge with the common bisector of the two adjacent cells

## THREE GEOMETRIC STRUCTURES


associahedron
pseudotriangulations polytope
multiassociahedron

$\longleftrightarrow$
crossing-free sets of internal edges pointed crossing-free sets of internal edges ( $k+1$ )-crossing-free sets of $k$-internal edges


## DUALITY



DUALITY


## DUALITY



VP \& M. Pocchiola, Multitriangulations, pseudotriangulations and primitive sorting networks, 2011.

## DUALITY




## NETWORKS \& PSEUDOLINE ARRANGEMENTS


network $\mathcal{N}=n$ horizontal levels and $m$ vertical commutators.
bricks of $\mathcal{N}=$ bounded cells.

## NETWORKS \& PSEUDOLINE ARRANGEMENTS



```
network \mathcal{N}=n}\mathrm{ horizontal levels and m}\mathrm{ vertical commutators.
bricks of \mathcal{N}=\mathrm{ bounded cells.}
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pseudoline $=x$-monotone path which starts at a level $l$ and ends at the level $n+1-l$. crossing $=$


$$
\text { contact }=
$$


pseudoline arrangement (with contacts) $=n$ pseudolines supported by $\mathcal{N}$ which have pairwise exactly one crossing, eventually some contacts, and no other intersection.

## FLIPS

flip $=$ exchange a contact with the corresponding crossing.


THEOREM. Let $\mathcal{N}$ be a sorting network with $n$ levels and $m$ commutators. The graph of flips $G(\mathcal{N})$ is $\left(m-\binom{n}{2}\right)$-regular and connected.

QUESTION. Is $G(\mathcal{N})$ the graph of a simple $\left(m-\binom{n}{2}\right)$-dimensional polytope?
A. Knutson \& E. Miller, Subword complexes in Coxeter groups, 2004.

## BRICK POLYTOPE

$\Lambda$ pseudoline arrangement supported by $\mathcal{N} \longmapsto$ brick vector $\omega(\Lambda) \in \mathbb{R}^{n}$. $\omega(\Lambda)_{j}=$ number of bricks of $\mathcal{N}$ below the $j$ th pseudoline of $\Lambda$.


Brick polytope $\Omega(\mathcal{N})=\operatorname{conv}\{\omega(\Lambda) \mid \Lambda$ pseudoline arrangement supported by $\mathcal{N}\}$.

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REMARK. The brick polytope is not full-dimensional:

$$
\Omega(\mathcal{N}) \subset\left\{\left.\left(\begin{array}{c}
x_{1} \\
\vdots \\
x_{n}
\end{array}\right) \right\rvert\, \sum_{i=1}^{n} x_{i}=\sum_{b \text { brick of } \mathcal{N}} \operatorname{depth}(b)\right\}
$$

## EXAMPLE: 2-LEVELS NETWORKS

$$
\mathcal{X}_{m}=\text { network with two levels and } m \text { commutators. }
$$



Graph of flips $G\left(\mathcal{X}_{m}\right)=$ complete graph $K_{m}$.

Brick polytope $\Omega\left(\mathcal{X}_{m}\right)=\operatorname{conv}\left\{\left.\binom{m-i}{i-1} \right\rvert\, i \in[m]\right\}=\left[\binom{m-1}{0},\binom{0}{m-1}\right]$.


## BRICK VECTORS AND FLIPS



REMARK. If $\Lambda$ and $\Lambda^{\prime}$ are two pseudoline arrangements supported by $\mathcal{N}$ and related by a flip between their $i$ th and $j$ th pseudolines, then $\omega(\Lambda)-\omega\left(\Lambda^{\prime}\right) \in \mathbb{N}_{>0}\left(e_{j}-e_{i}\right)$.

COROLLARY. The cone generated by the vector configuration
$\left\{e_{j}-e_{i} \mid\right.$ there is a contact between the $i$ th and $j$ th pseudolines of $\left.\Lambda\right\}$
is contained in the cone of the brick polytope $\Omega(S)$ at the brick vector $\omega(\Lambda)$.

## INCIDENCE CONE OF A DIRECTED MULTIGRAPH

$G$ directed (multi)graph $\longmapsto$ Incidence configuration $I(G)=\left\{e_{j}-e_{i} \mid(i, j) \in G\right\}$, $\longmapsto$ Incidence cone $C(G)=$ cone generated by $I(G)$.

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REMARK. independant sets in I(G) \longleftrightarrow forests in G
    spanning sets of }\langle\mathbb{1}|x\rangle=0\longleftrightarrow\mathrm{ connected spanning subgraphs of }
    basis of }\langle\mathbb{1}|x\rangle=0 \longleftrightarrow spanning trees of G
                circuits in I(G) \longleftrightarrow simple cycles in G
            cocircuits in I(G) \longleftrightarrow minimal cuts in G
        (and signs correspond to the orientations of the edges).
```

