

# PROPORTIONAL LUMPABILITY

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# STOCHASTIC SYSTEMS MODELING

## CONTEXT - CONTINUOUS TIME MARKOV CHAINS

- ▶ **Continuous Time Markov Chains** are the underlying semantics of many high-level formalisms for **modeling, analysing and verifying stochastic systems**, such as Stochastic Petri nets, Stochastic Automata Networks, Markovian process algebras
- ▶ High-level languages **simplify the specification task** thanks to compositionality and abstraction
- ▶ So, even very compact specifications can generate **very large stochastic systems** that are difficult/impossible to analyse

# STATE SPACE REDUCTION

## CONTEXT - LUMPABILITY

- ▶ In the non-deterministic setting **bisimulation** allows to quotient the state space and precisely characterizes **modal logic** [Van Benthem Th.]
- ▶ On Markov Chains **lumpability** [Kemeny-Snell 1976] (probabilistic bisimulation [Larsen-Skou 1991]) plays the same role, preserving **stationary quantities** [Buchholz 1994] and **stochastic/probabilistic modal logics** [Larsen-Skou 1991, Desharnais et al 2002, Bernardo et al. 2019]

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## ISSUE

**Lumpability is too demanding**

As a consequence it usually provides **poor reductions**

# APPROXIMATIONS

## CONTEXT - PSEUDO-METRICS ON PATHS

- ▶ **Distances** measuring the difference between states of probabilistic systems are introduced in [Desharnais et al. 1999]
- ▶ The distance evaluates the probabilities along **paths** allowing **discounts**
- ▶ Probabilistic bisimilar states have distance **0**
- ▶ **Behavioural properties** have been largely investigated [van Breugel et al. 2001, Wild et al. 2019]
- ▶ **Compositionality properties** have been proved [Gebler et al. 2015]
- ▶ **Algorithmic solutions** have been proposed [Bacci et al. Concur 2019]
- ▶ Stationary distribution bounds?

# APPROXIMATIONS

## CONTEXT - QUASI LUMPABILITY AND $\epsilon$ -BISIMULATION

- ▶ **Quasi Lumpability** relates states allowing  $\epsilon$  **perturbations** of the outgoing probabilities/rates [Franceschinis et al. 1994]
- ▶ **Bounds on the stationary distributions** have been proved
- ▶ **Behavioural properties** have been studied on  $\epsilon$ -Bisimulation [Desharnais et al. 2008, Tracol et al. 2011, Abate et al. 2014, Abate et al. 2017]
- ▶ **Algorithmic solutions** have been proposed [Milios et al. 2012]

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- ▶ **Algorithmic solutions** have been proposed [Milios et al. 2012]

## UNFORTUNATELY

It is **not possible to exactly reconstruct the stationary distribution** of the original system

# PROPORTIONAL LUMPABILITY

## MOTIVATION

We aim at **relaxing** the conditions of **lumpability** while **allowing to derive the exact stationary indices for the original system**



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## CONTRIBUTION

- ▶ We define the notion of **Proportional Lumpability** over **Continuous Time Markov Chains (CTMC)**
- ▶ We show that this allows to **exactly derive the original stationary distribution**
- ▶ We introduce the notion of **Proportional Bisimulation** over the stochastic process algebra **PEPA** and prove that it induces a proportional lumpability on the underlying semantics

# OUTLINE OF THE TALK

- ▶ The notions of **Lumpability** and **Quasi Lumpability** over CTMC
- ▶ The notion of **Proportional Lumpability** and its **properties**
- ▶ Proportional Lumpability over the **Process Algebra PEPA**
- ▶ Example
- ▶ Conclusions

# CONTINUOUS TIME MARKOV CHAINS

## CTMC

Let  $X(t)$  with  $t \in \mathbb{R}^+$  be a **stochastic process** taking values in a discrete space  $\mathcal{S}$ .

