## Combinatorial Depth Measures

Patrick Schnider, ETH Zürich
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## Introduction



## Tukey and Tverberg



Tukey depth:
Minimum number of data points in any closed halfspace containing query point $q$

## Centerpoint theorem:

$$
\forall S \exists q: \operatorname{TD}(S, q) \geq \frac{|S|}{d+1}
$$



Tverberg depth:
Max. number of vertex disjoint simplices whose intersection contains $q$

Tverbergs theorem:
$\forall S \exists q: \operatorname{TvD}(S, q) \geq \frac{|S|}{d+1}$

## Combinatorial Depth Measures

$$
\begin{gathered}
\rho: S^{\mathbb{R}^{d}} \times \mathbb{R}^{d} \rightarrow \mathbb{R}_{\geq 0} \\
(S, q) \mapsto \rho(S, q)
\end{gathered}
$$

"combinatorial": depends only on relative position of $S$ and $q$ (order type), not on distances

Standard depth in $\mathbb{R}^{1}$ :


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## A "bad" measure

Convex hull peeling depth:

depth 2

no deep query point

## Super-additive Depth Measures

$\rho$ is super-additive if
(1) $\forall S, q, p:|\rho(S, q)-\rho(S \cup\{p\}, q)| \leq 1$
(2) $\forall S, q: \rho(S, q)=0$ if $q \notin \operatorname{conv}(S)$
(3) $\forall S, q: \rho(S, q) \geq 1$ if $q \in \operatorname{conv}(S)$
(4) $\forall S_{1} \sqcup S_{2}=S: \rho(S, q) \geq \rho\left(S_{1}, q\right)+\rho\left(S_{2}, q\right)$

Theorem [S', '23]:
Let $\rho$ be a super-additive depth measure. Then $\forall S, q$
$\operatorname{TD}(S, q) \geq \rho(S, q) \geq \operatorname{TvD}(S, q) \geq \frac{1}{d} \operatorname{TD}(S, q)$.

## Central Depth Measures

$\rho$ is central if
(1) $\forall S, q, p:|\rho(S, q)-\rho(S \cup\{p\}, q)| \leq 1$
(2) $\forall S, q: \rho(S, q)=0$ if $q \notin \operatorname{conv}(S)$
(3) $\forall S \exists q: \rho(S, q) \geq \frac{|S|}{d+1}$
(4) $\forall S, p, q: \rho(S \cup\{p\}, q) \geq \rho(S, q)$

Theorem [S', '23]:
Let $\rho$ be a central measure. Then $\exists c(d)$ s.t. $\forall S, q$
$\operatorname{TD}(S, q) \geq \rho(S, q) \geq c(d) \cdot \operatorname{TD}(S, q)$.

## Enclosing Depth $k$



## Enclosing depth:

$\mathrm{ED}(S, q)=\max k$ s.t. $q$ is $k$-enclosed

Lemma:
$\rho$ central. Then
$\rho(S, q) \geq \operatorname{ED}(S, q)-(d+1)$.

Lemma:
$\exists c(d)$ s.t.
$\mathrm{ED}(S, q) \geq c(d) \cdot \operatorname{TD}(S, q)$.

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## The 2D case

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## The 2D case



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The 2D case


## ЕНzürich

The 2D case


for every halfline: $b \geq r$

## ЕНzürich

The 2D case

$\mathrm{ED}(S, q) \geq \frac{1}{3} \mathrm{TD}(S, q)$.

for every halfline: $b \geq r$

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## The general case

Assume points are on sphere around $q$
Take witness-hyperplane for Tukey depth $k$
Project to larger side, color red and blue

for each half-space: $b \geq r$

Want:


## The Same Type Lemma

Theorem [Bárány, Valtr, '98]:
Let $X_{1}, \ldots, X_{m}$ be point sets in $\mathbb{R}^{d}$. Then there is a constant $c(d, m)$ and subsets $Y_{i} \subseteq X_{i}$ s.t.

- $\left|Y_{i}\right| \geq c \cdot\left|X_{i}\right|$ and
- each selection $y_{1} \in Y_{1}, \ldots, y_{m} \in Y_{m}$ has the same order type.



## Constant Fraction Radon

Theorem [S', '23]:
Let $P=R \cup B$ be a point set in $\mathbb{R}^{d}$ s.t. for every halfspace we have $b \geq r$. Then there is a constant $c(d)$ and subsets $R_{1}, \ldots, R_{a} \subseteq R$ and $B_{1}, \ldots, B_{b} \subseteq B$ s.t.

- $a+b=d+2$,
- $\left|R_{i}\right| \geq c \cdot|R|$ and $\left|B_{j}\right| \geq c \cdot|R|$ and
- for each selection
$r_{1} \in R_{1}, \ldots, r_{a} \in R_{a}, b_{1} \in B_{1}, \ldots, b_{b} \in B_{b}$ we have the sme order type and $\operatorname{conv}\left(r_{1}, \ldots, r_{a}\right) \cap \operatorname{conv}\left(b_{1}, \ldots, b_{b}\right) \neq \emptyset$.

Constant Fraction Radon in $\mathbb{R}^{d} \Rightarrow$ Enclosing Depth in $\mathbb{R}^{d+1}$

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Radon in 2D

$|R| / 3$

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Radon in 2D


Radon in 2D


Radon in 2D


Radon in 2D


Kirchberger: if the colors have a common intersection, then some $d+2$ elements do.

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## Radon: the general case


$|R| / 3$

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## Radon: the general case



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## Radon: the general case



$\mathrm{ED} \geq c \cdot \mathrm{TD}$

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## Radon: the general case



## Conclusion

- Combinatorial depth measures are intimately related to fundamental results in discrete geometry
- Many related algorithmic questions:
- Complexity of finding a Centerpoint/Tverberg point
- Complexity of computing Enclosing depth in general dimension? (Known $O\left(n^{d^{2}}\right)$ )
- Many depth measures are an approximation of Tukey depth
- Improve factors?


## Thank you!

