

CAV 2013

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PSyHCoS

Parameter Synthesis for Hierarchical Concurrent Real-Time Systems

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Motivation

- Timed systems are characterized by a **set of timing constants**
 - “The packet transmission lasts for **50 ms**”
 - “The sensor reads the value every **10 s**”
 - etc.
- Verification for **one** set of constants does not guarantee the correctness for other values
- Challenges
 - **Numerous verifications**: is the system correct for any value within **[40; 60]**?
 - **Optimization**: until what value can we increase **10**?
 - **Robustness**: What happens if **50** is implemented with **49.99**?

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- **Parameter synthesis**
 - Consider that timing constants are unknown constants (**parameters**)
 - Find **good values** for the parameters

Parametric Stateful Timed CSP

- An **intuitive formal modeling language**
 - English-like description
 - Formal semantics allowing verification
 - Standard constructions of **CSP** [Hoare, 1978]
 - User-defined **data structures** (C#-like code)

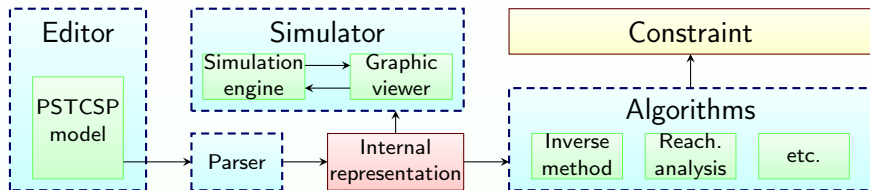
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 - Standard constructions of **CSP** [Hoare, 1978]
 - User-defined **data structures** (C#-like code)
- Parameterized **timed constructs** [André et al., 2012, Sun et al., 2013]
 - **Wait**[*u*]: waits exactly *u* time units
 - **P timeout**[*u*] **Q**: the first observable event of **P** shall occur before *u* time units; otherwise, behaves like **Q**
 - **P interrupt**[*u*] **Q**: behaves like **P** until *u* time units; then, like **Q**
 - **P within**[*u*]: the first observable event of **P** shall occur before *u* time units
 - **P deadline**[*u*]: **P** shall terminate before *u* time units

Algorithms Implemented

- Implementation in PSyHCoS
 - Parameter Synthesis for Hierarchical Concurrent Systems
- Computation of the reachability graph
 - 😊 Interesting for small examples
 - 😞 Often leads to an infinite state space
 - ↪ Does not terminate (no synthesis possible)
- Synthesis using the inverse method [André and Soulat, 2013]
 - 😊 Partial exploration of the state space only
 - 😊 Outputs a **constraint** on the **parameters**: avoiding numerous verification, allowing optimization and robustness analysis
 - 😊 Often terminate in practice
- And also: classical model checking algorithms
 - LTL / deadlock-checking, etc.

Architecture of PSyHCoS



- Implemented in C# (Microsoft .NET framework)
- Each syntactic construct (with its semantics) implemented in a different class
- Algorithms implemented in a modular way
 - ~ Easy **reusability** and addition of **new features**

Experiments

Case study	U	<i>reachAll</i>				<i>reachAll+</i>				<i>IM</i>			<i>IM+</i>		
		S	T	X	t	S	T	X	t	S	X	t	S	X	t
Bridge	4	-	-	-	M.O.	-	-	-	M.O.	2.8k	2	253	2.8k	2	455
Fischer ₄	2	-	-	-	M.O.	-	-	-	M.O.	11k	4	41.9	2k	4	8.65
Fischer ₅	2	-	-	-	M.O.	-	-	-	M.O.	133k	5	1176	13k	5	84.5
Fischer ₆	2	-	-	-	M.O.	-	-	-	M.O.	-	-	M.O.	86k	6	1144
Jobshop	8	14k	20k	2	21.0	12k	17k	2	18.1	1112	2	17.1	877	2	22.8
RCS ₅	4	5.6k	7.2k	4	10.5	5.6k	7.2k	4	9.54	5.6k	4	7.83	5.6k	4	16.7
RCS ₆	4	34k	43k	4	91.7	34k	43k	4	54.5	34k	4	60.4	34k	4	91.3
TrAHV	6	7.2k	13k	6	14.2	7.2k	13k	6	15.8	227	6	0.555	227	6	0.655

- *reachAll*: computation of the reachability graph
- *IM*: inverse method
- *reachAll+* (resp. *IM+*): version with optimized encoding

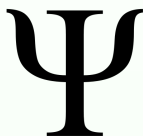
Perspectives

- Integration of further **state space reduction** techniques
[André et al., 2013]
- Improvement of the internal **representation of constraints** relying on the Parma Polyhedra Library [Bagnara et al., 2008]
- Parameterized **refinement checking**
- Extension to the **multi-core** setting [Laarman et al., 2013]

Try it!

Available under the GNU-GPL license

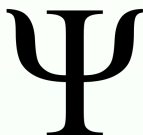
<http://lipn.univ-paris13.fr/~andre/software/PSyHCoS/>



Try it!

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Demo today at 5pm

FM 2014

19th International Symposium on Formal Methods (FM 2014)

- 12-16, May, 2014
- Singapore
- <http://www.comp.nus.edu.sg/~pat/FM2014/>



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Configuration encoding

Configuration encoding

- Encoding

- Process (ID)
- Value for variables
- List of clocks
- Constraint: definition of a normal form

- Example

- $(\text{Wait}[u_3]_{x_3} \parallel \text{Wait}[u_5]_{x_3} \parallel \text{Wait}[u_5]_{x_2}, x_3 \leq x_2)$

- Encoding:

- Process: $\text{Wait}[u_3] \parallel \text{Wait}[u_5] \parallel \text{Wait}[u_5]$
- List of clocks: $\{x_3, x_3, x_2\}$
- Constraint: $x_3 \leq x_2$

- Justification for the list of clocks

- Distinguishes between $(\text{Wait}[u_3]_{x_3} \parallel \text{Wait}[u_5]_{x_3} \parallel \text{Wait}[u_5]_{x_2}, x_3 \leq x_2)$
and $(\text{Wait}[u_3]_{x_2} \parallel \text{Wait}[u_5]_{x_3} \parallel \text{Wait}[u_5]_{x_2}, x_3 \leq x_2)$

Configuration encoding: optimization

- Actual equivalence between
 $(\text{Wait}[u_3]_{x_1} \parallel \text{Wait}[u_5]_{x_2}, x_1 \leq x_2)$ and
 $(\text{Wait}[u_3]_{x_2} \parallel \text{Wait}[u_5]_{x_1}, x_2 \leq x_1)$

Configuration encoding: optimization

- Actual equivalence between
 $(\text{Wait}[u_3]_{x_1} \parallel \text{Wait}[u_5]_{x_2}, x_1 \leq x_2)$ and
 $(\text{Wait}[u_3]_{x_2} \parallel \text{Wait}[u_5]_{x_1}, x_2 \leq x_1)$
- Idea \leadsto rename clocks
 - Example: $(\text{Wait}[u_3]_{x_3} \parallel \text{Wait}[u_5]_{x_3} \parallel \text{Wait}[u_5]_{x_2}, x_3 \leq x_2)$
 New encoding:
 - Process: $\text{Wait}[u_3] \parallel \text{Wait}[u_5] \parallel \text{Wait}[u_5]$
 - List of clocks: $\{x_1, x_1, x_2\}$
 - Constraint: $x_1 \leq x_2$
- This method is time consuming
 - Numerous string and list sorts
 - But often leads to efficient state space reduction

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