$X(t)$  is a CTMC if it is **stationary** and **markovian**

We focus on **finite**, **time-homogeneous**, **ergodic** Markov Chains

## INFINITESIMAL GENERATOR

A CTMC is given as a **matrix**  $Q$  of dim.  $|\mathcal{S}| \times |\mathcal{S}|$  such that:

- ▶ for  $i \neq j$  the **transition rate** from  $i$  to  $j$  is  $q(i, j) \geq 0$ , i.e.,

$$\text{Prob}(X(t+h) = j | X(t) = i) = q(i, j) * h + o(h)$$

- ▶  $q(i, i) = -\sum_{j \neq i} q(i, j)$

# STATIONARY ANALYSIS

## STATIONARY DISTRIBUTION

A distribution  $\pi$  over  $\mathcal{S}$  such that  $\pi(i)$  is the probability of being in  $i$  when time goes to  $\infty$

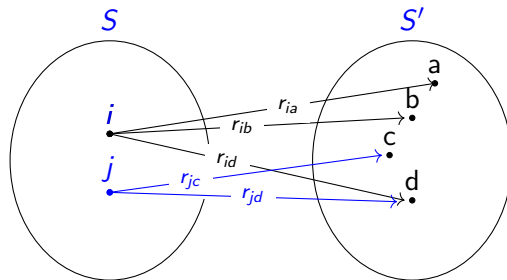
In our setting  $\pi$  is the unique distribution that solves

$$\pi Q = 0$$

## STATIONARY PERFORMANCES INDICES

Stationary performances indices, such as throughput, expected response time, resource utilization, can be computed from the steady state distribution  $\pi$

## LUMPABILITY - INTUITIVELY



$$r_{ia} + r_{ib} + r_{id} = r_{jc} + r_{jd}$$

# LUMPABILITY

## STRONG LUMPABILITY

The strong lumpability  $\sim$  is the largest equivalence over  $\mathcal{S}$  such that  $\forall S, S' \in \mathcal{S}/\sim$  and  $\forall i, j \in S$

$$\sum_{a \in S'} q(i, a) = \sum_{a \in S'} q(j, a)$$

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## PROPERTIES

- ▶ We can safely restrict to  $S \neq S'$
- ▶ There always exists a unique maximum lumpability
- ▶ The stationary distribution  $\Pi$  of the lumped chain is the aggregation of  $\pi$
- ▶ Probabilistic modal logic properties are preserved

# QUASI LUMPABILITY

QUASI LUMPABILITY [FRANCESCHINIS ET AL. '94, MILIOS ET AL. 2012]

An  $\epsilon$ -quasi lumpability  $\mathcal{R}$  is an equivalence over  $\mathcal{S}$  such that  $\forall S, S' \in \mathcal{S}/\mathcal{R}$  and  $\forall i, j \in S$

$$|\sum_{a \in S'} q(i, a) - \sum_{a \in S'} q(j, a)| \leq \epsilon$$



## QUASI LUMPABILITY

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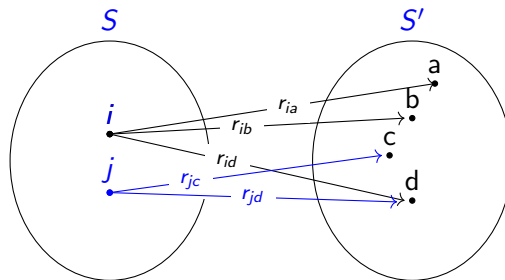
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$$|\sum_{a \in S'} q(i, a) - \sum_{a \in S'} q(j, a)| \leq \epsilon$$

## PROPERTIES

- ▶ It was originally defined splitting  $Q$  into  $Q^-$  and  $Q^\epsilon$  (perturbation)
- ▶ Bounds on the exact stationary distribution (indices) can be computed
- ▶ Algorithms for approximating an optimal aggregation have been proposed

## QUASI LUMPABILITY – EXAMPLE



$$r_{ia} + r_{ib} + r_{id} = 10 \quad r_{jc} + r_{jd} = 100$$

$$\epsilon \geq 90$$

# PROPORTIONAL LUMPABILITY

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Given  $\kappa : \mathcal{S} \rightarrow \mathbb{R}^+$ , a  $\kappa$ -proportional lumpability  $\mathcal{R}$  is an equivalence over  $\mathcal{S}$  such that  $\forall S, S' \in \mathcal{S}/\mathcal{R}$  and  $\forall i, j \in S$

$$\frac{\sum_{a \in S'} q(i, a)}{\kappa(i)} = \frac{\sum_{a \in S'} q(j, a)}{\kappa(j)}$$

