

# Substitutions on Multidimensional Sequences

Thomas Fernique

LIRMM - Université Montpellier II  
Poncelet Lab. - Independent University of Moscow

WORDS'05  
September 13-17, 2005

## 1 Local rules

- Main idea
- General definition

## 2 Substitutions on patterns

- $\Lambda$ -paths
- Good sets of local rules

## 3 Derived local rules

- Global rules
- Generalized substitutions

## 1 Local rules

- Main idea
- General definition

## 2 Substitutions on patterns

- $\Lambda$ -paths
- Good sets of local rules

## 3 Derived local rules

- Global rules
- Generalized substitutions

# Classic substitutions on words

$$\boxed{1} \longrightarrow \boxed{1} \boxed{2}$$

$$\boxed{2} \longrightarrow \boxed{1}$$

$$\boxed{1} \boxed{2} \boxed{1} \boxed{1} \boxed{2}$$



$$\boxed{1} \boxed{2}$$

# Classic substitutions on words

$$\boxed{1} \longrightarrow \boxed{1} \boxed{2}$$

$$\boxed{2} \longrightarrow \boxed{1}$$

**init**

$$\boxed{1} \boxed{2} \boxed{1} \boxed{1} \boxed{2}$$



$$\boxed{1} \boxed{2}$$

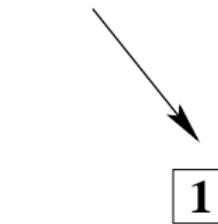
# Classic substitutions on words

$$\boxed{1} \longrightarrow \boxed{1} \boxed{2}$$

$$\boxed{2} \longrightarrow \boxed{1}$$

init

$$\boxed{1} \boxed{2} \boxed{1} \boxed{1} \boxed{2}$$



# Classic substitutions on words

$$\boxed{1} \longrightarrow \boxed{1} \boxed{2}$$

$$\boxed{2} \longrightarrow \boxed{1}$$

init

concat

$$\boxed{1} \boxed{2} \boxed{1} \boxed{1} \boxed{2}$$

$$\boxed{1} \boxed{2} \boxed{1}$$



# Classic substitutions on words

**1** → **1 2**

**1 2 1 1 2**

**2** → **1**

**init**

**concat**

**1 2 1 1 2 1 2 1**

# How to avoid the concatenation

$$\boxed{1} \longrightarrow \boxed{1 \mid 2}$$

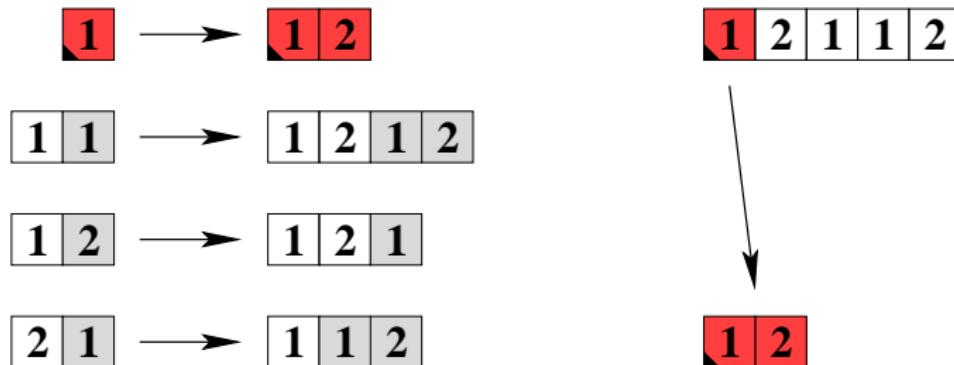
$$\boxed{1 \mid 2 \mid 1 \mid 1 \mid 2}$$

