

Bipartite subfamilies of planar graphs

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CONSEJO SUPERIOR
DE INVESTIGACIONES
CIENTÍFICAS



The material of this talk

- 1.— **Background**
- 2.— **Graph decompositions. First results**
- 3.— **The bipartite framework**

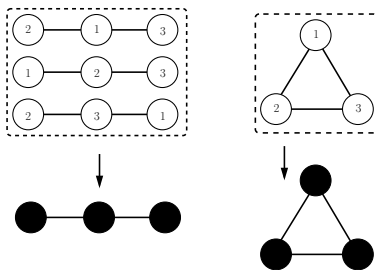
Background

Objects: graphs

Labelled Graph = labelled vertices + edges.

Unlabelled Graph = labelled one up to permutation of labels.

Simple Graph = NO multiples edges, NO loops.



Question: How many graphs with n vertices are in the family?

The counting series

Strategy: Encapsulate these numbers \rightarrow Counting series

- ▶ Labelled framework: exponential generating functions

$$A(x) = \sum_{a \in \mathcal{A}} \frac{x^{|a|}}{|a|!} = \sum_{n \geq 0} \frac{|\mathcal{A}_n|}{n!} x^n$$

- ▶ Unlabelled framework: cycle index sums

$$Z_{\mathcal{A}}(s_1, s_2, \dots) = \sum_{n \geq 0} \frac{1}{n!} \sum_{\substack{(\sigma, g) \in \mathfrak{S}_n \times \mathcal{A}_n \\ \sigma \cdot g = g}} s_1^{c_1} s_2^{c_2} \cdots s_n^{c_n},$$

$$\tilde{A}(x) = Z_{\mathcal{A}}(x, x^2, x^3, \dots) = \sum_{n \geq 0} |\tilde{\mathcal{A}}_n| x^n.$$

The symbolic method

COMBINATORIAL RELATIONS between CLASSES



EQUATIONS between GENERATING FUNCTIONS

Class	Labelled setting	Unlabelled setting
$\mathcal{C} = \mathcal{A} \cup \mathcal{B}$	$C(x) = A(x) + B(x)$	$\tilde{C}(x) = \tilde{A}(x) + \tilde{B}(x)$
$\mathcal{C} = \mathcal{A} \times \mathcal{B}$	$C(x) = A(x) \cdot B(x)$	$\tilde{C}(x) = \tilde{A}(x) \cdot \tilde{B}(x)$
$\mathcal{C} = \text{Set}(\mathcal{B})$	$C(x) = \exp(B(x))$	$\tilde{C}(x) = \exp\left(\sum_{i \geq 1} \frac{1}{i} \tilde{B}(x^i)\right)$
$\mathcal{C} = \mathcal{A} \circ \mathcal{B}$	$C(x) = A(B(x))$	$\tilde{C}(x) = Z_{\mathcal{A}}(\tilde{B}(x), \tilde{B}(x^2), \dots)$

Singularity analysis on generating functions

GFs: analytic functions in a neighbourhood of the origin.

The smallest singularity of $A(z)$ determines the asymptotics of the coefficients of $A(z)$.

- ▶ **POSITION:** exponential growth ρ .
- ▶ **NATURE:** subexponential growth
- ▶ **Transfer Theorems:** Let $\alpha \notin \{0, -1, -2, \dots\}$. If

$$A(z) = a \cdot (1 - z/\rho)^{-\alpha} + o((1 - z/\rho)^{-\alpha})$$

then

$$a_n = [z^n]A(z) \sim \frac{a}{\Gamma(\alpha)} \cdot n^{\alpha-1} \cdot \rho^{-n}(1 + o(1))$$