REMARK. $H$ subgraph of $G$. Then $I(H)$ forms a $k$-face of $C(G) \Longleftrightarrow H$ has $n-k$ connected components and $G / H$ is acyclic. In particular:
$C(G)$ is pointed $\longleftrightarrow G$ is acyclic facets of $C(G) \longleftrightarrow$ complements of the minimal directed cuts of $G$

## CONTACT GRAPH OF A PSEUDOLINE ARRANGEMENT

Contact graph $\Lambda^{\#}$ of a pseudoline arrangement $\Lambda=$

- a node for each pseudoline of $\Lambda$, and
- an arc for each contact point of $\Lambda$ oriented from top to bottom.


THEOREM. The cone of the brick polytope $\Omega(S)$ at the brick vector $\omega(\Lambda)$ is the incidence cone $C\left(\Lambda^{\#}\right)$ of the contact graph of $\Lambda$.

## COMBINATORIAL DESCRIPTION

THEOREM. The cone of the brick polytope $\Omega(S)$ at the brick vector $\omega(\Lambda)$ is the incidence cone $C\left(\Lambda^{\#}\right)$ of the contact graph of $\Lambda$.

## VERTICES OF $\Omega(\mathcal{N})$

The brick vector $\omega(\Lambda)$ is a vertex of $\Omega(\mathcal{N}) \Longleftrightarrow$ the contact graph $\Lambda^{\#}$ is acyclic.

## GRAPH OF $\Omega(\mathcal{N})$

The graph of the brick polytope is a subgraph of $G(\mathcal{N})$ whose vertices are the pseudoline arrangements with acyclic contact graphs.

## FACETS OF $\Omega(\mathcal{N})$

The facets of $\Omega(\mathcal{N})$ correspond to the minimal directed cuts of the contact graphs of the pseudoline arrangements supported by $\mathcal{N}$.

## BRICK POLYTOPES AND MINKOWSKI SUMS

$\mathcal{N}$ network with $n$ levels, $b$ a brick of $\mathcal{N}, \Lambda$ pseudoline arrangement supported by $\mathcal{N}$. $\omega(\Lambda, b) \in \mathbb{R}^{n}$ characteristic vector of the pseudolines of $\Lambda$ passing above $b$. $\Omega(\mathcal{N}, b)=\operatorname{conv}\{\omega(\Lambda, b) \mid \Lambda$ pseudoline arrangement supported by $\mathcal{N}\} \subset \mathbb{R}^{n}$.


THEOREM. The brick polytope $\Omega(\mathcal{N})$ is the Minkowski sum of the polytopes $\Omega(\mathcal{N}, b)$ associated to the bricks of $\mathcal{N}$ :

$$
\Omega(\mathcal{N})=\sum_{b \text { brick of } \mathcal{N}} \Omega(\mathcal{N}, b)
$$

## BRICK POLYTOPES AND GENERALIZED PERMUTOHEDRA

Generalized permutohedra $=$ polytope whose inequality description is of the form

$$
Z\left(\left\{z_{I}\right\}_{I \in[n]}\right)=\left\{\left.\left(\begin{array}{c}
x_{1} \\
\cdots \\
x_{n}
\end{array}\right) \in \mathbb{R}^{n} \right\rvert\, \sum_{i=1}^{n} x_{i}=z_{[n]} \text { and } \sum_{i \in I} x_{i} \geq z_{I} \text { for } I \subset[n]\right\}
$$

for some family $\left\{z_{I}\right\}_{I \subset[n]} \in \mathbb{R}^{2^{[n]}}$.
A. Postnikov, Permutohedra, associahedra and beyond, 2009.

THEOREM. Any generalized permutahedron is a Minkowski sum of simplices:

$$
Z\left(\left\{z_{I}\right\}_{I \in[n]}\right)=\sum_{I \subset[n]} y_{I} \Delta_{I} \quad \text { where } \quad y_{I}=\sum_{J \subset I}(-1)^{|\backslash J|} z_{J} \quad\left(\text { ie. } \quad z_{I}=\sum_{J \subset I} y_{J}\right) .
$$

F. Ardila, C. Benedetti \& J. Doker, Matroid polytopes and their volumes, 2010.

REMARK. All brick polytopes are generalized permutohedra. Compute $\left\{y_{I}\right\}_{I \subset[n]}$. Which generalized permutohedra are brick polytopes?