$$\boxed{1 \mid 1} \longrightarrow \boxed{1 \mid 2 \mid 1 \mid 2}$$

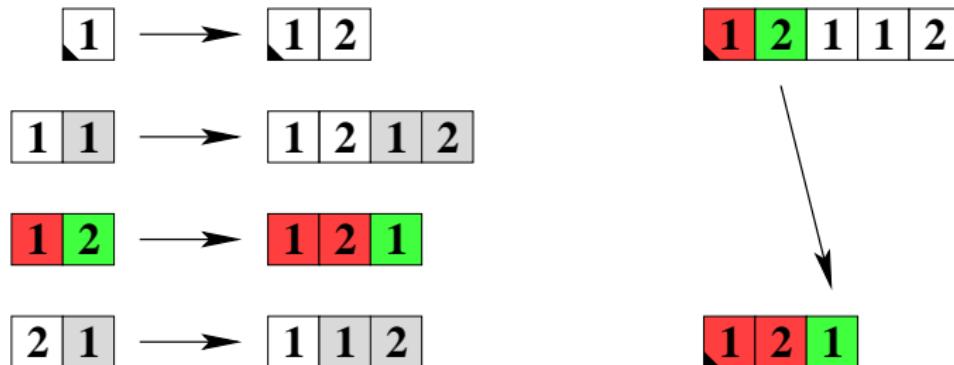
$$\boxed{1 \mid 2} \longrightarrow \boxed{1 \mid 2 \mid 1}$$

$$\boxed{2 \mid 1} \longrightarrow \boxed{1 \mid 1 \mid 2}$$

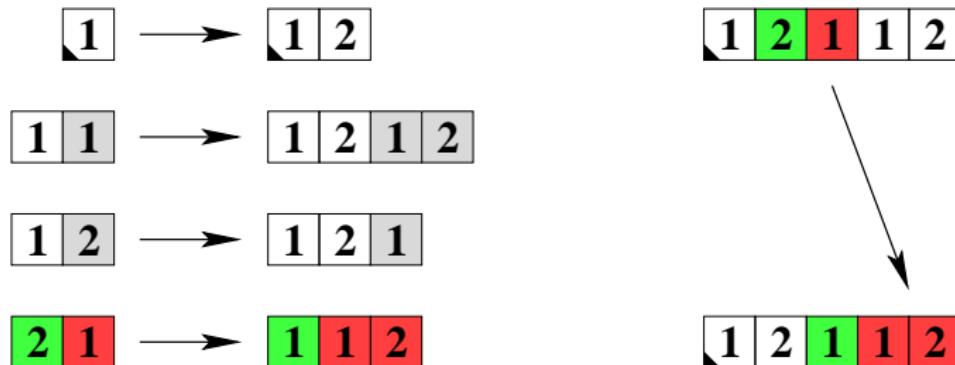
# How to avoid the concatenation



# How to avoid the concatenation



# How to avoid the concatenation



# How to avoid the concatenation

$$\boxed{1} \longrightarrow \boxed{1 \mid 2}$$

$$\boxed{1 \mid 2 \mid 1 \mid \color{red}{1} \mid \color{green}{1} \mid 2}$$

$$\boxed{1 \mid \color{red}{1}} \longrightarrow \boxed{1 \mid 2 \mid \color{red}{1} \mid \color{green}{1}}$$

$$\boxed{1 \mid 2} \longrightarrow \boxed{1 \mid 2 \mid 1}$$

$$\boxed{2 \mid 1} \longrightarrow \boxed{1 \mid 1 \mid 2}$$

$$\boxed{1 \mid 2 \mid 1 \mid \color{red}{1} \mid \color{red}{2} \mid \color{green}{1} \mid 2}$$

# How to avoid the concatenation

$$\boxed{1} \longrightarrow \boxed{1 \mid 2}$$

$$\boxed{1 \mid 1} \longrightarrow \boxed{1 \mid 2 \mid 1 \mid 2}$$

$$\boxed{1 \mid 2} \longrightarrow \boxed{1 \mid 2 \mid 1}$$

$$\boxed{2 \mid 1} \longrightarrow \boxed{1 \mid 1 \mid 2}$$

$$\boxed{1 \mid 2 \mid 1 \mid \color{green}{1} \mid \color{red}{2}}$$

$$\boxed{1 \mid 2 \mid 1 \mid 1 \mid 2 \mid \color{green}{1} \mid \color{green}{2} \mid \color{red}{1}}$$



