Efficient Unfolding of Coloured Petri Nets using Interval Decision Diagrams

Martin Schwarick, Christian Rohr, Fei Liu, George Assaf, Jacek Chodak and Monika Heiner

Brandenburg Technical University Petri Nets 2020 - Paris

25 June 2020

Schwarick, Rohr, Liu, Assaf, Chodak, Heiner

25 June 2020 1 / 37

Outline

- 1 Why coloured Petri Nets?
- 2 State of the Art
 - 3 The Problem
 - What are IDD?
- IDD basic principles
- IDD-based unfolding algorithm
 - D Experiments



Powerful modelling - Coloured Petri nets







\\ Simple CS
enum Nodes = {A,B,C,D,F,G,H,I,J };
/\ Product CS
Matrix = Prod(Nodes,Nodes);
/\ Subset CS
Connections = Matrix[(a=A &
 (b=B|b=C |b=D | b=G | b=1)) |
 (a=B & (b=A|b=G|b=H|b=J)) |
 (a=C & (b=F | b=H))];
variables:
Nodes : a:

Nodes : a; Nodes : b; functions: bool IsConnected(Node

Image: Image:

bool IsConnected(Node a1, Node b1)

{(a1,b1) elemOf Connections};



\\ Simple CS
enum Nodes = {A,B,C,D,F,G,H,I,J };
/\ Product CS
Matrix = Prod(Nodes,Nodes);
/\ Subset CS
Connections = Matrix[(a=A &
(b=B|b=C |b=D | b=G | b=I)) |
(a=B & (b=A|b=G|b=H|b=J)) |
(a=C & (b=F | b=H))];
variables:

Nodes : a; Nodes : b; functions: bool IsConnected(Node a1, Node b1) {(a1,b1) elemOf Connections};



\\ Simple CS
enum Nodes = {A,B,C,D,F,G,H,I,J };
\\ Product CS
Matrix = Prod(Nodes,Nodes);
\\ Subset CS
Connections = Matrix[(a=A &
(b=B|b=C |b=D | b=G | b=I)) |
(a=B & (b=A|b=G|b=H|b=J)) |
(a=C & (b=F | b=H))];
variables:

Nodes : a; Nodes : b; functions: bool IsConnected(Node a1, Node b1)

{(a1,b1) elemOf Connections};



\\ Simple CS
enum Nodes = {A,B,C,D,F,G,H,I,J };
\\ Product CS
Matrix = Prod(Nodes,Nodes);
\\ Subset CS
Connections = Matrix[(a=A &
 (b=B|b=C |b=D | b=G | b=I)) |
 (a=B & (b=A|b=G|b=H|b=J)) |
 (a=C & (b=F | b=H))];
variables:
Nodes : a;
Nodes : a;
Nodes : b;
functions:

bool IsConnected(Node a1, Node b1)

{(a1,b1) elemOf Connections};

Image: A matrix



\\ Simple CS
enum Nodes = {A,B,C,D,F,G,H,I,J };
\\ Product CS
Matrix = Prod(Nodes,Nodes);
\\ Subset CS
Connections = Matrix[(a=A &
(b=B|b=C |b=D | b=G | b=I)) |
(a=B & (b=A|b=G|b=H|b=J)) |
(a=C & (b=F | b=H))];
variables:
Nodes : a;
Nodes : b;
functions:

bool IsConnected(Node a1, Node b1)

{(a1,b1) elemOf Connections};

Image: Image:

Powerful modelling - Coloured Petri nets



coloured stochastic Petri net

Image: A matrix A

• Coloured Petri nets are in use for a wide range of applications, covering natural/engineering/life sciences.

- Coloured Petri nets are in use for a wide range of applications, covering natural/engineering/life sciences.
- Currently, most analysis and simulation techniques require unfolding: coloured Petri net → plain Petri net.

- Coloured Petri nets are in use for a wide range of applications, covering natural/engineering/life sciences.
- Currently, most analysis and simulation techniques require unfolding: coloured Petri net → plain Petri net.
- Unfolding tends to be time consuming.

State of the Art

• Example of a scaleable model to adjust grid size.

State of the Art

- Example of a scaleable model to adjust grid size.
- 2D Diffusion in space.



The Problem

• The core problem of efficient unfolding is to determine the transition instances, e.g, all bindings of the involved variables.



The Problem

• The core problem of efficient unfolding is to determine the transition instances, e.g, all bindings of the involved variables.