Our starting point

Asymptotic enumeration and limit laws of planar graphs (Giménez, Noy)

$$g_1 \cdot n^{-7/2} \cdot \gamma_1^n \cdot n! \cdot (1 + o(1))$$

Asymptotic enumeration and limit laws of series-parallel graphs (Bodirsky, Giménez, Kang, Noy)

$$g_2 \cdot n^{-5/2} \cdot \gamma_2^n \cdot n! \cdot (1 + o(1))$$

Our starting point

$$g_1 \cdot n^{-7/2} \cdot \gamma_1^n \cdot n! \cdot (1 + o(1))$$

$$g_2 \cdot n^{-5/2} \cdot \gamma_2^n \cdot n! \cdot (1 + o(1))$$

Our starting point

$$g_1 \cdot n^{-7/2} \cdot \gamma_1^n \cdot n! \cdot (1 + o(1))$$

$$g_2 \cdot n^{-5/2} \cdot \gamma_2^n \cdot n! \cdot (1 + o(1))$$



**THE SUBEXPONENTIAL TERM GIVES THE
“PHYSICS” OF THE GRAPHS**



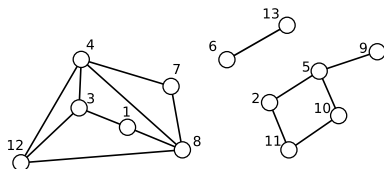
**GENERAL FRAMEWORK TO UNDERSTAND THIS
EXPONENT**

Graph decompositions. First results

General graphs from connected graphs

Let \mathcal{C} be a family of *connected* graphs.

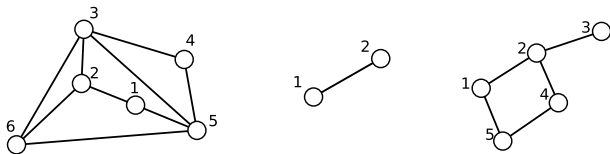
\mathcal{G} : graphs such that their *connected components* are in \mathcal{C} .



General graphs from connected graphs

Let \mathcal{C} be a family of *connected* graphs.

\mathcal{G} : graphs such that their *connected components* are in \mathcal{C} .

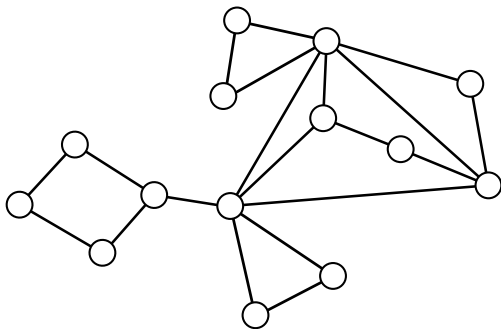


$$\mathcal{G} = \text{Set}(\mathcal{C}) \implies G(x) = \exp(C(x))$$

Connected graphs from 2-connected graphs

Let \mathcal{B} be a family of 2-connected graphs.

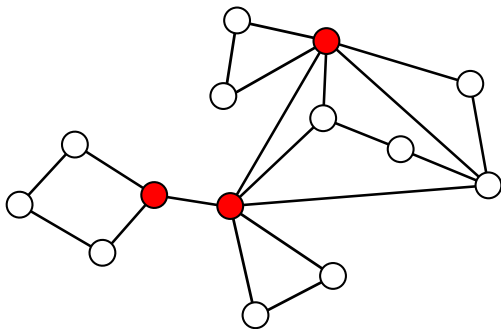
\mathcal{C} : connected graphs with blocks in \mathcal{B} .



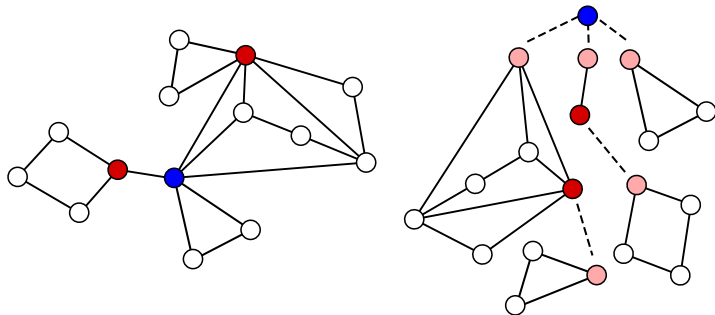
Connected graphs from 2-connected graphs

Let \mathcal{B} be a family of 2-connected graphs.

\mathcal{C} : connected graphs with blocks in \mathcal{B} .