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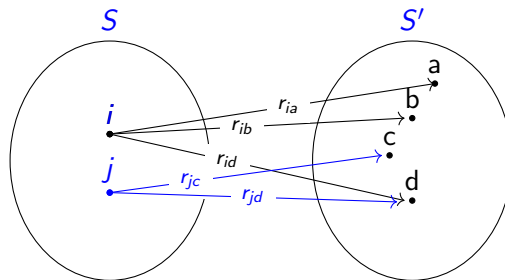
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## PROPERTIES

- ▶ We can safely restrict to  $S \neq S'$
- ▶ There exists a unique maximum  $\kappa$ -proportional lumpability  $\sim_\kappa$
- ▶ More properties ... thanks to one of FORMATS reviewers

# PROPORTIONAL LUMPABILITY – EXAMPLE



$$r_{ia} + r_{ib} + r_{id} = 10 \quad r_{jc} + r_{jd} = 100$$

$$\kappa(i) = 1 \quad \kappa(j) = 10$$

# PERTURBED SYSTEMS

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It is any CTMC  $X'(t)$  over the state space  $\mathcal{S}$  having generator  $Q'$  such that  $\forall i \in \mathcal{S}$  and  $\forall S' \in \mathcal{S}/\sim$

$$\sum_{a \in S', a \neq i} q'(i, a) = \frac{\sum_{a \in S', a \neq i} q(i, a)}{\kappa(i)}$$

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## EXAMPLE

$X'(t)$  defined by

$$q'(i, a) = \frac{q(i, a)}{\kappa(i)} \quad \text{for any } a \neq i$$



# STATIONARY DISTRIBUTIONS OF PERTURBED SYSTEMS

## PROPOSITION

The stationary distributions of  $X(t)$  and  $X'(t)$  are related as follows

$$\pi(i) = \frac{\pi'(i)}{K\kappa(i)}$$

where the normalization factor is  $K = \sum_{i \in \mathcal{S}} \pi'(i)/\kappa(i)$

# AGGREGATED SYSTEM

## AGGREGATED SYSTEM

It is the CTMC  $\tilde{X}(t)$

- ▶ over the state space  $\mathcal{S}/\sim$
- ▶ it has infinitesimal generator  $\tilde{Q}$  with  $\tilde{q}(S, S') = \frac{\sum_{a \in S'} q(i, a)}{\kappa(i)}$  with  $i \in S$

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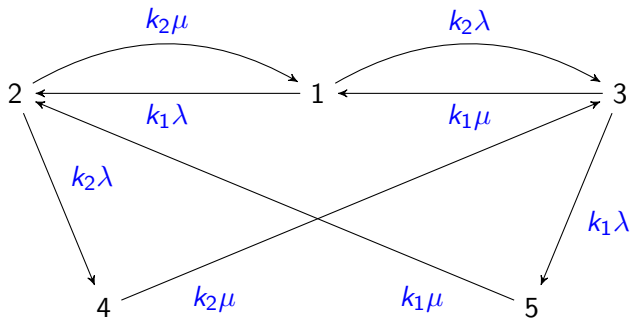
## PROPOSITION

The stationary distributions of  $X(t)$  and  $\tilde{X}(t)$  are related as follows

$$\tilde{\pi}(S) = \frac{\sum_{i \in S} \pi(i) \kappa(i)}{\tilde{K}}$$

where the normalization factor is  $\tilde{K} = \sum_{i \in \mathcal{S}} \pi(i) \kappa(i)$

## EXAMPLE - CPUs SYSTEM



$$\kappa(1) = 1 \quad \kappa(2) = k_2 \quad \kappa(3) = k_1 \quad \kappa(4) = k_2 \quad \kappa(5) = k_1$$