The Problem

• The core problem of efficient unfolding is to determine the transition instances, e.g, all bindings of the involved variables.



Interval decision diagrams \rightarrow CSP.



• Directed acyclic graphs (DAGs) to encode interval logic functions in the form of symbolic data structure.



• They have two types of nodes: non-terminal (ellipses) and terminal ones (boxes).





• They have two types of nodes: non-terminal (ellipses) and terminal ones (boxes).







• Non-terminal nodes may have an arbitrary number of outgoing arcs labelled with intervals of natural numbers in the form [a,b).



• The set of all paths going from: the root → the terminal node 1 describes all solutions of the given constraint problem.

- The set of all paths going from: the root → the terminal node 1 describes all solutions of the given constraint problem.
- Typically, one path encodes more than one solution.

- The set of all paths going from: the root → the terminal node 1 describes all solutions of the given constraint problem.
- Typically, one path encodes more than one solution.
- Thus, we can easily pick all CSP solutions from the constraint IDD.

Example (x1 \geq 8) \lor (x1 \in [6, 8) \land x2>0)

• Variable ordering







Schwarick, Rohr, Liu, Assaf, Chodak, Heiner

Example (x1 \geq 8) \vee (x1 \in [6, 8) \wedge x2>0)

• $x1 \ge 8$.



Example (x1 \geq 8) \lor (x1 \in [6, 8) \land x2>0)

• some intermediate screenshots.



25 June 2020 14 / 37

Example (x1 \geq 8) \lor (x1 \in [6, 8) \land x2>0)

• One final solution.



Reducing IDD

 Interval partitions labelling the outgoing arcs of each non-terminal node are reduced. For example, [6,7) and [7,8]

Reducing IDD

- Interval partitions labelling the outgoing arcs of each non-terminal node are reduced. For example, [6,7) and [7,8]
- Each non-terminal node has at least two different children.

Reducing IDD

- Interval partitions labelling the outgoing arcs of each non-terminal node are reduced. For example, [6,7) and [7,8]
- Each non-terminal node has at least two different children.
- There exist no two nodes with isomorphic subgraphs.



not Reduced



Reduced

• Preparation step: registration of constants, color sets, variables and functions.

- Preparation step: registration of constants, color sets, variables and functions.
- Unfold places and determine initial marking: this may involve CSP if we have subset mechanism and to determine initial marking.

- Preparation step: registration of constants, color sets, variables and functions.
- Unfold places and determine initial marking: this may involve CSP if we have subset mechanism and to determine initial marking.
- Unfold transitions: and their adjacent arcs and add the result to the unfolded net.

- Preparation step: registration of constants, color sets, variables and functions.
- Unfold places and determine initial marking: this may involve CSP if we have subset mechanism and to determine initial marking.
- Unfold transitions: and their adjacent arcs and add the result to the unfolded net.
- Add all unfolded places which are involved in the transition unfoldings to the unfolded net. This implicitly removes isolated places.

- Preparation step: registration of constants, color sets, variables and functions.
- Unfold places and determine initial marking: this may involve CSP if we have subset mechanism and to determine initial marking.
- Unfold transitions: and their adjacent arcs and add the result to the unfolded net.
- Add all unfolded places which are involved in the transition unfoldings to the unfolded net. This implicitly removes isolated places.
- The entire pseudo code of the algorithm is given in the paper.



 $cs = \{1,8,3..6,10,9,11,20..23\};$ enum $ab = \{A,C,D\};$ variables: cs : x;

< ロ > < 同 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < 回 > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ >

3

Example (no constraints)



Color definitions:

 $cs = \{1,8,3..6,10,9,11,20..23\};$ enum $ab = \{A,C,D\};$ variables: cs : x;

47 ▶

Schwarick, Rohr, Liu, Assaf, Chodak, Heiner

25 June 2020 18 / 37

Example (no constraints)



3

イロト イヨト イヨト イヨト

Example (with Guard)



Color definitions:

 $cs = \{1,8,3..6,10,9,11,20..23\};$ enum ab = {A,C,D}; variables: cs : x;

Schwarick, Rohr, Liu, Assaf, Chodak, Heiner

25 June 2020 19 / 37

Encoding the entire $cs = \{1, 8, 3, ..6, 10, 9, 11, 20, ..23\}$



B → B

< 4 ₽ >

Constraining the color set cs to $6 \le x$



ъ.

