# Distribution of parameters in certain fragments of the linear and planar $\lambda$ -calculus

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Laboratoire Informatique de Paris Nord - CALIN Seminar

#### What is the $\lambda$ -calculus?

- · A *universal* formal system for expressing computation.
- Its terms are formed using the following grammar:
  - A variable is a valid term.
  - If x a variable and t is a valid term, then so is  $(\lambda x.t)$ .
- $\cdot$  If s and t are valid terms, then so is (s t).
- The  $\lambda$  calculus also provides us with tools to transform terms, including the operation of  $\beta$ -reduction:

$$((\lambda x.t) s) \stackrel{\beta}{\to} t[x := s]$$

Some examples of terms:

$$(\lambda x.(xx))(\lambda x.(xx) \lambda x.\lambda y.(x (x y)) \lambda x.(z (\lambda y.y))$$

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#### Combinatorics of the $\lambda$ -calculus

- General terms are quite complicated. Growth is super-exponential, generating functions are not analytic. <sup>1</sup> Asymptotic number of general terms still (?) unresolved!
- We focus on *linear* terms: bound variables must appear exactly once:  $\lambda x.(x x), \lambda x. \lambda y.(a (y x)).$
- We also consider planar terms: bound variables must appear in the order they are introduced:
   λx.λy.(y x), λx.λy.(a (x y)).

<sup>&</sup>lt;sup>1</sup>For the notion of term size given recursively by: |var|=1, |(s t)|=|s|+|t|+1,  $|\lambda x.t|=|t|+1$ .

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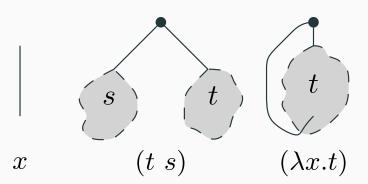
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#### The $\lambda$ -calculus and maps

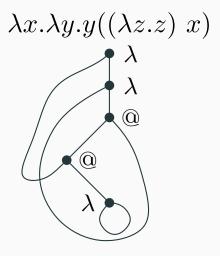
- Maps: graphs embedded in an oriented surface without boundary.
- Closed linear terms are combinatorially intriguing: they correspond to rooted connected trivalent maps! [1, 2]
   Closed planar terms correspond to planar such maps. Open terms allow for univalent vertices too.

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#### An example of a term and its corresponding map



Where  $\lambda$  annotates abstractions and  $\mathbf{0}$  applications.

#### Purpose of this work

- How do "typical" (random, of large size) linear and planar terms behave?
- How many free variables do they have? How often is a typical term an abstraction?
- Using tools from analytic combinatorics to obtain parameter distributions.

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#### In this talk

#### We'll sketch the following results:

Linear  $\lambda$ -Terms (Differentially Algebraic, Divergent)

- Limit distribution of free variables
- Limit distribution of identity-subterms in closed terms.
- Limit distribution of closed subterms in closed terms.
- Probability that term is an abstraction.

Planar  $\lambda$ -Terms (Algebraic, Analytic)

- Limit distribution of free variables for regular and bridgless terms.
- Probability that regular or bridgless open term is an abstraction.

#### Free variables in closed linear terms

• Free variables are those not bound by an abstraction. For example:  $\lambda x.(a x)$ 

#### Proposition

The limit distribution of free variables in linear  $\lambda$ -terms of size n is Gaussian with mean and variance  $\mu = \sigma^2 \sim \sqrt[3]{n}$ .

Starting point (follows from definition of combinatorial maps):

$$L(z^2, u) = uz^2 + z^4 + z^5 \frac{\partial}{\partial z} \left( \ln \left( \exp(z^2/2) \odot \exp(z^3/3 + uz) \right) \right)$$

where L counts open linear  $\lambda$ -terms with u tagging free variables

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#### Free variables in closed linear terms

#### **Proof Sketch**

Saddle-point analysis of Hadamard product yields:

$$[z^n] \exp\left(z^3/3 + uz\right) = \left(\frac{1}{6} \frac{\sqrt{2}\sqrt{3} \, n^{-\frac{1}{2}}}{\sqrt{\pi}} - \frac{1}{36} \frac{\sqrt{2}\sqrt{3}u^2 n^{-\frac{5}{6}}}{\sqrt{\pi}} + O\left(n^{-\frac{7}{6}}\right)\right) e^{un^{1/3} + n/3} n^{-n/3}$$

$$[z^n] \exp\left(z^2/2\right) \sim \frac{1}{2} \frac{e^{1/2 + n/2}}{\left(\sqrt{1 + n}\right)^{1 + n} \sqrt{\pi}} - \frac{1}{2} \frac{e^{1/2 + n/2}}{\left(-\sqrt{1 + n}\right)^{1 + n} \sqrt{\pi}}$$