# How to avoid the concatenation

$$\boxed{1} \longrightarrow \boxed{1 \mid 2}$$

$$\boxed{1 \mid 2 \mid 1 \mid 1 \mid 2}$$

$$\boxed{1 \mid 1} \longrightarrow \boxed{1 \mid 2 \mid 1 \mid 2}$$

$$\boxed{1 \mid 2} \longrightarrow \boxed{1 \mid 2 \mid 1}$$

$$\boxed{2 \mid 1} \longrightarrow \boxed{1 \mid 1 \mid 2}$$

$$\boxed{1 \mid 2 \mid 1 \mid 1 \mid 2 \mid 1 \mid 2 \mid 1}$$

## 1 Local rules

- Main idea
- General definition

## 2 Substitutions on patterns

- $\Lambda$ -paths
- Good sets of local rules

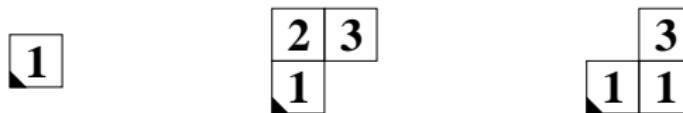
## 3 Derived local rules

- Global rules
- Generalized substitutions

# Basic notions 1/2

*Pointed letter:*  $L = (x, \ell) \in \mathbb{Z}^n \times \mathcal{A}$ ,  $x = \text{Supp}(L)$ .  $\rightsquigarrow$  Set  $\mathcal{L}$ .

*Pointed pattern:* pointed letters with distinct supports.  $\rightsquigarrow$  Set  $\mathcal{P}$ .



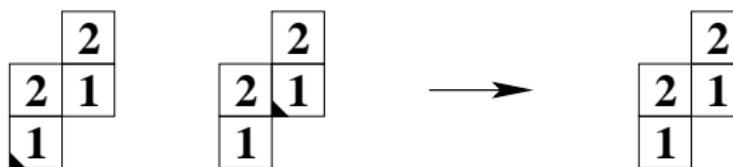
*Consistence* of  $P, Q \in \mathcal{P}$ :  $((x, \ell) \in P \wedge (x, \ell') \in Q) \Rightarrow \ell = \ell'$ .

## Basic notions 2/2

$\mathbb{Z}^n$  acts on  $\mathcal{L}$  and  $\mathcal{P}$  by translation:  $y + (x, \ell) = (y + x, \ell)$ .

↪ Non-pointed letters and patterns:  $\overline{\mathcal{L}} = \mathcal{L}/\mathbb{Z}^n$  and  $\overline{\mathcal{P}} = \mathcal{P}/\mathbb{Z}^n$ .

$P$  and  $Q$  congruent:  $\overline{P} = \overline{Q}$ , i.e.  $P = Q + y$ .

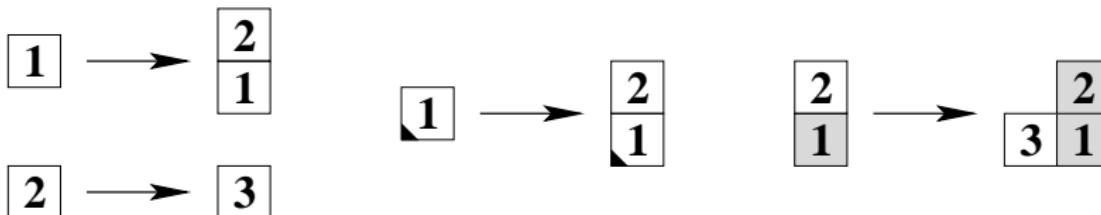


## Definition

A *non-pointed substitution* is a map  $\bar{\sigma} : \overline{\mathcal{L}} \rightarrow \overline{\mathcal{P}}$ .

### Definition (Local rules for $\bar{\sigma}$ )

- ◊ *Initial rule*  $\lambda^*$ : def. on  $I(\lambda^*) = \{L\}$  s.t.  $\overline{\lambda^*(L)} = \bar{\sigma}(\overline{L})$ ;
- ◊ *extension rule*  $\lambda$ : def. on  $E(\lambda) = \{L_1, L_2\}$  s.t.  $\overline{\lambda(L_i)} = \bar{\sigma}(\overline{L_i})$ .



## 1 Local rules

- Main idea
- General definition

## 2 Substitutions on patterns

- $\Lambda$ -paths
- Good sets of local rules

## 3 Derived local rules

- Global rules
- Generalized substitutions

Definition ( $\Lambda$ -path)