Connected graphs from 2-connected graphs



A *vertex-rooted* connected graph is a *tree* of rooted blocks.

$$\mathcal{C}^\bullet = v \times \text{Set}(\mathcal{B}'(v \leftarrow \mathcal{C}^\bullet)) \implies \mathcal{C}^\bullet(x) = x \exp B'(\mathcal{C}^\bullet(x))$$

2-connected graphs from 3-connected graphs

Decomposition in 3-connected components is slightly harder.

Let \mathcal{T} be a family of 3-connected graphs: $T(x, z)$.

We define \mathcal{B} as those 2-connected graphs such that can be obtained from *series*, *parallel*, and \mathcal{T} -compositions.

$$D(x, y) = (1 + y) \exp \left(\frac{x D^2}{1 + x D} + \frac{1}{2x^2} \frac{\partial T}{\partial z}(x, D) \right) - 1$$

$$\frac{\partial B}{\partial y}(x, y) = \frac{x^2}{2} \left(\frac{1 + D(x, y)}{1 + y} \right)$$

D is the GF for networks (essentially edge-rooted 2-connected graphs without the edge root).


A set of equations

$$\left\{ \begin{array}{l} \frac{1}{2x^2D} \frac{\partial T}{\partial z}(x, D) - \log\left(\frac{1+D}{1+y}\right) + \frac{x D^2}{1+xD} = 0 \\ \frac{\partial B}{\partial y}(x, y) = \frac{x^2}{2} \left(\frac{1+D(x, y)}{1+y} \right) \end{array} \right.$$

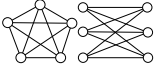
$$\left\{ \begin{array}{l} C^\bullet(x) = x \exp(B'(C^\bullet(x))) \\ G(x) = \exp(C(x)) \end{array} \right.$$

Examples of families & excluded minors (I)

- ▶ Series-parallel graphs

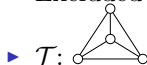
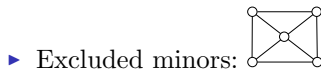
- ▶ Excluded minors: 
- ▶ \mathcal{T} : None.
- ▶ $T(x, z) = 0$.

- ▶ Planar graphs

- ▶ Excluded minors: 
- ▶ \mathcal{T} : 3-connected planar graphs.
- ▶ $T(x, z)$: *The number of labelled 2-connected planar graphs*
(Bender, Gao, Wormald, 2002)

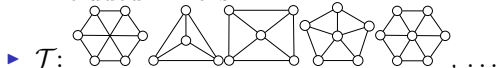
Examples of families & excluded minors (II)

- ▶ W_4 -free



- ▶ $T(x, z) = \frac{1}{4!}x^4z^6$.

- ▶ K_5^- -free



- ▶ $T(x, z) = \frac{70}{6!}x^6z^9 - \frac{1}{2}x(\log(1 - xz^2) + 2xz^2 + x^2z^4)$.

Examples of families & excluded minors (III)

- ▶ $K_{3,3}$ -free (Gerke, Giménez, Noy, Weibl, 2006)



- ▶ Excluded minors:



- ▶ 3-connected components: , 3-connected planar graphs.
- ▶ $T(x, z) = \dots$

- ▶ If $\mathcal{G} = \text{Ex}(\mathcal{M})$ and all the excluded minors \mathcal{M} are 3-connected, then \mathcal{G} can be expressed in terms of its 3-connected graphs.

RESULT: asymptotic enumeration

If either $\frac{\partial T}{\partial z}(x, z)$

- ▶ has no singularity, or
- ▶ the singularity type is $(1 - z/z_0)^\alpha$ with $\alpha < 1$,

then the situation is alike to the **series-parallel case**:

$$D(x) \sim d \cdot (1 - x/x_0)^{1/2}$$

$$d_n \sim d \cdot n^{-3/2} \cdot x_0^{-n} \cdot n!$$

$$B(x) \sim b \cdot (1 - x/x_0)^{3/2}$$

$$b_n \sim b \cdot n^{-5/2} \cdot x_0^{-n} \cdot n!$$

$$C(x) \sim c \cdot (1 - x/\rho)^{3/2}$$

$$c_n \sim c \cdot n^{-5/2} \cdot \rho^{-n} \cdot n!$$

$$G(x) \sim g \cdot (1 - x/\rho)^{3/2}$$

$$g_n \sim g \cdot n^{-5/2} \cdot \rho^{-n} \cdot n!$$

RESULT: asymptotic enumeration (II)

If $\frac{\partial T}{\partial z}(x, z)$ has singularity type $(1 - z/z_0)^{3/2}$, then 3 different situations may happen.