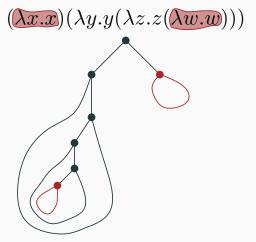
While an application of Bender's theorem [3, Theorem 1] gives

$$2[z^n]\frac{d}{dz}\ln(A(z^{1/2},u)) = n\left([z^n]h(z,u) - \frac{1}{2}[z^{n-2}]h(z,u)\right) + O\left([z^{n-4}]h(z,u)\right)$$

for 
$$A(x, u) = \exp(z^2/2) \odot \exp(z^3/3 + uz)$$

- Identity terms: terms which are  $\alpha$ -equivalent to  $\lambda x.x$ . For example:  $\lambda x.(x(\lambda y.y))$ .
- · They appear as loops in the corresponding map.

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#### Proposition

The limit distribution of identity-subterms in closed linear  $\lambda$ -terms is Poisson of parameter  $\lambda = 1$ .

Proof Sketch: Use moment pumping on

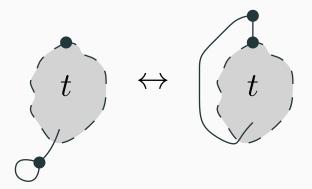
$$G = (u - 1)z^2 + zG^2 + \frac{\partial}{\partial u}G$$

where G counts closed linear terms with u tagging identity-subterms.

Justification for

$$G = (u - 1)z^2 + zG^2 + \frac{\partial}{\partial u}G$$

Terms are either identity-terms, applications, or



For the pumping, note that the *k*-th derivative of the eq. may be written as

$$\frac{\partial^k}{\partial u^k}G - S - 2z G \frac{\partial^k}{\partial u^k}G = \frac{\partial^{k+1}}{\partial u^{k+1}}G$$

with S, depending on the parity of k, being as follows

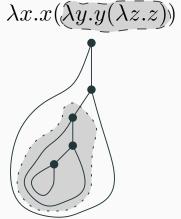
$$\begin{split} &\sum_{l=1}^{\lfloor \frac{k}{2} \rfloor} 2z \binom{k}{l} \frac{\partial^{l}}{\partial u^{l}} G \frac{\partial^{k-l}}{\partial u^{k-l}} G, \text{ for odd } k \\ &\sum_{l=1}^{\lfloor \frac{k}{2} \rfloor - 1} 2z \binom{k}{l} \frac{\partial^{l}}{\partial u^{l}} G \frac{\partial^{k-l}}{\partial u^{k-l}} G + z \binom{k}{\lfloor \frac{k}{2} \rfloor} \left( \frac{\partial^{\lfloor \frac{k}{2} \rfloor}}{\partial u^{\lfloor \frac{k}{2} \rfloor}} G \right)^{2}, \text{ for even } k \end{split}$$

#### Distribution of closed subterms in closed linear terms

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- Coresponding to non-root-containing connected components resulting from the deletion of some bridge in the respective map.

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#### Distribution of closed subterms in closed linear terms

#### **Proposition**

The limit distribution of closed subterms in closed linear  $\lambda$ -terms is Poisson of parameter  $\lambda = 1$ .

Proof Sketch: Use moment pumping on

$$\frac{\partial W}{\partial v} = \frac{-(zv^2W^2 + z^2 - W)W}{zv^2(v - 1)W^2 + (1 - v)W + vz^2}$$

where *W* counts closed linear terms with *v* tagging identity-subterms.

#### Probability that a closed linear term is an abstraction

#### Proposition

Asymptotically almost surely a random closed linear  $\lambda$ -term is an abstraction.

#### Proof Sketch:

It can be shown that  $[z^n]L_c \sim k \cdot 6^n \cdot n!$  for some constant k. Compare the coefficients of  $L_c$  and  $2z^4 \frac{\partial}{\partial z} L_c$  in

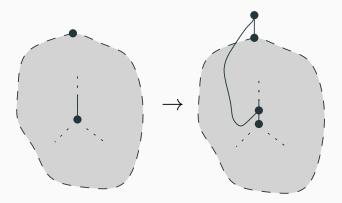
$$L_c = z^2 + zL_c^2 + 2z^4 \frac{\partial}{\partial z} L_c.$$

where  $L_c$  enumerates closed linear  $\lambda$ -terms.