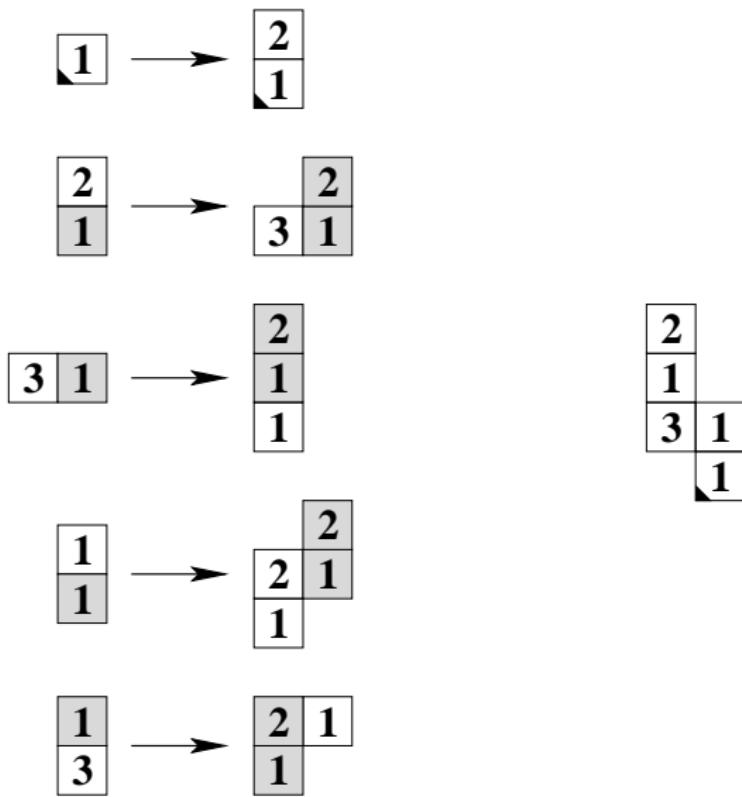
$R = (R_1, \dots, R_k) \in \mathcal{L}^k$  is a  $\Lambda$ -path if there exist:

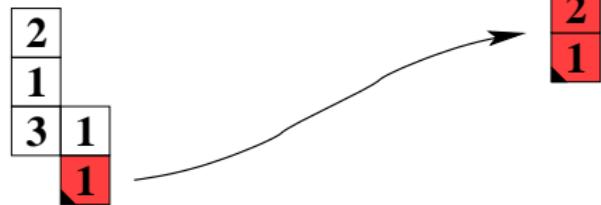
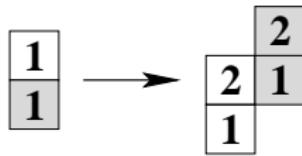
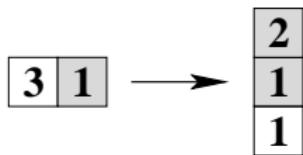
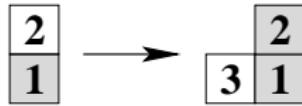
- ◊ an initial rule  $\lambda^* \in \Lambda$  s.t.  $I(\lambda^*) = R_1$ ;
- ◊ for  $1 \leq i < k$ , an extension rule  $\lambda_i \in \Lambda$  and  $x_i \in \mathbb{Z}^n$  s.t.  
 $E(\lambda_i) = \{L_1^i, L_2^i\}$  and  $(R_i, R_{i+1}) = (L_1^i, L_2^i) + x_i$ .

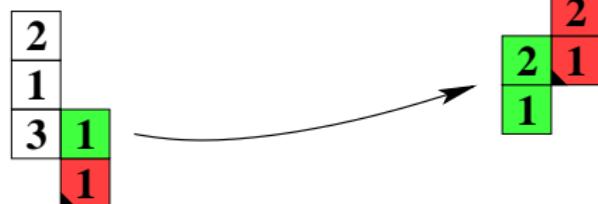
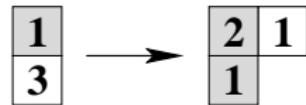
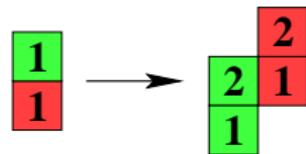
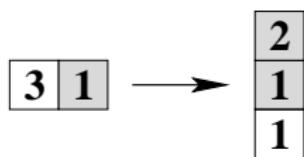
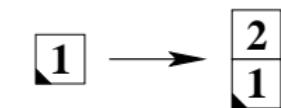
Definition (Substitution on a  $\Lambda$ -path)