Case 1 (Planar case)

$$D(x) \sim d \cdot (1 - x/x_0)^{3/2}$$

$$d_n \sim d \cdot n^{-5/2} \cdot x_0^{-n} \cdot n!$$

$$B(x) \sim b \cdot (1 - x/x_0)^{5/2}$$

$$b_n \sim b \cdot n^{-7/2} \cdot x_0^{-n} \cdot n!$$

$$C(x) \sim c \cdot (1 - x/\rho)^{5/2}$$

$$c_n \sim c \cdot n^{-7/2} \cdot \rho^{-n} \cdot n!$$

$$G(x) \sim g \cdot (1 - x/\rho)^{5/2}$$

$$g_n \sim g \cdot n^{-7/2} \cdot \rho^{-n} \cdot n!$$

RESULT: asymptotic enumeration (II)

If $\frac{\partial T}{\partial z}(x, z)$ has singularity type $(1 - z/z_0)^{3/2}$, then 3 different situations may happen.

Case 2 (**Series-parallel case**)

$$D(x) \sim d \cdot (1 - x/x_0)^{1/2}$$

$$d_n \sim d \cdot n^{-3/2} \cdot x_0^{-n} \cdot n!$$

$$B(x) \sim b \cdot (1 - x/x_0)^{3/2}$$

$$b_n \sim b \cdot n^{-5/2} \cdot x_0^{-n} \cdot n!$$

$$C(x) \sim c \cdot (1 - x/\rho)^{3/2}$$

$$c_n \sim c \cdot n^{-5/2} \cdot \rho^{-n} \cdot n!$$

$$G(x) \sim g \cdot (1 - x/\rho)^{3/2}$$

$$g_n \sim g \cdot n^{-5/2} \cdot \rho^{-n} \cdot n!$$

RESULT: asymptotic enumeration (II)

If $\frac{\partial T}{\partial z}(x, z)$ has singularity type $(1 - z/z_0)^{3/2}$, then 3 different situations may happen.

Case 3 (Mixed case)

$$D(x) \sim d \cdot (1 - x/x_0)^{3/2}$$

$$d_n \sim d \cdot n^{-5/2} \cdot x_0^{-n} \cdot n!$$

$$B(x) \sim b \cdot (1 - x/x_0)^{5/2}$$

$$b_n \sim b \cdot n^{-7/2} \cdot x_0^{-n} \cdot n!$$

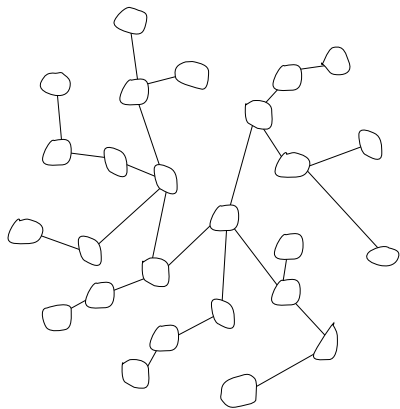
$$C(x) \sim c \cdot (1 - x/\rho)^{3/2}$$

$$c_n \sim c \cdot n^{-5/2} \cdot \rho^{-n} \cdot n!$$

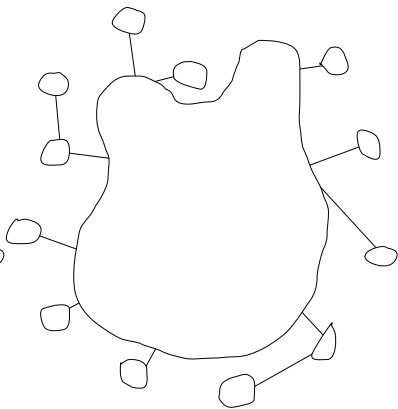
$$G(x) \sim g \cdot (1 - x/\rho)^{3/2}$$