## DUPLICATED NETWORKS: PERMUTAHEDRA

Reduced network $=$ network with $n$ levels and $\binom{n}{2}$ commutators. It supports only one pseudoline arrangement.

Duplicated network $\Pi=$ network with $n$ levels and $2\binom{n}{2}$ commutators obtained by duplicating each commutator of a reduced network.


Any pseudoline arrangement supported by $\Pi$ has one contact and one crossing among each pair of duplicated commutators.

## DUPLICATED NETWORKS: PERMUTAHEDRA

Graph of flips $G(\Pi)=\binom{n}{2}$-dimensional cube.


## DUPLICATED NETWORKS: PERMUTAHEDRA



Any pseudoline arrangement supported by $\Pi$ has one contact and one crossing among each pair of duplicated commutators. $\Longrightarrow$ The contact graph $\Lambda^{\#}$ is a tournament.

Vertices of $\Omega(\Pi) \Longleftrightarrow$ acyclic tournaments $\Longleftrightarrow \quad$ permutations of $[n]$
Facets of $\Omega(\Pi) \Longleftrightarrow$ cuts in a tournament $\Longleftrightarrow$ ordered bipartitions of $[n]$

Brick polytope $\Omega(\Pi)=$ permutahedron

## DUPLICATED NETWORKS: PERMUTAHEDRA

Brick polytope $\Omega(\Pi)=$ permutahedron


## DUPLICATED NETWORKS: PERMUTAHEDRA



Minkowski sum decomposition

$$
\Omega(\Pi)=\sum_{b \text { brick of } \Pi} \Omega(\Pi, b)=\sum_{i<j} \text { segment }\left[e_{i}-e_{j}\right]+\sum \text { vertices }=\text { permutahedron }
$$

$$
\begin{aligned}
P(0,1, \ldots, n-1) & =\operatorname{Newton}\left(\operatorname{det}\left[t_{i}^{j-1}\right]_{i, j \in[n]}\right)=\operatorname{Newton}\left(\prod_{1 \leq i<j \leq n}\left(t_{j}-t_{i}\right)\right) \\
& =\sum_{1 \leq i<j \leq n} \operatorname{Newton}\left(t_{j}-t_{i}\right)=\sum_{1 \leq i<j \leq n}\left[e_{j}-e_{i}\right]
\end{aligned}
$$

## ALTERNATING NETWORK: ASSOCIAHEDRA

For $x \in\{a, b\}^{n-2}$, we define a reduced alternating network $\mathcal{N}_{x}$ and a polygon $\mathcal{P}_{x}$.

$\mathcal{N}_{x}$ is the dual pseudoline arrangement of the polygon $\mathcal{P}_{x}$.

## ALTERNATING NETWORK: ASSOCIAHEDRA

THEOREM. There is a duality between the pseudoline arrangements supported by $\mathcal{N}_{x}^{1}$ and the triangulations of the polygon $\mathcal{P}_{x}$.

$T$ triangulation of $\mathcal{P}_{x} \longleftrightarrow T^{*}$ pseudoline arrangement supported by $\mathcal{N}_{x}^{1}$ $\Delta$ triangle of $T \longleftrightarrow \Delta^{*}$ pseudoline of $T^{*}$
$e$ common edge of $\Delta$ and $\Delta^{\prime} \longleftrightarrow e^{*}$ contact between $\Delta^{*}$ and $\Delta^{*}$ $f$ common bissector of $\Delta$ and $\Delta^{\prime} \longleftrightarrow f^{*}$ crossing between $\Delta^{*}$ and $\Delta^{*}$

COROLLARY. (i) The graph of flips $G\left(\mathcal{N}_{x}^{1}\right)$ is (isomorphic to) the graph of flips $G\left(\mathcal{P}_{x}\right)$.
(ii) The contact graph $\left(T^{*}\right)^{\#}$ is (isomorphic to) the dual binary tree of $T$.

THEOREM. For any word $x \in\{a, b\}^{n-2}$, the simplicial complex of crossing-free sets of internal diagonals of the convex $n$-gon $\mathcal{P}_{x}$ is (isomorphic to) the boundary complex of the polar of the brick polytope $\Omega\left(\mathcal{N}_{x}^{1}\right)$.


REMARK. Up to translation, we obtain Hohlweg \& Lange's associahedra.
C. Hohlweg \& C. Lange, Realizations of the associahedron and cyclohedron, 2007.
arXiv:1103.2731