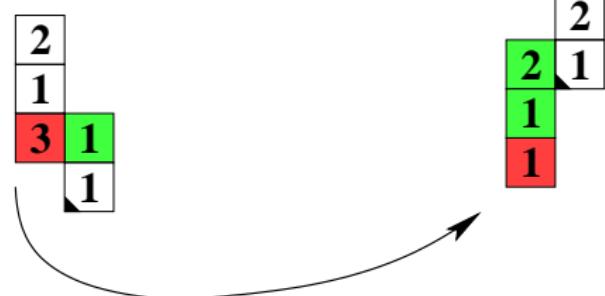
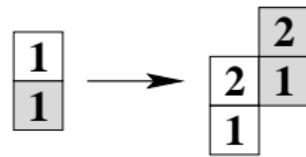
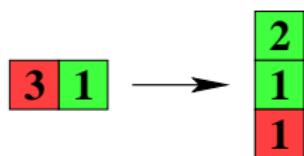
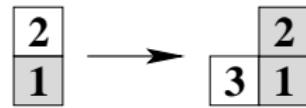
$(\bar{\sigma}, \Lambda, R)$  defined on  $R$  by induction:

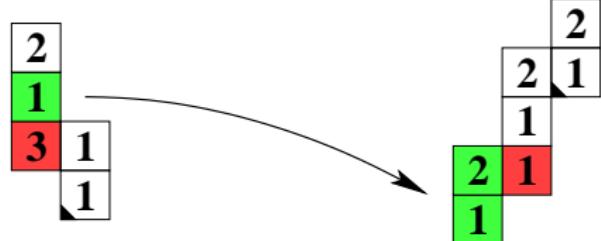
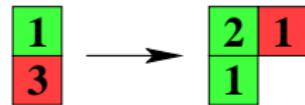
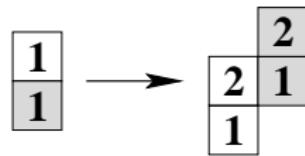
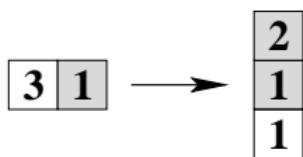
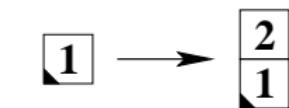
- ◊  $(\bar{\sigma}, \Lambda, R)(R_1) = \lambda^*(R_1)$ ;
- ◊  $(\bar{\sigma}, \Lambda, R)(R_{i+1}) = \lambda_i(L_2^i) + \vec{v}(\lambda_i(L_1^i), (\bar{\sigma}, \Lambda, R)(R_i))$ .

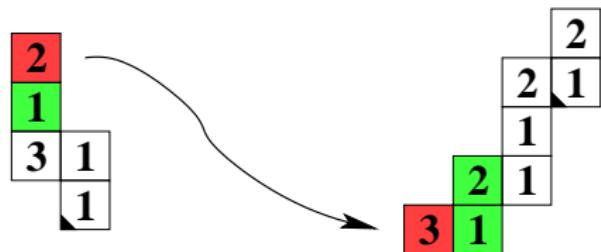
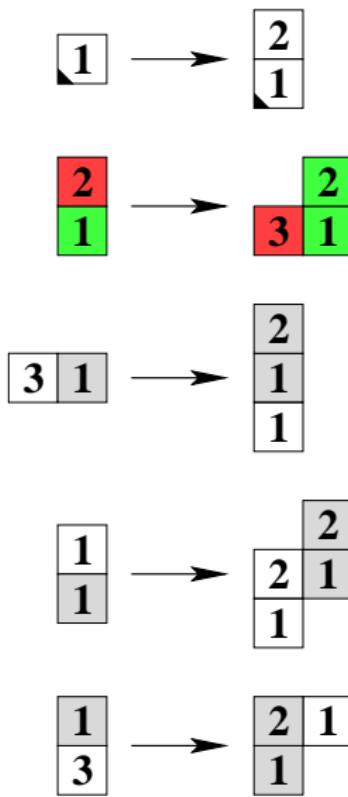


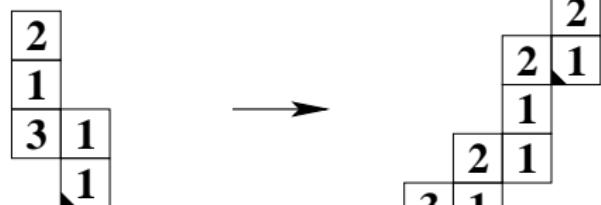
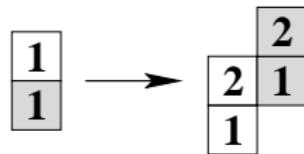
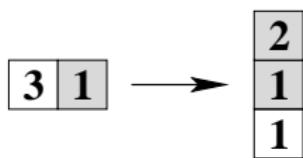
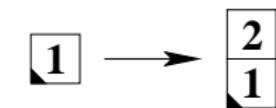