$$g_n \sim g \cdot n^{-5/2} \cdot \rho^{-n} \cdot n!$$

2 different pictures



Series-parallel-like situation

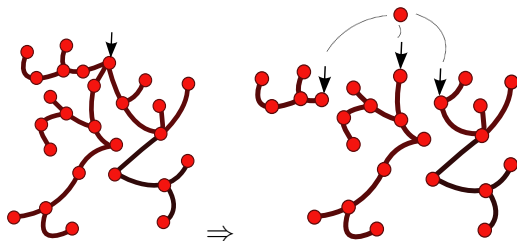


Planar-like situation

The bipartite framework

A key example: Trees

We count *rooted* trees



$$\mathcal{T} = \bullet \times \text{Set}(\mathcal{T}) \rightarrow T(x) = xe^{T(x)}$$

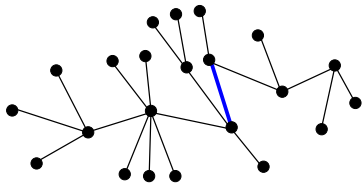
To forget the root, we just integrate: $(xU'(x) = T(x))$

$$\int_0^x \frac{T(s)}{s} ds = \left\{ \begin{array}{l} T(s) = u \\ T'(s) ds = du \end{array} \right\} = \int_{T(0)}^{T(x)} 1-u du = T(x) - \frac{1}{2}T(x)^2$$

Question: can we interpret this formula **combinatorially**?

The dissymmetry theorem

Let \mathcal{T} a class of unrooted trees \Rightarrow canonical root (their centers).



Dissymmetry Theorem for trees:

$$\mathcal{T} \cup \mathcal{T}_{\bullet \rightarrow \bullet} \simeq \mathcal{T}_{\bullet - \bullet} \cup \mathcal{T}_{\bullet},$$

For trees:

$$\mathcal{T}_{\bullet \rightarrow \bullet} \rightarrow T(x)^2; \mathcal{T}_{\bullet - \bullet} \rightarrow \frac{1}{2}T(x)^2; \mathcal{T}_{\bullet} \rightarrow T(x).$$

Dissymmetry Theorem \equiv Combinatorial **Integration**.

Returning to the equations

$$\frac{\partial B}{\partial y}(x, y) = \frac{x^2}{2} \left(\frac{1 + D(x, y)}{1 + y} \right) \leftrightarrow 2(1 + y) \frac{\partial B}{\partial y}(x, y) = \frac{x^2}{2} (1 + D(x, y))$$

↓

$$B(x, y) = \frac{x^2}{2} \int_0^y \left(\frac{1 + D(x, s)}{1 + s} \right) ds$$

Amazingly, an **EXACT** formula exists!

$$B(x, y) = T(x, D(x, y)) - \frac{1}{2} x D(x, y) + \frac{1}{2} \log(1 + x D(x, y)) + \frac{x^2}{2} \left(D(x, y) + \frac{1}{2} D(x, y)^2 + (1 + D(x, y)) \log \left(\frac{1 + y}{1 + D(x, y)} \right) \right).$$

Is there a “tree-like” argument to explain this formula?

The complete grammar for graphs

A Grammar for Decomposing a Family of Graphs into 3-connected Components; (Chapuy, Fusy, Kang, Shoilekova)

