



Context: Formal Verification of Real-Time Systems

- Critical systems involve timing constants and concurrency
- Bugs can be dramatic (risk of loss of lives or huge financial loss)



⇒ Need for formal verification

Problem: what if the system constants are **uncertain** or are **not yet known**?

Solution: **parametric verification**

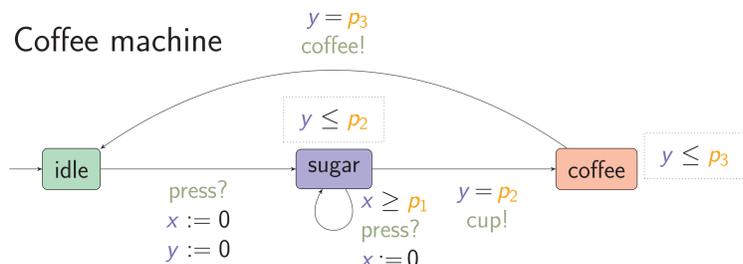
- Timing constants become **parameters**

Objective: derive values for these parameters ensuring the **absence of bug** (usually under the form of a set of constraints)

Parametric Timed Automata (PTA) [Alur et al., 1993]

- Finite automata (sets of locations and actions) extended with:
 - Clocks: real-valued variables evolving linearly
 - Parameters: unknown constants

- Example: Coffee machine



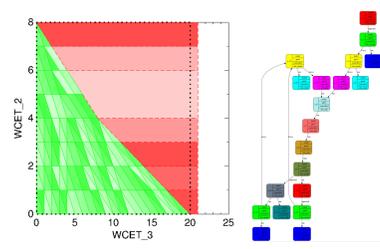
IMITATOR: Parameter Synthesis for Critical Systems

Input: a real-time system modeled by a network of PTA

Output: a constraint over the **parameters** guaranteeing the system correctness (e.g., non-reachability of some unsafe state)

Several algorithms:

- Non-reachability synthesis
- Parametric language preservation
- Behavioral cartography



Try IMITATOR! [André et al., 2012]

- Entirely written in OCaml
- Graphical outputs (behaviors, parameter constraints, etc.)
- Large repository of benchmarks
 - Asynchronous hardware circuits, scheduling problems, communication protocols, train controllers... and more!
- Available for free under the GNU-GPL license

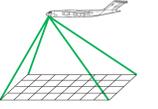
www.imitator.fr

What's next?

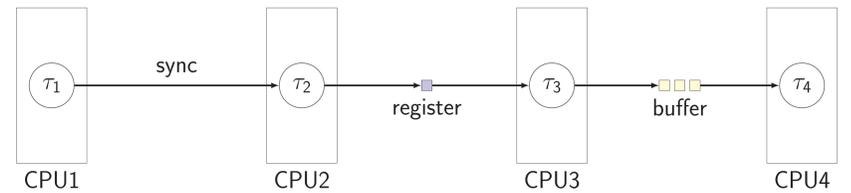
- Improved optimizations to address scalability
- Distributed and multi-core algorithms
- An input language for IMITATOR dedicated to real-time systems
 - Followed by a translation to PTA

A Case Study: The FMTV Challenge

- A problem proposed by Thales Research & Technology for the video capture in an aerial video system (2014)



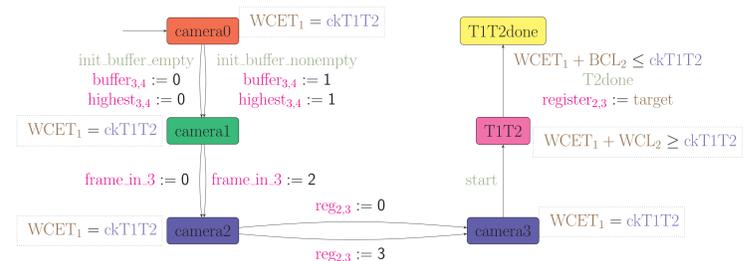
- A distributed video processing system (abstract view)



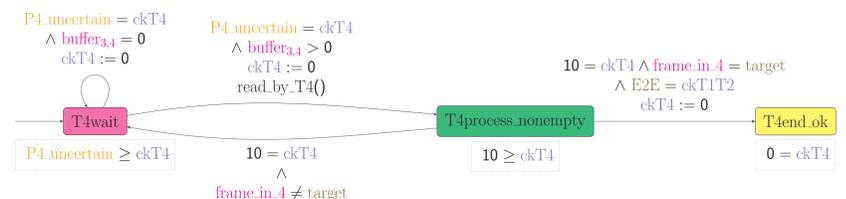
- τ_1 , τ_3 and τ_4 are periodic tasks
 - The exact value for each task's period is constant but **unknown**
 - $P1 \in [40 - 0.004 \text{ ms}, 40 + 0.004 \text{ ms}]$
 - $P3 \in [\frac{40}{3} - \frac{1}{150} \text{ ms}, \frac{40}{3} + \frac{1}{150} \text{ ms}]$
 - $P4 \in [40 - 0.004 \text{ ms}, 40 + 0.004 \text{ ms}]$
- τ_2 is triggered by the completion of τ_1
- The FIFO buffer between τ_3 and τ_4 has a size $n = 1$ or $n = 3$
- Challenge:** find the min/max end-to-end latency that a frame may experience in this system

Our Solution: Parametric Analysis [André et al., 2015]

- Task periods are modeled as **parameters**
 - E.g., $P4_{\text{uncertain}} \in [40 - 0.004 \text{ ms}, 40 + 0.004 \text{ ms}]$
- Another parameter: the end-to-end latency **E2E**
 - To focus on the **E2E** of an arbitrary frame (denoted as target)
- Some of the PTA modeling the system (for $n = 1$)
 - The system status is initialized to be arbitrary so that the worst-case and best-case scenarios for **E2E** will be included



- PTA model for task τ_4



- The end-to-end latency results returned by IMITATOR
 - $E2E \in [63 \text{ ms}, 145.008 \text{ ms}]$ (for $n = 1$)
 - $E2E \in [63 \text{ ms}, 225.016 \text{ ms}]$ (for $n = 3$)
- Runtime costs: 7.908 s with $n = 1$ and 115.247 s with $n = 3$

Conclusion

- Solved a problem with **uncertain timing constants** using **parametric analysis**, which turned out to be an efficient option

References

- Alur, R., Henzinger, T. A., and Vardi, M. Y. (1993). Parametric real-time reasoning. In *STOC*, pages 592–601. ACM.
- André, É., Fribourg, L., Kühne, U., and Soulat, R. (2012). IMITATOR 2.5: A tool for analyzing robustness in scheduling problems. In *FM*, volume 7436 of *Lecture Notes in Computer Science*, pages 33–36. Springer.
- André, É., Lipari, G., and Sun, Y. (2015). Verification of two real-time systems using parametric timed automata. In *WATERS*.