## 1 Local rules

- Main idea
- General definition

## 2 Substitutions on patterns

- $\Lambda$ -paths
- Good sets of local rules

## 3 Derived local rules

- Global rules
- Generalized substitutions

## Definition (Good set of local rules)

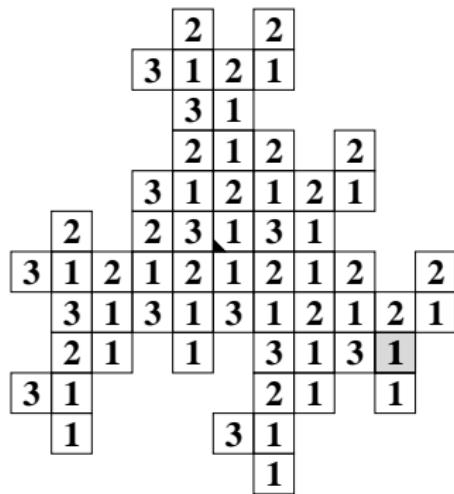
$\Lambda$  is a *good* set of local rules on a pointed pattern  $U$  if:

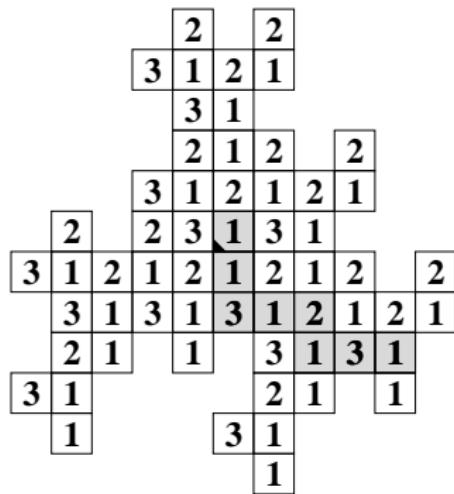
- ◊ any pointed letter of  $U$  belongs to a  $\Lambda$ -path of  $U$ ;
- ◊ the images of any two  $\Lambda$ -paths of  $U$  are consistent.

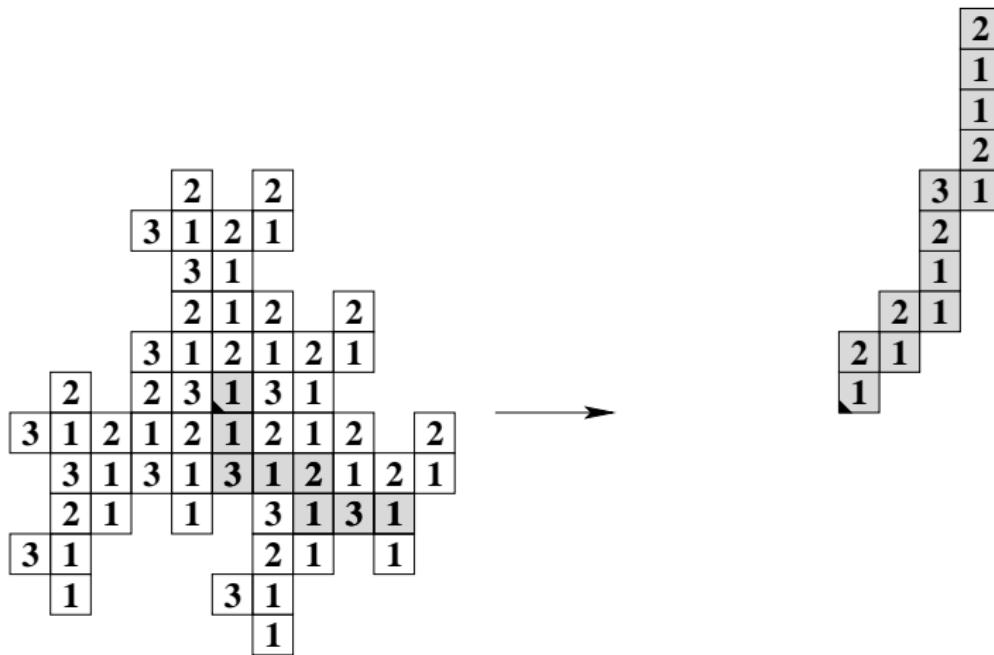
## Definition (Substitution on a pattern)