(1). General from Connected (folklore). $\mathcal{G} = \text{Set}(\mathcal{G}_1)$
(2). Connected from 2-Connected (Bergeron, Labelle, Leroux). $\mathcal{G}_1 = \mathcal{C} - \mathcal{C}_v + \mathcal{C}_D - \mathcal{C}_{Dv}$ [dissymmetry theorem] $\mathcal{C}_v = v * \mathcal{C}'$ $\mathcal{C}' = \text{Set}(\mathcal{G}'_v \circ_v \mathcal{C}')$ $\mathcal{C}_D = \mathcal{G}'_2 \circ_v \mathcal{C}'$ $\mathcal{C}_{Dv} = (v * \mathcal{G}'_2) \circ_v \mathcal{C}'$
(3). 2-Connected from 3-Connected. (i) Networks $\mathcal{D} = e + \mathcal{S} + \mathcal{P} + \mathcal{H}$ $\mathcal{S} = (\mathcal{D} - \mathcal{S}) * v * \mathcal{D}$ $\mathcal{P} = \begin{cases} \text{Set}_{\geq 2}(\mathcal{D} - \mathcal{P}), & \text{[Multi-edges allowed]} \\ e * \text{Set}_{\geq 1}(\mathcal{D} - \mathcal{P} - e) + \text{Set}_{\geq 2}(\mathcal{D} - \mathcal{P} - e), & \text{[No Multi-edges]} \end{cases}$ $\mathcal{H} = \overline{\mathcal{G}}_3 \circ_e \mathcal{D}$ (ii) Unrooted 2-Connected $\mathcal{G}_2 = e + \mathcal{B}, \quad e = \begin{cases} \ell_1 + \ell_2 & \text{[Multi-edges allowed]} \\ \ell_1 & \text{[No multi-edges]} \end{cases}$ $\mathcal{B} = \mathcal{B}_R + \mathcal{B}_M + \mathcal{B}_T - \mathcal{B}_{R-M} - \mathcal{B}_{R-T} - \mathcal{B}_{M-T} - \mathcal{B}_{R-T} + \mathcal{B}_{R-T}$ $\mathcal{B}_R = \mathcal{R} \circ_e (\mathcal{D} - \mathcal{S})$ $\mathcal{B}_M = \begin{cases} \mathcal{M} \circ_e (\mathcal{D} - \mathcal{P}) = (v^2 * \text{Set}_{\geq 3}(\mathcal{D} - \mathcal{P})) / \bullet \bullet \bullet, & \text{[Multi-edges allowed]} \\ (v^2 * e * \text{Set}_{\geq 2}(\mathcal{D} - \mathcal{P} - e) + v^2 * \text{Set}_{\geq 3}(\mathcal{D} - \mathcal{P} - e)) / \bullet \bullet \bullet, & \text{[No multi-edges]} \end{cases}$ $\mathcal{B}_T = \overline{\mathcal{G}}_3 \circ_e \mathcal{D}$ $\mathcal{B}_{R-M} = (v^2 * \mathcal{S} * \mathcal{P}) / \bullet \bullet \bullet, \quad [\text{Up to pole exchange, denoted by } / \bullet \bullet \bullet]$ $\mathcal{B}_{R-T} = (v^2 * \mathcal{S} * \mathcal{H}) / \bullet \bullet \bullet$ $\mathcal{B}_{M-T} = (v^2 * \mathcal{P} * \mathcal{H}) / \bullet \bullet \bullet$ $\mathcal{B}_{T-T} = (v^2 * \mathcal{H} * \mathcal{H}) / \bullet \bullet \bullet$ $\mathcal{B}_{R-T} = (v^2 * \mathcal{H} * \mathcal{H}) / (\bullet \bullet \bullet, H \equiv H), \quad [\text{Up to pole and component exchange}]$ (iii) Vertex-pointed 2-Connected $\mathcal{G}'_2 = e' + \mathcal{B}', \quad e' = \begin{cases} \ell'_1 + \ell'_2 = v * (e + \text{Set}_2(e)) & \text{[Multi-edges allowed]} \\ \ell'_1 = v * e & \text{[No multi-edges]} \end{cases}$ $\mathcal{B}' = \mathcal{V} - \mathcal{V}_R + \mathcal{V}_M + \mathcal{V}_T - \mathcal{V}_{R-M} - \mathcal{V}_{R-T} - \mathcal{V}_{M-T} - \mathcal{V}_{R-T} + \mathcal{V}_{R-T}$ $\mathcal{V}_R = \mathcal{R}' \circ_e (\mathcal{D} - \mathcal{S})$ $\mathcal{V}_M = \begin{cases} v * \text{Set}_{\geq 3}(\mathcal{D} - \mathcal{P}), & \text{[Multi-edges allowed]} \\ v * e * \text{Set}_{\geq 2}(\mathcal{D} - \mathcal{P} - e) + v * \text{Set}_{\geq 3}(\mathcal{D} - \mathcal{P} - e), & \text{[No multi-edges]} \end{cases}$ $\mathcal{V}_T = \overline{\mathcal{G}}'_2 \circ_e \mathcal{D}$ $\mathcal{V}_{R-M} = v * \mathcal{S} * \mathcal{D}$ $\mathcal{V}_{R-T} = v * \mathcal{S} * \mathcal{H}$ $\mathcal{V}_{M-T} = v * \mathcal{P} * \mathcal{H}$ $\mathcal{V}_{T-T} = v * \mathcal{H} * \mathcal{H}$ $\mathcal{V}_{R-T} = (v * \mathcal{H} * \mathcal{H}) / H \equiv H$