Justification for

$$G = L_c = z^2 + zL_c^2 + 2z^4 \frac{\partial}{\partial z} L_c.$$

Terms are either identity-terms, applications, or



#### Proposition

The limit distribution of free variables in planar  $\lambda$ -terms of size n is Gaussian with mean  $\mu = \frac{n}{8}$  and variance  $\sigma^2 = \frac{9n}{32}$ .

#### **Proposition**

The limit distribution of free variables in bridgeless planar  $\lambda$ -terms of size n is Gaussian with mean  $\mu = \frac{n}{5}$  and variance  $\sigma^2 = \frac{9n}{25}$ .

Both results follow similar steps.

Our starting points are the following two equations

$$P(z, u) = uz + zQ(z, u)^{2} + \frac{z(P(z, u) - P(z, 0))}{u}$$

$$Q(z, u) = uz + zQ(z, u)^{2} + \frac{z(Q(z, u) - u[u^{1}]Q(z, u))}{u}$$

with *P* and *Q* counting planar and bridgless planar terms respectively and *u* tagging free variables.

Sketch: use elimination and the quadratic method to obtain closed form solutions. Proceed by applying, [4, Proposition IX.17].

$$Q(z, u) = 1/2 z^{-1} - 1/2 u^{-1}$$

$$+ \frac{1}{2} \frac{1}{uz} \left( \frac{1}{3} u^2 \sqrt[3]{-1458 z^6 + 6 \sqrt{3} \sqrt{19683 z^8 - 4374 z^5 + 324 z^2 - 8 z^{-1} z^2 - 270 z^3 + 1} \right)$$

$$+ 36 u^2 z^3 \frac{1}{\sqrt[3]{-1458 z^6 + 6 \sqrt{3} \sqrt{19683 z^8 - 4374 z^5 + 324 z^2 - 8 z^{-1} z^2 - 270 z^3 + 1}}$$

$$+ \frac{1}{3} u^2 \frac{1}{\sqrt[3]{-1458 z^6 + 6 \sqrt{3} \sqrt{19683 z^8 - 4374 z^5 + 324 z^2 - 8 z^{-1} z^2 - 270 z^3 + 1}}$$

$$+ \frac{1}{3} u^2 - 4 u^3 z^2 - 2 uz + z^2 \right)^{1/2}.$$

While 
$$P(z) = A(z, u) + B(z, u) \cdot C(z, u)^{-1/2}$$
 with
$$A(z, u) = \frac{1}{2z} - \frac{1}{2u}, \ B(z, u) = \frac{1}{2uz}$$

$$C(Z, u) = -4u^3z^2$$

$$+ \frac{1}{48} \frac{u\sqrt[3]{1492992}z^{12} + 8640z^6 + 96\sqrt[3]{80621568}z^{18} - 559872z^{12} + 1296z^6 - 1z^3 - 1}{z^2}$$

$$+ 72 \frac{uz^4}{\sqrt[3]{1492992}z^{12} + 8640z^6 + 96\sqrt[3]{80621568}z^{18} - 559872z^{12} + 1296z^6 - 1z^3 - 1}$$

$$+ \frac{1}{48} \frac{u}{z^2}\sqrt[3]{1492992}z^{12} + 8640z^6 + 96\sqrt[3]{80621568}z^{18} - 559872z^{12} + 1296z^6 - 1z^3 - 1}$$

$$- \frac{1}{48} \frac{u}{z^2} + u^2 + z^2$$

## Probability that an open planar or bridgeless planar term is an abstraction

#### Proposition

Asymptotically, the probability that an random open planar (bridgeless planar) term is an abstraction is  $\rho_P = \frac{\sqrt{2}}{4}$  ( $\rho_{PB} = \frac{2}{5}$ ).

Proof Sketch: Estimate

$$\frac{[z^n] z(P(z,1) - P(z,0))}{[z^n] P(z,1)} \text{ and } \frac{[z^n] z(Q(z,1) - ([u^1]Q(z,u))|_{u=1})}{[z^n] Q(z,1)}$$

Both P(z,0) and  $[u^1]Q(z,u)|_{u=1}$  are analytic at the respective singularities  $\rho_P$  and  $\rho_{PB}$  of P and Q. Use the singular expansions of P,Q at the corresponding singularities to obtain the desired result.

#### Conclusions

- Clear distinctions between the divergent/differentially-algebraic case of linear terms and the algebraic one of planar terms.
- Need for: more tools to handle divergent combinatorial classes, algebraicity results for closed planar terms.
- Future directions: study of  $\beta$ -reduction, typing of linear terms.

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