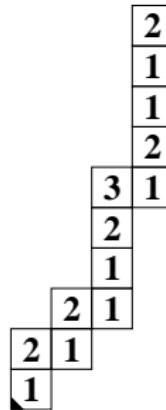
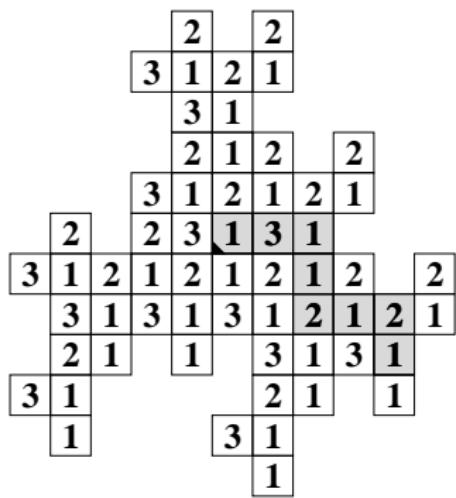
If  $\Lambda$  is a good set of local rules on  $U$ , then one sets:

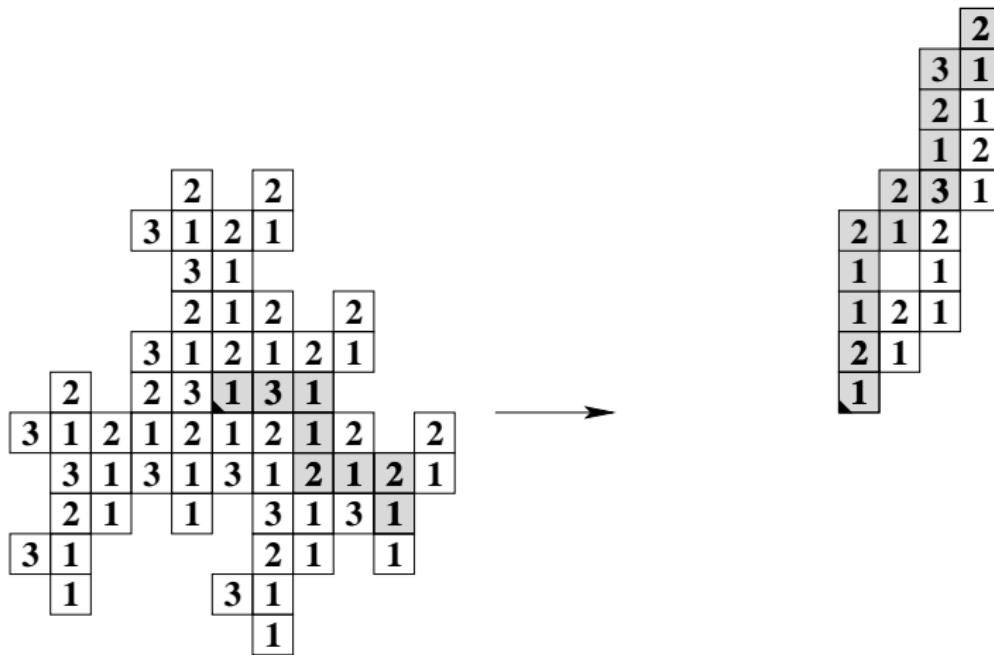
$$(\bar{\sigma}, \Lambda)(U) = \bigcup \{ (\bar{\sigma}, \Lambda, R)(U) \mid R \text{ is a } \Lambda\text{-path of } U \}.$$

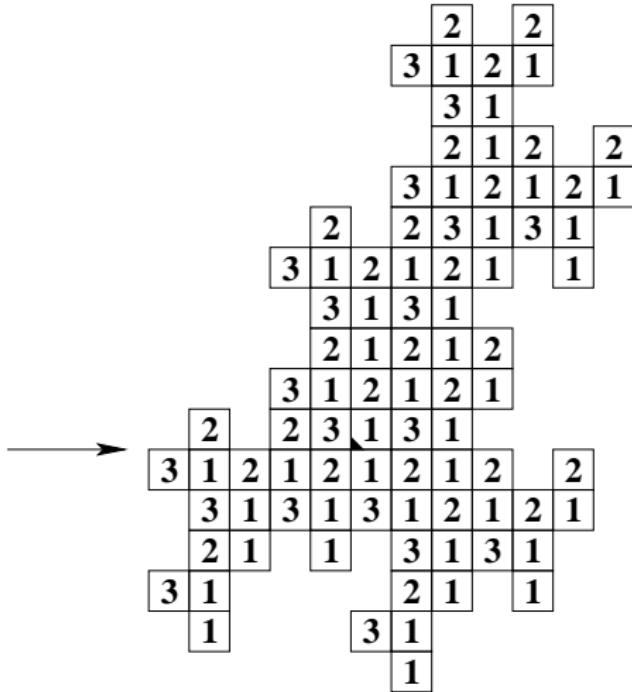
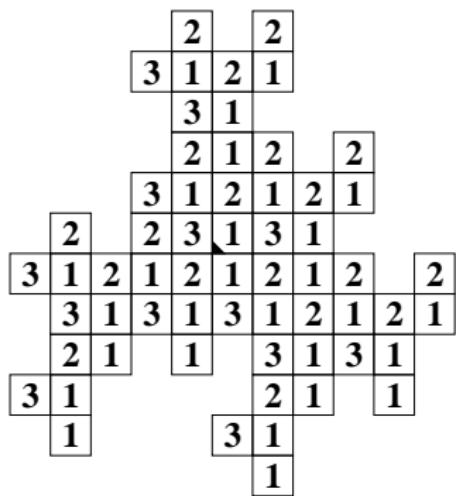