# PERFORMANCES EVALUATION PROCESS ALGEBRA - PEPA

## PEPA SYNTAX

Let  $\mathcal{A}$  be a set of actions with  $\tau \in \mathcal{A}$

Let  $\alpha \in \mathcal{A}$ ,  $A \subseteq \mathcal{A}$ , and  $r \in \mathbb{R}$

$$S ::= \mathbf{0} \mid (\alpha, r).S \mid S + S \mid X$$

$$P ::= P \boxtimes_A P \mid P/A \mid P \setminus A \mid S$$

Each variable  $X$  is associated to a definition  $X \stackrel{\text{def}}{=} P$

## PEPA SEMANTICS

It defines Labeled Continuous Time Markov Chains

# PERFORMANCES EVALUATION PROCESS ALGEBRA - PEPA

$$\frac{P \xrightarrow{(\alpha, r)} P'}{P \boxtimes_A Q \xrightarrow{(\alpha, r)} P' \boxtimes_A Q} \quad (\alpha \notin A) \qquad \frac{Q \xrightarrow{(\alpha, r)} Q'}{P \boxtimes_A Q \xrightarrow{(\alpha, r)} P \boxtimes_A Q'} \quad (\alpha \notin A)$$

$$\frac{P \xrightarrow{(\alpha, r_1)} P' \quad Q \xrightarrow{(\alpha, r_2)} Q'}{P \boxtimes_A Q \xrightarrow{(\alpha, R)} P' \boxtimes_A Q'} \quad (\alpha \in A)$$

where  $R = \frac{r_1}{r_\alpha(P)} \frac{r_2}{r_\alpha(Q)} \min(r_\alpha(P), r_\alpha(Q))$

# LUMPABLE BISIMILARITY

LUMPABLE BISIMILARITY [HILLSTON ET AL. 2013, ALZETTA ET AL. 2018]

A lumpable bisimilarity is an equivalence  $\mathcal{R}$  such that for each action  $\alpha$ ,  $\forall S, S' \in \mathcal{C}/\mathcal{R}$ , and  $\forall P, Q \in S$

- ▶ either  $\alpha \neq \tau$ ,
- ▶ or  $\alpha = \tau$  and  $S \neq S'$ ,

it holds

$$\sum_{P' \in S', P \xrightarrow{(\alpha, r_\alpha)} P'} r_\alpha = \sum_{Q' \in S', Q \xrightarrow{(\alpha, r_\alpha)} Q'} r_\alpha$$

## PROPERTIES

There exists a unique maximum lumpable bisimilarity  $\approx_l$ , it is *contextual*, *action preserving*, and induces a *lumpability*

# PROPORTIONAL BISIMILARITY

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Given  $\kappa : \mathcal{C} \rightarrow \mathbb{R}^+$  a  $\kappa$ -proportional bisimilarity is an equivalence  $\mathcal{R}$  such that for each action  $\alpha$ ,  $\forall S, S' \in \mathcal{C}/\mathcal{R}$ , and  $\forall P, Q \in S$

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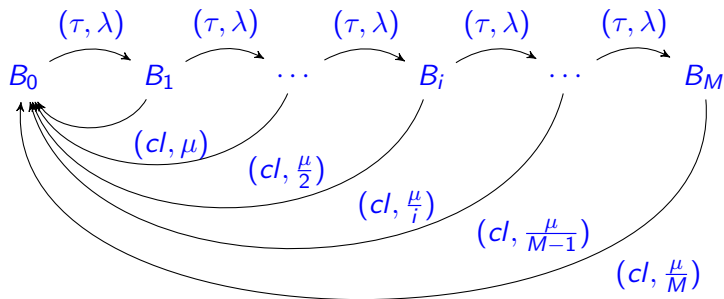
$$\frac{\sum_{P' \in S', P \xrightarrow{(\alpha, r_\alpha)} P'} r_\alpha}{\kappa(P)} = \frac{\sum_{Q' \in S', Q \xrightarrow{(\alpha, r_\alpha)} Q'} r_\alpha}{\kappa(Q)}$$

## PROPERTIES

There exists a unique maximum proportional bisimilarity  $\approx_l^\kappa$ , it induces a *proportional lumpability*



## EXAMPLE



# CONCLUSIONS

- ▶ The notion of **proportional lumpability** has been introduced
- ▶ It “preserves” the **stationary distribution**
- ▶ It can be applied for **PEPA components reduction**
- ▶ We are looking at its **computation** and **compositionality** properties