Image: Image:

Constraining the color set cs to $x \leq 10$



3)) J

Image: A matrix

Combining $6 \leq x$ and $x \leq 10$ using & operator



Encoding the entire color set $ab = \{A, C, D\}$



Schwarick, Rohr, Liu, Assaf, Chodak, Heiner

25 June 2020 24 / 37

B → B

Image: Image:

Constraining the color set ab to y = A



Schwarick, Rohr, Liu, Assaf, Chodak, Heiner

25 June 2020 25 / 37

∃ ⊳

< □ > < ^[] >

Merging the result of ($6\leqslant x$ and $x\leqslant 10)$ and (y=A) using & operator



Two-path solution:

1st path :
$$(y = A, x = 6)$$

2nd path : $(y = A, x = 8 ..10)$

Unfolded Net



25 June 2020 27 / 37

글 > 글

• • • • • • • • •

Unfolded Net



글 > 글

• • • • • • • • •

• We compared our IDD unfolding with an unfolding employing the popular constraint solver library Gecode.

- We compared our IDD unfolding with an unfolding employing the popular constraint solver library Gecode.
- 22 MCC models (PNML format) \rightarrow https://mcc.lip6.fr/models.php
 - 1st group: requires no substantial unfolding time.
 - 2nd group: requires substantial unfolding time.

- We compared our IDD unfolding with an unfolding employing the popular constraint solver library Gecode.
- 22 MCC models (PNML format) \rightarrow https://mcc.lip6.fr/models.php
 - 1st group: requires no substantial unfolding time.
 - 2nd group: requires substantial unfolding time.
- We used also two biological test cases from our own collection: https://www-dssz.informatik.tu-cottbus.de/DSSZ/Software/Examples?dir=IddUnfolding
 - 3D Diffusion.
 - Brusselator.

the model

- a lane bridge with limited capacity.
- used by two types of vehicles.
- coloured model has 15 places, 11 transition and 57 arcs.



Bridges and Vehicles (MCC)



Figure: Bridges and vehicles (MCC); requires no substantial unfolding time

the model

- reunification process.
- the coloured model has 104 places, 66 transition and198 arcs.
- it is scaled by the number of legal residents.



Family reunion (MCC)



Figure: Family Reunion (MCC); requires substantial unfolding time

Diffusion in space (3D)



Schwarick, Rohr, Liu, Assaf, Chodak, Heiner

2

Diffusion in space (3D)



Figure: Diffusion (3D); N - Grid size

- ∢ ศ⊒ ▶



Schwarick, Rohr, Liu, Assaf, Chodak, Heiner

25 June 2020 32 / 37

3

< ∃⇒

Image: A math a math



Figure: Brusselator ; N – Grid size of a 2D square



3

・ロト ・四ト ・ヨト ・ヨト

- Gecode, when:
 - models with a few guards or no guards.
 - models with simple colour sets.
 - e.g, 12 MCC models (no substantial time).



- Gecode, when:
 - models with a few guards or no guards.
 - models with simple colour sets.
 - e.g, 12 MCC models (no substantial time).
- IDDs, when:
 - parametrized models with a large scaling factor, e.g, Diffusion in space.
 - models with sophisticated guards.
 - e.g, most models in our own collection.



- Gecode, when:
 - models with a few guards or no guards.
 - models with simple colour sets.
 - e.g, 12 MCC models (no substantial time).
- IDDs, when:
 - parametrized models with a large scaling factor, e.g, Diffusion in space.
 - models with sophisticated guards.
 - e.g, most models in our own collection.
- The complete performance report is available:

https://www-dssz.informatik.tu-

cottbus.de/DSSZ/Software/Examples?dir=IddUnfolding



SNOOPY

MARCIE SPIKE

https://www-dssz.informatik.tu-cottbus.de/DSSZ/Software/

Schwarick, Rohr, Liu, Assaf, Chodak, Heiner

25 June 2020 35 / 37



https://www-dssz.informatik.tu-cottbus.de/DSSZ/Software/

25 June 2020 35 / 37



https://www-dssz.informatik.tu-cottbus.de/DSSZ/Software/

- Performance:
 - Considering memory.
 - Power consumption.

- Performance:
 - Considering memory.
 - Power consumption.
- Implementation efficiency:
 - Multi-threading: unfolding the coloured places and transition is currently done sequentially.
 - Reuse of already computed solutions.
 - Choosing among several variable order strategies.

Thank You For Your Attention