## 1 Local rules

- Main idea
- General definition

## 2 Substitutions on patterns

- $\Lambda$ -paths
- Good sets of local rules

## 3 Derived local rules

- Global rules
- Generalized substitutions

$\bar{\sigma}$  non-pointed substitution,  $\mathcal{H} \subset \mathcal{P}$  and  $\mathcal{L}(\mathcal{H}) = \{L \in U \mid U \in \mathcal{H}\}$ .

### Definition (Context-Free Global Rule)

A *CF global rule* on  $\mathcal{H}$  for  $\bar{\sigma}$  is a map  $\Gamma : \mathcal{L}(\mathcal{H}) \rightarrow \mathcal{P}$  such that:

- ◊  $\overline{\Gamma(L)} = \bar{\sigma}(\overline{L})$ ;
- ◊  $(\overline{U} = \overline{V}) \Rightarrow (\overline{\Gamma(U)} = \overline{\Gamma(V)})$ ;
- ◊  $Supp(L) \neq Supp(L') \Rightarrow Supp(\Gamma(L)) \cap Supp(\Gamma(L')) = \emptyset$ ;
- ◊  $\forall L \in \mathcal{L}(\mathcal{H}), \exists L' \text{ s.t. } L \in \Gamma(L')$ .

Hyp:  $\exists U_0 \in \mathcal{P}$  s.t.  $\forall U \in \mathcal{H}$ ,  $U$  is 4-connected and  $U \cap U_0 \neq \emptyset$ .

### Theorem

$\exists$  CF global rule  $\Gamma \iff \exists$  finite good set  $\Lambda$  s.t.  $(\bar{\sigma}, \Lambda) = \Gamma$  on  $\mathcal{H}$ .

$\Lambda$  computable from  $\Gamma$ .

## 1 Local rules

- Main idea
- General definition

## 2 Substitutions on patterns

- $\Lambda$ -paths
- Good sets of local rules

## 3 Derived local rules

- Global rules
- Generalized substitutions

*Sturmian hyperplane sequence:*  $(n - 1)$ -dimensional sequence over  $\{1, \dots, n\}$  which generalizes Sturmian words.

*Generalized substitution:* acts on such sequences.

## Proposition

A generalized substitution is a CF global rule acting on the Sturmian hyperplane sequences.

↔ Good sets of local rules

# And now?

Defined: substitutions on multidimensional sequences.

◊ Fixed point sequences?

Classic substitutions: characterization by *return words* (Durand)

Local rules in 1D: nothing more (Fogg)

◊ Fixed points among *some* sequences?

Sturmian words: *iff* quadratic slope (Crisp & Moran)

Sturmian hyperplane sequences: sufficient condition (F.)

# And now?

Defined: substitutions on multidimensional sequences.

◊ Fixed point sequences?

Classic substitutions: characterization by *return words* (Durand)

Local rules in 1D: nothing more (Fogg)

◊ Fixed points among *some* sequences?

Sturmian words: *iff* quadratic slope (Crisp & Moran)

Sturmian hyperplane sequences: sufficient condition (F.)

## And now?

Defined: substitutions on multidimensional sequences.

◊ Fixed point sequences?

Classic substitutions: characterization by *return words* (Durand)

Local rules in 1D: nothing more (Fogg)

◊ Fixed points among *some* sequences?

Sturmian words: *iff* quadratic slope (Crisp & Moran)

Sturmian hyperplane sequences: sufficient condition (F.)