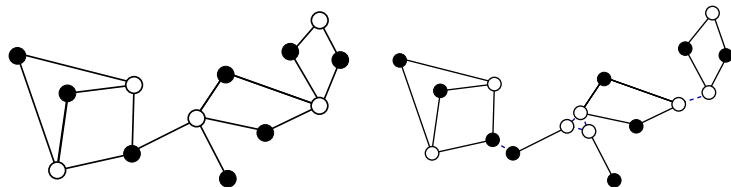
This system is obtained applying the **dissymmetry theorem for trees** in an ingenious way.

The key step is the one which translates combinatorially the integration!

Bipartite Graphs: the strategy (I)

Can we apply the same decomposition for bipartite graphs?

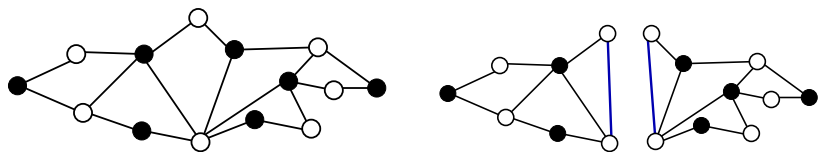
1-sums are easy!



The 2-connected components are also bipartite

Bipartite graphs: the strategy (II)

For 2-sums we have problems



We need to study something more general **Ising Model**.

Bipartite graphs: the strategy (III)

PROBLEM: going from 3-connected level to 2-connected level.

Networks in the general case:

$$\begin{cases} D(x, y) = y + S(x, y) + P(x, y) + H(x, y) \\ S(x, y) = D(x, y)x (D(x, y) - S(x, y)) \\ P(x, y) = (1 + y) (\exp(S(x, y) + H(x, y)) - 1 - S(x, y) - H(x, y)) \\ H(x, y) = \frac{2}{x^2} T_y(x, D(x, y)). \end{cases}$$

Networks in the Ising model:

$$\begin{cases} S_{\circ-\bullet} = x D_{\circ-\bullet} \frac{x(D_{\circ-\circ}^2 - D_{\circ-\bullet}^2) + 2D_{\circ-\circ}}{(1+x(D_{\circ-\circ} + D_{\circ-\bullet})) (1+x(D_{\circ-\circ} - D_{\circ-\bullet}))} \\ S_{\circ-\circ} = x \frac{D_{\circ-\circ}^2 + D_{\circ-\bullet}^2 + D_{\circ-\circ}^3 - x D_{\circ-\circ} D_{\circ-\bullet}^2}{(1+x(D_{\circ-\circ} + D_{\circ-\bullet})) (1+x(D_{\circ-\circ} - D_{\circ-\bullet}))} \end{cases}$$

and

$$2(1 + y_{\circ-\bullet}) \frac{\partial}{\partial y_{\circ-\bullet}} B(x, y_{\circ-\bullet}, y_{\circ-\circ}) + 2(1 + y_{\circ-\circ}) \frac{\partial}{\partial y_{\circ-\circ}} B(x, y_{\circ-\bullet}, y_{\circ-\circ}) = x^2(1 + D_{\circ-\circ} + D_{\circ-\bullet})$$

We do not have any choice: Combinatorial Integration!

The Program (Coming soon!)

One needs to rephrase the grammar for graphs including the colours.

Once we have this (+ Singularity analysis), we can:

- i.- Study SP-graphs.
- ii.- Study families of graphs defined by “easy” 3-connected components.
- ii.- Study limit laws for several parameters

What we **CANNOT** do (for the moment!):

**STUDY GENERAL
PLANAR BIPARTITE PLANAR GRAPHS**



**OBTAIN GF FOR 3-CONNECTED MAPS
(BERNARDI & BOUSQUET-MÉLOU)**

Merçi



Bipartite subfamilies of planar graphs

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